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# A comparative study of image coding techniques : filter banks vs. discrete cosine transform

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### Abstract

Title of Thesis: A Comparative Study of Image Coding Techniques: Filter Banks vs. Discrete Cosine Transform.

Hosam Mutlag, Master of Science, 1991

Thesis directed by: Dr. Ali N. Akansu

Subband coding of still image frames using Binomial PR-QMF has been presented in this thesis. Simulation results have shown that the performance of subband coding using a low complexity coder is practically the same as the performance of the industry standard ( $8 \times 8$ ) DCT based image coder.

A low bit rate adaptive video coding technique is also introduced in this thesis. The redundancy within adjacent video frames is exploited by motion compensated interframe prediction. The Motion Compensated Frame Difference (MCFD) signals are filtered by employing Binomial PR-QMF structure into four subbands. Then, the subbands are quantized using an efficient Motion Based Adaptive Vector Quantization (MBAVQ) algorithm. Here, the adaptation scheme is based on block motion vectors rather than local signal energy which was used in earlier works of several researchers. The new technique results in a reduction in bit rate by nearly (40%) due to the drop of the extra bits used for local variances. Moreover, for the video test sequances considered, MBAVQ method gives superior SNR results over local Variance Based Adaptive Vector Quantiztion (VBAVQ) scheme especially for high motion frames. A Comparative Study of Image Coding Techniques: Filter Banks vs. Discrete Cosine Transform

> by Hosam Fawzi Mutlag

Thesis submitted to the Faculty of the Graduate School of the New Jersey Institute of Technology in partial fulfillment of the requirements for the degree of Master of Science in Electrical Engineering

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# Chapter 1 Introduction

Typical video has spatial resolution of approximately  $512 \times 512$  pixels per frame. At 8 bits per pixel per color channel and 30 frames per second, this translates into a rate of nearly  $180 \times 10^6$  bits/s. The large channel capacity and memory requirements for digital image transmission and storage makes it desirable to consider data compression techniques. Image data compression is concerned with minimizing the number of bits required to represent an image with an acceptable visual quality. Compression can be achieved by transforming the signal, projecting it on a basis of functions, and then encoding the transform coefficients. These transforms vary from the conventional block transforms to ideal subband filter banks. The principle of subband coding has recently been successfully applied to data compression of both still images and video. In subband coding, the signal to be coded is decomposed into narrow band pass signals (subbands). Each subband is then accordingly subsampled and encoded with a bit rate matched to the signal statistics in that subband.

Subband coding was first introduced by Crochiere, et. al.[1] in speech coding and then extended to multidimensional signals by Vetterli[2]. Then, this concept was applied to the coding of images [3] [4]. Perfect Reconstruction Quadrature Mirror Filters (PR-QMF) have been proposed as structures suitable for subband coding [5][6]. These filters employing a tree decomposition structure provide a basis for a multiresolution signal representation [7] [8]. Recently, discrete wavelet transforms have been proposed as a new approach for multi-resolution signal decomposition. More recently, Binomial QMF-Wavelet Transform has been proposed for multiresolution signal decomposition [10] [11].

In this thesis, we have studied the performance of Binomial QMF in subband coding with comparison to the Discrete Cosine Transform (DCT). Simulation programs were developed for a complete subband coding system (codec) consisting of the following four distinct parts:

- An analysis filter bank splitting the input signal into subbands.
- An encoder which consists of:
  - 1. Bit allocation algorithm.
  - 2. Quantizers.
  - 3. Huffman encoder.
- A decoder whose purpose is to produce an approximation to the original subband signals.
- The synthesis filter bank that combines the decoded subband signals, to reconstruct the received signal.

A similar DCT based codec is also developed for comparison purposes.

We also propose a new Motion Based Adaptive Vector Quantization (MBAVQ) approach for subband video coding. In video signals, interframe images have significant frame to frame redundancy. For that reason, in video coding techniques, video frames are motion compensated using an efficient search algorithm to remove temporal redundancies. The resulting prediction error signals, Motion Compensated Frame Difference (MCFD), are coded using an efficient coding scheme. In our video coding scheme, the MCFD signals are divided into four subbands. Least significant band (highest frequency), is neglected and the other bands are vector quantized adaptively. This thesis is organized as follows. Chapter 2 describes the Binomial-QMF which have been used. In the following chapter, Discrete Cosine Transform is discussed. In chapter 4, optimum bit allocation and gain of transform coding over PCM is explained. Quantization is reviewed in chapter 5. The next chapter discusses source entropy and Huffman coding. In chapter 7, Motion Based Adaptive Vector Quantization (MBA-VQ) video coding scheme is introduced. Also, a comprehensive set of simulation results is presented in this chapter. Finally, we conclude the thesis with a discussion of the results and future research.

# Chapter 2 PR\_QMF

#### 2.1 Theoretical Derivations

The QMF bank is a multirate digital filter bank. There are decimators in the system which down-sample a signal sequence, and there are interpolators which perform upsampling. The input-output relation for a two-fold decimator can be written in the transform domain as

$$Y(e^{jw}) = \frac{1}{2} \left[ X(e^{jw/2}) + X(-e^{jw/2}) \right]$$
(2.1)

where  $Y(e^{jw})$  has a period of  $2\pi$ . The effect of compression in time domain is an expansion (or stretching) in frequency domain. The transform domain relation for a two-fold interpolator is

$$Y(e^{jw}) = X(e^{jw^2})$$
  

$$Y(z) = X(z^2)$$
(2.2)

The interpolator causes compression in the frequency domain.

Let us consider Figure (2.1) in which a two-channel QMF system is shown. Based on relations (2.1) and (2.2), it is possible to express  $\hat{X}(z)$  in Figure 1 as

$$\hat{X}(z) = \frac{1}{2} [H_0(z)F_0(z) + H_1(z)F_1(z)]X(z) + \frac{1}{2} [H_0(-z)F_0(z) + H_1(-z)F_1(z)]X(-z)$$
(2.3)

The reconstructed signal can in general be subject to three types of distortions:

- Aliasing distortion.
- Phase distortion.
- Amplitude distortion.

The second term in (2.3) represents the effects of aliasing and imaging. This term can be made to drop simply by choosing the synthesis filters to be

$$F_0(z) = -H_1(-z) (2.4)$$

$$F_1(z) = H_0(-z) (2.5)$$

The QMF bank becomes a linear and time-invariant system with transfer function

$$T(z) = \frac{\hat{X}(z)}{X(z)} = \frac{1}{2} \left[ H_0(z) H_1(-z) - H_1(z) H_0(-z) \right]$$
(2.6)

If  $|T(e^{jw})| = \text{constant}$  for all w, then there is no amplitude distortion. Also, if T(z) is a linear-phase FIR function, then  $\arg[T(e^{jw})] = kw$ , and there is no phase distortion. This means T(z) is a delay, i.e.,  $T(z) = Cz^{-n_0}$ , so that the reconstructed signal  $\hat{x}(n)$ is a delayed version of x(n).

Smith and Barnwell [6] have shown that one can simultaneously eliminate both amplitude and phase distortions by choosing (2N - 1) odd and

$$H_1(z) = z^{(2N-1)} H_0(-z^{-1})$$
(2.7)

Therefore,

$$T(z) = \frac{1}{2} z^{-(2N-1)} \left[ H_0(z) H_0(z^{-1}) - H_0(-z) H_0(-z^{-1}) \right] = C z^{-(2N-1)}$$
(2.8)

The perfect reconstruction requirement of a two-band QMF reduces to[5]

$$Q(z) = H(z)H(z^{-1}) + H(-z)H(-z^{-1}) = C$$
  
=  $R(z) + R(-z)$  (2.9)

where H(z) is a low-pass filter of length 2N.

The PR requirement, Eq.(2.9), can be readily recast in an alternate, time domain form [10]. First, one notes that R(z) is a spectral density function and hence is representable by a finite series of the form

$$R(z) = \gamma_{2N-1} z^{2N-1} + \gamma_{2N-2} z^{2N-2} + \ldots + \gamma_0 z^0 + \ldots + \gamma_{2N-1} z^{-(2N-1)}$$
(2.10)

Then

$$R(-z) = -\gamma_{2N-1}z^{2N-1} + \gamma_{2N-2}z^{2N-2} - \ldots + \gamma_0 z^0 - \gamma_1 z^{-1} \ldots - \gamma_{2N-1} z^{-(2N-1)}$$
(2.11)

Therefore Q(z) consists only of even powers of z. To force Q(z) = C which is a constant, it suffices to make all even indexed coefficients in R(z) equal to zero, except for n=0. However, the  $\gamma_n$  coefficients in R(z) are simply the samples of the autocorrelation  $\rho(n)$  given by

$$\rho(n) = \sum_{k=0}^{2N-1} h(k)h(k+n) = \rho(-n)$$

$$\stackrel{\text{def}}{=} h(n) \odot h(n)$$
(2.12)

where  $\odot$  indicates a correlation operation. This follows from the z-transform relationships

$$R(z) = H(z)H(z^{-1}) \longleftrightarrow h(n) * h(-n) = \rho(n)$$
(2.13)

where  $\rho(n)$  is the convolution of h(n) with h(-n), or equivalently, the time autocorrelation, Eq.(2.12). Hence, we need to set  $\rho(n) = 0$  for n even, and  $n \neq 0$ . Therefore,

$$\rho(2n) = \sum_{k=0}^{2N-1} h(k)h(k+2n) = 0, \qquad n \neq 0$$
(2.14)

If the normalization is imposed,

$$\sum_{k=0}^{2N-1} |h(k)|^2 = 1$$
(2.15)

one obtains the PR requirement in time

$$\sum_{k=0}^{2N-1} h(k)h(k+2n) = \delta_n \tag{2.16}$$

### 2.2 The Binomial Family

The binomial family of orthogonal sequences [13] [14] is generated by successive differencing of the binomial sequence, which is defined on the finite interval [0, N] by

$$x_{0}(k) = \begin{cases} \begin{pmatrix} N \\ k \end{pmatrix} = \frac{N!}{(N-k)!k!} & 0 \le k \le N \\ 0 & \text{otherwise} \end{cases}$$
(2.17)

The other members of the binomial family are obtained from

$$x_r(k) = \nabla^r \left( \begin{array}{c} N-r\\ k \end{array} \right) \quad r = 0, 1, \dots, N \tag{2.18}$$

where

$$\nabla f(n) = f(n) - f(n-1)$$

this is the backward difference operator. Taking successive differences yields Bonomial function family of length (N+1)

$$x_r(k) = \binom{N}{k} \sum_{\nu=0}^r (-2)^{\nu} \binom{r}{\nu} \frac{k^{(\nu)}}{N^{(\nu)}}$$
(2.19)

where  $k^{(\nu)}$  is the forward factorial function, a polynomial in k of degree  $\nu$ 

$$k^{(\nu)} = \begin{cases} k(k-1)\dots(k-\nu+1) & \nu \ge 1\\ 1 & \nu = 1 \end{cases}$$
(2.20)

This family of binomially-weighted polynomials has a number of properties. Taking z transform, we obtain

$$X_0(z) = (1+z^{-1})^N$$
  

$$X_r(z) = (1-z^{-1})^r (1+z^{-1})^{N-r}$$
(2.21)

The binomial matrix X is the  $(N+1) \times (N+1)$  matrix

$$X = [x_r(k)]$$

where  $x_r(k)$  is the entry in the  $r^{th}$  row and  $k^{th}$  column. The salient property of this matrix is that the rows are orthogonal to the columns,

$$\sum_{k=0}^{N} x_r(k) x_k(s) = (2)^N \delta_{r-s}$$
(2.22)

or

$$X^2 = (2)^N I (2.23)$$

Additionally, the Bionomial filters are linear phase quadrature mirror filters. From eq. (2.21), we see that

$$X_r(-z) = X_{N-r}(z)$$
 (2.24)

which implies

$$(-1)^{k} x_{r}(k) = x_{N-r}(k) \quad r = 0, 1, \cdots, N$$
(2.25)

Also,

$$z^{-N}X_r(z^{-1}) = (-1)^r X_r(z)$$
(2.26)

implies

$$x_r(N-k) = (-1)^r x_r(k)$$
(2.27)

Equations (2.25) and (2.27) demonstrate the symmetry and asymmetry of the rows and columns of the Bionomial matrix X. From equation (2.25), we can infer that the complementary filters  $X_r$  and  $X_{N-r}(z)$  have magnitude responses which are mirror images about  $\omega = \pi/2$ 

$$|X_r(\exp j(\frac{\pi}{2} - \omega))| = |X_{N-r}(\exp j(\frac{\pi}{2} + \omega))|$$
(2.28)

Moreover, the cross-correlation of the sequence  $x_r(n)$ , and  $x_s(n)$  is defined as

$$\rho_{rs} = x_r(n) * x_s(-n) = \sum_{k=0}^N x_r(k) x_s(n+k) \rightleftharpoons R_{rs}(z)$$
(2.29)

and

$$R_{rs}(z) = X_r(z^{-1})X_s(z)$$
(2.30)

Now for any real crosscorrelation,

$$\rho_{rs}(-n) = \rho(n) \qquad \forall s, r \qquad (2.31)$$

Also, the quadrature mirror property of Eq.(2.25) implies that

$$\rho_{rs}(n) = -\rho_{rs}(n) \qquad (s-r) \text{ is odd}$$
  
 $\rho_{rs}(n) = \rho_{rs}(n) \qquad (s-r) \text{ is even} \qquad (2.32)$ 

These properties are subsequently used in driving the perfect reconstruction Bionomial QMF in the next section.

#### 2.3 The Binomial QMF

Now, it is a straight forward matter to impose PR condition of Eq.(2.16) on the binomial family [10]. First, the half-band filter is

$$h(n) = \sum_{r=0}^{\frac{N-1}{2}} \theta_r x_r(n)$$

or in the z transform

$$H(z) = \sum_{r=0}^{\frac{N-1}{2}} \theta_r (1+z^{-1})^{N-r} (1-z^{-1})^r = (1+z^{-1})^{(N+1)/2} F(z)$$
(2.33)

where F(z) is FIR filter of order (N-1)/2. For convenience, take  $\theta_0 = 1$ , and later impose the normalization of Eq.(2.15). Substituting (2.33) into (2.12) gives

$$\rho(n) = \left(\sum_{r=0}^{\frac{N-1}{2}} \theta_r x_r(n)\right) \odot \sum_{s=0}^{\frac{N-1}{2}} \theta_s x_s(n) \\
= \sum_{r=0}^{\frac{N-1}{2}} \sum_{s=0}^{\frac{N-1}{2}} \theta_r \theta_s [x_r(n) \odot x_s(n)] \\
= \sum_{r=0}^{\frac{N-1}{2}} \sum_{s=0}^{\frac{N-1}{2}} \theta_r \theta_s \rho_{rs}(n) \\
= \sum_{r=0}^{\frac{N-1}{2}} \theta_r^2 \rho_{rr}(n) + \sum_{r=0, r \neq s}^{\frac{N-1}{2}} \sum_{s=0}^{\frac{N-1}{2}} \theta_r \theta_s \rho_{rs}(n) \qquad (2.34)$$

Eq.(2.32) implies that the second summation in Eq.(2.34) has only terms where the indices differ by an even integer. Therefore the autocorrelation for the binomial half-band low-pass filter is

$$\rho(n) = \sum_{n=0}^{\frac{N-1}{2}} \theta_r^2 \rho_{rr}(n) + 2 \sum_{l=1}^{\frac{N-3}{2}} \sum_{\nu=0}^{\frac{N-1}{2}-2l} \theta_{\nu} \theta_{\nu+2l} \rho_{\nu,\nu+2l}(n)$$
(2.35)

Finally, the PR requirement is

$$\rho(n) = 0, \qquad n = 2, 4, \dots, N-1$$
(2.36)

This condition gives a set of  $\frac{N-1}{2}$  nonlinear algebraic equations, in the  $\frac{N-1}{2}$  unknowns  $\theta_1, \theta_2, \ldots, \theta_{\frac{N-1}{2}}$ . From these, we can obtain the corresponding Binomial PR-QMFs.

These filters can be implemented using either the purely FIR structure, or the pole-zero cancellation configuration. The latter is shown in Fig.2 for N = 5. Wherein both low-pass and high-pass filters are simultaneously realized. Coefficient  $\theta_0$  can be taken equal to unity, leaving only  $\theta_1$  and  $\theta_2$  as tap weights. These are the only multiplications needed when using the Binomial network as the half-band QMF rather than the six h(n) weights in a transversal structure.

#### 2.4 M-Band Tree Decomposition

After a given signal x(t) is sampled at  $f_s$  to give a signal X(n), and split it into two signals  $X_L(n)$  and  $X_H(n)$ , with the reduction of the sampling rate to  $f_s/2$ , the decomposition can be extended to more than two subbands by processing these signals in the same manner as the initial signal X(n). Four signals are thus obtained with a reduction of the sampling rate to  $f_s/4$ . The decomposition/reconstruction can be generalized by repeating n times the previously described splitting using n-stage tree decomposition. In figure (2.3), a four subband decomposition is shown. The signal X(n) is split into four signals  $X_{LL}$ ,  $X_{LH}$ ,  $X_{HL}$ , and  $X_{HH}$ .

### 2.5 Two Dimensional Separable Case

The extention of the QMF decomposition/reconstruction concept to two dimensional case is relatively straight forward in separable filter case. The conditions required for splitting two dimensional signals into more than two bands, with alias-free recon-



Figure 2.1: Two Channel QMF Bank.



Figure 2.2: Low-pass and high-pass QMF filters from Binomial Network.



Figure 2.3: Four-band tree decomposition for one dimensional signal X(n)

struction of the original signal, is the separability feature of the filters, that is

$$h(n_1, n_2) = h_1(n_1)h_2(n_2) \tag{2.37}$$

These separability characteristics of the filter provide an alternative method of implementation for 2D-QMF banks. The analysis and synthesis are done as shown in figure (2.4) and figure (2.5) where the structure consists of one-dimensional filters. The computation is performed first along one axis (rows) and then along the other axis (columns). It can be shown that application of this filter structure will permit an alias-free reconstruction of the input signal at the receiver. The detailed analysis/synthesis and perfect reconstruction are given in [3]. Figure (2.4) shows four-band tree decomposition of a two-dimensional signal X(n,m) while figure (2.5) shows the construction tree.



Figure 2.4: Four-band tree decomposition of a two-dimensional signal X(n,m)



Figure 2.5: Four-band reconstruction of a two-dimensional signal X(n,m)

# Chapter 3 Discrete Cosine Transform

The term image transform usually refers to a class of unitary matrices used for representing images. Just as one-dimensional signal can be represented by an orthogonal series of basis functions, an image can also be expanded in terms of a discrete set of basis arrays called basis images. These basis images can be generated by unitary matrices. To make transform coding practical, a given image is divided into small rectangular blocks, and each block is transform coded independently. For an  $N \times M$ image divided into NM/pq blocks, each of size  $p \times q$ , the main storage requirements for implementing the transform are reduced by afactor of MN/pq. The computational load is reduced by a factor of  $\log_2 MN/\log_2 pq$  for fast transform requiring  $\propto N \log_2$ operations to transform an  $N \times 1$  vector. For 256  $\times$  256 images divided into 8  $\times$  8 blocks, these factors are 1024 and 2.66, respectively. Although the operation count is not greatly reduced, the complexity of the hardware for implementing small size transform is reduced significantly. However, smaller block sizes yield lower compression. The  $N \times N$  discrete cosine transform (DCT) matrix  $C = \{c(k, n)\}$ , is defined as[16]

$$c(k,n) = \begin{cases} \frac{1}{\sqrt{N}} & k = 0, 0 \le n \le N - 1\\ \sqrt{\frac{2}{N}} \cos \frac{\pi(2n+1)}{2N} & 1 \le k \le N - 1, 0 \le n \le N - 1 \end{cases}$$
(3.1)

The two-dimensional discrete cosine transform of a two-dimensional sequence

 $\{u(m,n), 0 \leq m \leq N-1, 0 \leq n \leq N-1\}$  , is obtained by

$$v(k,l) = \sum_{m=0}^{N-1} \sum_{n=0}^{N-1} c(k,m)u(m,n)c(l,n)$$
(3.2)

The inverse transformation is given by

$$u(m,n) = \sum_{k=0}^{N-1} \sum_{l=0}^{N-1} c(k,m)v(k,l)c(l,n)$$
(3.3)

or in vector notation

$$V = CUC^T \tag{3.4}$$

and

$$U = C^T V C \tag{3.5}$$

The cosine transform is real and orthogonal, that is,

$$C = C^* \to C^{-1} = C^T \tag{3.6}$$

To implement one-dimensional DCT of size N using fast transform algorithms, it requires  $O(N \log N)$  operations, where one operation is a real multiplication and a real addition. For two-dimensional DCT of size  $N \times N$ , the transformation is equivalent to 2N one-dimensional DCT. Hence, the total number of operation is  $O(N^2 \log_2 N)$ .

## Chapter 4

## Optimum Bit Allocation and Gain of Transform Coding

### 4.1 Bit Allocation

Bit allocation is a major concern in any coding scheme where a given quota of bits must be efficiently distributed among a number of different quantizers. Encoding in subbands offers several advantages. By appropriately allocating the bits in different bands, the number of quantizer levels and hence reconstruction error variance can be separateley controlled in each band and the shape of the overall reconstruction error spectrum can be controlled as a function of frequency. Using this approach, the noise spectrum can be shaped according to the subjective noise perception of the human eye. Moreover, the error in coding a subband is confined to that subband.

The overall MSE incurred in an orthonormal, equal bandwidth, subband coding scheme with K subbands is given by

$$D = \sum_{i=1}^{K} D_i(r_i),$$
(4.1)

where  $D_i(r_i)$  is the distortion rate performance of the encoder operating on the *i*th subband at  $r_i$  bits/sample. The average bit rate is given by

$$R = \frac{1}{K} \sum_{i=1}^{K} r_i \tag{4.2}$$

In equation (4.2) we have assumed that all subbands have the same number of pixels.

If this is not the case, the  $r_i$ 's should be multiplied by appropriate weighting coefficients to account for the variability in the number of pixels in any subband. The bit allocation we have used is based on the Lagrange multiplier technique in which it is assumed that all bands have the same pdf type and the distortion rate performance of the quantizers are given by

$$D_i(r_i) = \epsilon_i^2 \sigma_i^2 2^{-2r_i} \tag{4.3}$$

where,  $\sigma_i^2$  is the variance of the *i*th subband and  $\epsilon_i^2$  is the quantization correction factor for that band. If we assume equal band distortions  $D_i = D_j$ , the following bit allocation is obtained [16]

$$r_{i} = r_{ave} + \frac{1}{2}\log_{2}\left[\frac{\sigma_{i}^{2}}{\left[\prod_{i=1}^{K}\sigma_{i}^{2}\right]^{\frac{1}{K}}}\right]$$
(4.4)

where  $r_{ave}$  is the average bit rate.

 $\{r_i\}$  are not restricted to be non-negative here. In practice, they are truncated to zero if they become negative. A negative bit allocation result implies that if that band is completely discarded its reconstruction error contribution is still less than the corresponding distortion for the given rate. Equation (4.4) holds only for regular tree structures of subband coding. For irregular tree structure, with  $N_1$  bands in the first level of the tree and only band p is decomposed further into  $N_2$  bands in the second level of the tree, the corresponding optimum bit allocation expressions are found as [18]

$$r_{k1} = r_{ave} + \frac{1}{2} \log_2 \frac{\sigma_{k1}^2}{\left[\prod_{k1=1,k1\neq p}^{N_1} \sigma_{k1}^2\right]^{\frac{1}{N_1}} \left[\prod_{k2=1}^{N_2} \sigma_{pk2}^2\right]^{\frac{1}{N_1N_2}}}$$
(4.5)

and,

$$r_{pk2} = r_{ave} + \frac{1}{2} \log_2 \frac{\sigma_{pk2}^2}{\left[\prod_{k1=1,k1\neq p}^{N_1} \sigma_{k1}^2\right]^{\frac{1}{N_1}} \left[\prod_{k2=1}^{N_2} \sigma_{pk2}^2\right]^{\frac{1}{N_1N_2}}}$$
(4.6)

#### 4.2 Gain of Transform Coding

An N band orthonormal transform implies the variance preservation condition:

$$\sigma_x^2 = \frac{1}{N} \sum_{k=0}^{N-1} \sigma_k^2 \tag{4.7}$$

where  $\sigma_x^2$  is the input signal variance with zero mean and  $\sigma_k^2$  are the band variances. All orthonormal, variance preserving, signal decomposition techniques can be evaluated by employing the energy compaction criterion, namely the gain of transform coding over pulse code modulation  $G_{TC}$  [16]. If we assume that all the bands and the input signal have the same pdf type, the distortion ratio of PCM over transform coding at the same bit rate can be easily obtained as

$$G_{TC} = \frac{D_{PCM}}{D_{TC}} = \frac{\sigma_x^2}{\left[\prod_{k=1}^N \sigma_k^2\right]^{\frac{1}{N}}}$$
(4.8)

This equation holds for regular binary tree structures of subband technique. For irregular tree case equation (4.8) can be modified properly. Assuming an  $N_1$  band in the first level of the tree, and only band p is decomposed further into  $N_2$  bands in the second level of the tree, the corresponding gain is [18]

$$G_{TC} = \frac{\sigma_x^2}{\left(\left[\prod_{k_1=1, k_1 \neq p}^{N_1} \sigma_{k_1}^2\right] \left[\prod_{k_2=1}^{N_2} \sigma_{pk_2}^2\right]^{\frac{1}{N_2}}\right)^{\frac{1}{N_1}}}$$
(4.9)

# Chapter 5 Quantization

#### 5.1 Introduction

The subsequent step is quantization. Quantization is the most common form of data compression and is fundamental to any digitization scheme or data compression system. Quantization causes entropy reduction and hence it is an irreversible operation. A quantizer maps a continuous variable u into a discrete  $\dot{u}$ , which takes values from a finite set  $\{r_1, \dots, r_L\}$  of numbers. This mapping is generally a stair case function and the quantization rule is as follows:

- Define  $\{t_k, k = 1, \dots, L+1\}$  as a set of increasing transition or decision levels with  $t_1$  and  $t_{L+1}$  as the minimum and maximum values, respectively, of u.
- If u lies in interval  $[t_k, t_{k+1})$ , then it is mapped to  $r_k$ , the kth reconstruction level.

We will consider only zero memory quantizers, which operate on one input sample at a time, and the output value depends only on that input. Each quantizer is operating at a different rate by using different quantization tables. The rate is usually measured by

$$R_k = \log_2 N_k \tag{5.1}$$

where  $R_k$  and  $N_k$  are the rate and the size of the table for the kth quantizer respectively. In zero memory quantizer (scalar quantization), the quantization is performed such that quantization noise is minimized, with a given number of quantization levels.

### 5.2 The Optimum Mean Square or Lloyd-Max Quantizer

This quantizer minimizes the mean square error for a given number of quantization levels. Let u be a real scalar random variable with a continuous probability density function  $p_u(u)$ . It is desired to find the decision level  $t_k$  and the reconstruction levels  $r_k$  for an L level quantizer such that the mean square error [17]

$$\mathcal{E} = E[(u - \dot{u})^2] = \int_{t_1}^{t_{L+1}} (u - \dot{u})^2 p_u(u) du$$
(5.2)

is minimized. Rewriting this as

$$\mathcal{E} = \sum_{i=1}^{L} \int_{t_i}^{t_{i+1}} (u - r_i)^2 p_u(u) du$$
 (5.3)

By differentiating with respect to  $t_k$  and  $r_k$  and equating the result to zero, we get

$$\frac{\partial \mathcal{E}}{\partial t_k} = (t_k - r_{k-1})^2 p_u(t_k) - (t_k - r_k)^2 p_u(t_k) = 0$$
(5.4)

$$\frac{\partial \mathcal{E}}{\partial r_k} = 2 \int_{t_k}^{t_{k+1}} (u - r_k) p_u(u) du = 0 \quad 1 \le k \le L$$
(5.5)

Using the fact that  $t_{k-1} \leq t_k$ , this gives

$$t_k = \frac{(r_k + r_{k-1})}{2} \tag{5.6}$$

$$r_{k} = \frac{\int_{t_{k}}^{t_{k+1}} u p_{u}(u) du}{\int_{t_{k}}^{t_{k+1}} p_{u}(u) du} = E[u|u \epsilon \mathcal{J}_{k}]$$
(5.7)

where  $\mathcal{J}_k$  is the *k*th interval  $[t_k, tk + 1)$ . These results state that the optimum transition levels lie halfway between the optimum reconstruction levels, which, in turn, lie at the center of mass of the probability density in between the transition levels. Both (5.6) and (5.7) are nonlinear equations that have to be solved simultaneously. In practice, these equations can be solved by an iterative scheme such as the Newton method. A commonly used probability densities for quantization of image-related data is the Laplacian densities, which is defined as,

$$p_u = \frac{\alpha}{2} \exp(-\alpha |u - \mu|) \tag{5.8}$$

where  $\mu$  denotes the mean of u. The variance is given by

$$\sigma^2 = \frac{2}{\alpha} \tag{5.9}$$

#### 5.2.1 Properties of the Optimum Mean Square Quantizer

This quantizer has several interesting properties.

• The quantizer output is an unbiased estimate of the input, that is,

$$E[\dot{u}] = E[u] \tag{5.10}$$

• The quantization error is orthogonal to the quantizer output, that is,

$$E[(u - \dot{u})\dot{u}] = 0. \tag{5.11}$$

• The variance of the quantizer output is reduced by the factor 1 - f(B), where f(B) denotes the mean square distortion of the *B*-bit quantizer for unity variance inputs, that is,

$$\sigma_{\dot{u}}^2 = [1 - f(B)]\sigma_u^2 \tag{5.12}$$

• It is sufficient to design mean square quantizers for zero mean unity variance distributions.

### 5.3 Vector Quantization

When Shannon [20] proved that signals from an information source could be coded at a bit rate no greater than the entropy of the source, he showed that this would be
achieved, not by coding individual samples but by collecting samples into groups and encoding the groups. This is the basis of vector quantization.

A vector quantization scheme involves an encoder, a decoder and a codebook. The codebook is a lookup table with a k-bit address and  $2^k$  entries. Each entry in the table is a vector of samples. Both the encoder and the decoder have copies of the codebook. The encoder looks at each incoming vector of samples  $\bar{x}$  and selects the code-word  $\bar{y} = \bar{C}(i)$  which is the best match to  $\bar{x}$ . Formally, let  $d(\bar{x}, \bar{y})$  be the distortion measure. The Euclidean distance  $(\bar{x} - \bar{y})^T (\bar{x} - \bar{y})$  is frequently used as the measure but any positive definite quadratic form in  $(\bar{x} - \bar{y})$  will serve. The encoder transmits that vector index i for which  $d(\bar{x}, C(i))$  is minimum. The decoder receives i and presents C(i) as its output. Clearly the success of this technique depends on having a well chosen codebook. This codebook is, in general, not complete: i.e., the number of vectors in the codebook is finite while the number of possible vectors is usually, for all practical purposes, infinite. The distortion we must live with is that which arises from assigning an incoming vector to a codebook entry that does not quite match. The quantization process is reduced to a simple search and comparison procedure. The major difficulty in implementing vector quantizers is the time required to search a very large codebook. There are several algorithms for designing a codebook. We employed LBG algorithm which is summarized in the next subsection.

#### Lloyd's Algorithm

A procedure for designing codebooks was developed by Lloyd [21]. This algorithm was designed as a clustering technique for use in pattern recognition and related fields. It was extended to vector quantization by Linde, Buzo, and Gray [22] and therefore called LBG algorithm. The algorithm consists of an iterative technique for refining an initial codebook. The algorithm is as follows:

- 1. Encode a selection of test data with the current codebook and measure the average distortion. If the distortion is small enough the algorithm terminates.
- 2. For each address i in the codebook, find the centroid of all the input vectors which were mapped into i and make this centroid the new  $\bar{C}(i)$ . Go to step 1.

The test data used in this process are refered to as the training set. Each word in the codebook is used to represent a cluster of possible input vectors from this set.

# Chapter 6

# Source Entropy and Huffman Coding

### 6.1 Entropy

The entropy of a source with L symbols is defined as the average information generated by the source, i.e., Entropy

$$H = -\sum_{k=1}^{L} p_k \log_2 p_k \quad \text{bits/symbol} \tag{6.1}$$

where  $p_k$  is the probability of having kth symbol. For a digital image considered as a source of independent pixels, its entropy can be estimated from its histogram. For the given L levels, the entropy of a source is maximized for uniform distributions, i.e.,  $p_k = 1/L, k = 1, \dots, L$ . In that case

$$\max H = -\sum_{k=1}^{L} \frac{1}{L} \log_2 \frac{1}{L} = \log_2 L \quad \text{bits}$$
 (6.2)

The entropy of a source gives the lower bound on the number of bits required to encode its output. Obviously, this definition assumes a stationary source.

### 6.2 Entropy Coding: The Huffman Coding Algorithm

The quantizer output is generally coded by a fixed-length binary code word having B bits. If the quantized pixels are not uniformaly distributed, then their entropy will be

less than B, and there exists a code that uses less than B bits per pixel. In entropy, the goal is to encode a block of M pixels containing MB bits with probabilities  $p_i, i = 0, 1, \dots, L-1, L = 2^{MB}$ , by  $(-\log_2 p_i)$  bits, so that the average bit rate is

$$\sum_{i} p_i(-\log_2 p_i) = H \tag{6.3}$$

This gives a variabe-length code for each block, where highly probable blocks (or symbols) are represented by small length codes, and vice versa. If  $(-\log_2 p_i)$  is not an integer, the achieved rate exceeds H but approaches it asymptotically with increasing block size. For a given block size, a technique called *Huffman coding* is the most commonly used fixed to variable length encoding method [17].

#### 6.2.1 The Huffman Coding Algorithm

- 1. Arrange the symbol probabilities  $p_i$  in a decreasing order and consider them as leaf nodes of a tree.
- 2. While there is more than one node:
  - Merge the two nodes with smallest probability to form a new node whose probability is the sum of the two merged nodes.
  - Arbitrarily assign 1 and 0 to each pair of branches merging into a node.
- 3. Read sequentially from the root node to the leaf node where the symbol is located.

The preceding algorithm gives the *Huffman code book* for any given set of probabilities. Coding and decoding is done simply by looking up values in a table.

# Chapter 7 Experimental Studies

Since we are emphasizing the low complexity of the Binomial filter banks introduced in[10], it would be attractive to apply these filters in complete coder structures characterized by low overall complexity. Both still image coders and video coders are implemented and tested.

### 7.1 Subband Coding of Still Images

In a still image coder, the original image constitutes the input to the analysis filter bank. The subband signals are represented in a bit efficient manner using optimum bit allocation expressions given in equations (4.5) and (4.6). A problem with these two equations is that they may give negative values. Moreover, the  $r_{kj}$  may not match any of the existing quantizers bit rate. Actually, we need  $2^{r_{kj}}$  to be an integer equales to the number of levels in one of the used quantizers.

To solve these problems, subbands with negative  $r_{kj}$  are truncated to zero and dropped from our future calculations. Then the average bit rate  $r_{ave}$  is recalculated for the remaining bands. After that, a quickly converging iterative algorithm has been developed to find the integer number of levels corresponding to each of the remaining bands.

• First, we start with the band which has the highest frequency. The bit rate corresponding to this band is calculted using (4.5), (4.6) and the new average

bit rate  $r_{ave}$  recalculated previously.

- Then, choose the quantizer with the number of levels equals or less than  $2^{r_{k_j}}$ .
- Excessive bit rate resulting from this integer truncation is used to reevaluate average bit rate  $r_{ave}$  for the remaining bands and this band is dropped from our next calculations.
- This new average bit rate  $r_{ave}$  is used in equations (4.5) and (4.6) to calculate the number of levels for the next lower band by repeating the same preceding procedure.

We should notice that by assigning the excessive bits, resulting from truncation, to lower bands, we are putting more emphasis on low frequencies. This feature enhances the image quality because of the mechanism of the human visual system which is more sensitive to low frequencies. For different average bit rates and the "Lady" test image, this adaptive bit allocation results are shown in Table (7.1). These numbers reflect the iterative adjustment to remove negative and noninteger values.

Next step is to quantize each band using the corresponding number of levels found using the Bit Allocation Algorithm. The bands are quantized using normalized Laplacian quantizers [17] with the following numbers of levels:  $3, 5, 7, \dots, 29, 31, 33, 35, 64, 128$ . In this thesis, we have considered two methods of encoding the lowest frequency band; differential quantization such as DPCM and PCM quantization. In our experiments, we observed that the pixel to pixel correlation of upper band signals is very low. Consequently for these bands, PCM quantizer schemes are preferred to differential coding methods. We have therefore applied differential coding to the lowest band of the input signal in our experiments. In the DPCM encoder used here, a linear predictor constructs its predicted pixel as the weighted summation of previous pixels. Thus:

$$\bar{u}(m,n) = \frac{1}{4} [u(m-1,n) + u(m,n-1) + u(m-1,n-1) + u(m-1,n+1)] \quad (7.1)$$

The prediction error is then quantized by a symmetric non-uniform quantizer.

After quantization, the bands are entropy coded using Huffman coding scheme. To reduce complexity, only first order Huffman coding is used to code each quantized band.

Now, we can transmit the coded image over the channel. Here, we assume error free channel. It is suitable to use progressive transmission scheme for subband coding. The main objective of progressive transmission is to allow the receiver to recognize a picture as quickly as possible at minimum cost, by sending a low resolution level picture first. Further details of the picture are obtained by sequentially receiving the remaining encoded bands. At the receiver, the bands are decoded and reconstructed using the synthesizing filter bank.

A similar  $8 \times 8$  two-dimensional DCT based coder is also developed for comparison purposes.

The peak-to-peak signal to noise ratio is used as the objective performance criterion and defined as

$$SNR_{pp}(dB) = 10 \log_{10} \left[ \frac{255^2}{mse} \right]$$
 (7.2)

where mse is the mean square coding error.

### 7.2 Subband Video Coding with Motion Based Adaptive Vector Quantization

In our low complexity video coder, we employ block matching for motion compensation using Brute-force method. The block size is set to  $8 \times 8$  and the maximum displacement, both horizontally and vertically, is set to  $\pm 6$  pixels. The motion compensated frame difference (MCFD) signal is encoded with subband coding. The MCFD signal is split into four subbands, namely LL - HL - LH - HH bands, using 4-tap separable Binomial PR-QMF filters. The subbands are vector quantized adaptively after discarding the highest frequency band HH. Each  $4 \times 4$  block in the subbands corresponds to a motion vector of a block with size  $8 \times 8$  in the original resolution. From experiments, we have found that there is a relation between motion vector magnitude and prediction error variance of the corresponding block. In general, large magnitude motion vectors represent high variance blocks in MCFD while blocks with no motion have small variances. Using this relation, vector quantizers (codebook) were classified depending on the magnitude of block motion. The magnitude of block motion  $\hat{m}$  is defined as follow:

$$\hat{m} = \max(|i|, |j|) \tag{7.3}$$

where i and j are the horizontal and vertical motion respectively. We classified motion into three groups as follow:

- Group 1:  $\hat{m} = 1$  or 2.
- Group 2:  $\hat{m} = 3$  or 4.
- Group 3:  $\hat{m} = 5$  or 6.

Using blocks corresponding to these groups for each of the three bands, codebooks were generated using LBG algorithm. As a result of this, we got 9 codebooks each of 512 length. This technique were compared to the work done by Kadur in his thesis [23][19]. Kadur used local variances of each block in subbands to classify them. This means that we need to transmit extra bits for local variances for each  $4 \times 4$  block. The total bit rate for this technique can be written as

$$B = B_M + B_{SB} + B_G \tag{7.4}$$

where

- $B_M$  = average bits/pixel for motion information.
- $B_{SB}$  = average bits/pixel for subband signal.

•  $B_G$  = average bits/pixel for the transmission of local variances.

Using our MBAVQ scheme,  $B_G$  is not needed. This means a significant reduction in bit rate (40%) in general.



Figure 7.1: SNR versus entropy for 7 subband decomposition using Laplacian quantizers for all the subbands



Figure 7.2: SNR versus entropy for 10 subband decomposition using Laplacian quantizers for all the subbands



Figure 7.3: SNR versus entropy for 7 subband decomposition using DPCM for the low frequency band and Laplacian quantizers for the remaining subbands



Figure 7.4: SNR versus entropy for 7 subband decomposition using different quantization schemes for the lowest frequency subband



Figure 7.5: SNR versus entropy for DCT, 7 and 10 subband decomposition using Laplacian quantization scheme for the lowest frequency subband



Figure 7.6: SNR versus entropy for DCT and 10 subband decomposition using DPCM quantization scheme for the lowest frequency subband



Figure 7.7: Variance vs Motion for MCFD frames 10, 24 and 30 of CINDY



Figure 7.8: SNR vs frame index of MBAVQ and VBAVQ for CINDY



Figure 7.9: Entropy vs frame index of MBAVQ and VBAVQ for CINDY



Figure 7.10: Energy Compaction Gain,  $G_{TC}$  vs frame index for  $(8 \times 8)$  DCT and 4-tap Binomial QMF subband structure for MCFD of CINDY



Figure 7.11: SNR vs frame index of MBAVQ and VBAVQ for MONO



Figure 7.12: Entropy vs frame index of MBAVQ and VBAVQ for MONO



Figure 7.13: Motion vs frame index of MBAVQ and VBAVQ for MONO



Figure 7.14: SNR vs frame index of MBAVQ and VBAVQ for DUO



Figure 7.15: Entropy vs frame index of MBAVQ and VBAVQ for DUO



Figure 7.16: Motion vs frame index of MBAVQ and VBAVQ for DUO

	number of levels in the quantizer						
bit rate	band 1	band 2	band 3	band 4	band 5	band 6	band 7
0.5	35	3	7	0	0	0	0
0.75	64	5	11	3	0	0	0
1.0	64	7	15	5	0	3	0

Table 7.1: Results from bit allocation algorithm

# Chapter 8 Discussions and Future Research

Subband coding of digital images using Binomial PR-QMF has been presented for still images. Simulation results has shown that the performance of subband coding using a low complexity coder is practically the same as the performance of the industry standard  $(8 \times 8)$  DCT based codecs.

For video signals an efficient Motion Based Adaptive Vector Quantization (MBAVQ) subband coding method has been introduced. This new approach is compared with the Variance Based Adaptive Vector Quantization of subband video [19]. The new technique resulted in a reduction in bit rate by nearly (40%) due to the drop of the extra bits used for local variances. Moreover, this method gave superior SNR results especially for high motion frames. This clearly demonstrates that subband coding with MBAVQ should be considered as an atractive and powerful method for video coding. The modelling of MCFD signal and its relation with motion compensation techniques are open problems for future research.

## APPENDIX A

```
С
     SOURCE CODE FOR THE PROPOSED ADAPTIVE SUBBAND VIDEO CODING WI
С
     MOTION COMPENSATION using MBAVQ
С
    nx
            Row size of the picture
С
    ny
            Column size of the picture
С
    frame1 Previous Frame
    frame2 Current Frame
С
            Search frame from the previous frame
    pics
С
            Prediction of the current frame with motion compensati
С
    recon
С
    ibs
            Block size (8 is used)
            Assumed maximum displacement (Max of 8)
С
    ip
С
    frm2msk ibs*ibs size mask of the current frame to be
              motion compensated
С
    frm1msk ibs*ibs size mask of the previous frame in the
С
              same geometrical position (used for motion detection
С
            Motion Compensated Frame Difference Signal
С
    errl
С
            Search Region (ibs+ip)*(ibs+ip)
    searg
    mask
            Same as frm2msk
С
parameter(nx=400,ny=512)
      character*80 input_file
       common /ina/ input file
       real coff1(-20:20),coff2(-20:20),coff3(-20:20),coff4(-20:20)
       common /a/ coff1,coff2,coff3,coff4,ltap1,mtap1,ltap2,mtap2
    &
                ,ltap3,mtap3,ltap4,mtap4
       integer motionv(50,64), ifld
       common /motionv/ motionv
       common /ifld/ ifld
       real blv1(512,16),b1v2(512,16),b1v3(512,16)
       real b2v1(512,16),b2v2(512,16),b2v3(512,16)
       real b3v1(512,16),b3v2(512,16),b3v3(512,16)
       common /vqcodebook/ blv1,b1v2,b1v3,b2v1,b2v2,b2v3,b3v1,b3v2
    &
                          b3v3
       integer mcvector(50,64)
       common /bitrat/ entropy2,entropy
common /gtotal/ gtotal
С
         real frame1(nx,ny),frame2(nx,ny),pics(416,528),
       recon2(nx,ny),frmlmsk(8,8),frm2msk(8,8),errl(nx,ny),
    +
    +
       errtemp(nx,ny)
       integer ifrm1(nx,ny),ifrm2(nx,ny)
       character*1 pim(nx,ny)
```

character\*1 pim1(nx,ny)

common /AAA/ searg(24,24),mask(8,8)

```
С
```

С

call readf

c write(\*,\*) (coff1(i),i=ltap1,mtap1)

```
call readvq
```

```
c mfld: final field to be read
c ifld: starting field number
     mfld = 80
     ifld = 1
     call read frm(ifld,pim)
С
  Frame One is read into frame1 array
C****************
       do 10 i=1,nx
     do 10 j=1,ny
         ifrml(i,j)=ichar(pim(i,j))
         if(ifrm1(i,j).lt.0) ifrm1(i,j)=256+ifrm1(i,j)
         framel(i,j)=float(ifrm1(i,j))
 10
     continue
   call write_in_frm(ifld,frame1)
   call write out frm(ifld, framel)
     write(35,*) 'Original Image'
The loop to process mfld number of frames begins here
С
6000
     ifld = ifld+1
     write(*,*) 'Frame Number = ', ifld
      write(35,*) 'Frame Number =',ifld
     call read frm(ifld,pim1)
Current frame is read into frame2
С
```

```
do 11 i=1,nx
      do 11 j=1,ny
          ifrm2(i,j)=ichar(pim1(i,j))
          if(ifrm2(i,j).lt.0) ifrm2(i,j)=256+ifrm2(i,j)
          frame2(i,j)=float(ifrm2(i,j))
 11
      continue
   call write in frm(ifld, frame2)
      if(ifld.eq.11) then
С
      call write 1frm(frame2,'frame11')
С
С
      endif
Auto-Correlation, Mean, Variance are calculated in
C
    the subroutine autocor
C
c ip: displacement
      ip=6
     the mask block size
C
  ibs:
      ibs=8
  imthd: Enter 1 for Brute-force method and 2 for Orthogonal src
C
   imthd=1
  imdetect: Enter 1 if motion-detection is required'
С
   imdetect=1
Search Array pics is Initialized
do 100 i=1,nx+2*ip
     do 100 j=1,ny+2*ip
     pics(i,j)=0.0
100 continue
Search Array is generated from the previous frame.
C
С
  Borders are filled with first(or last) ip
  rows(or clums) of the previous frame
С
do 110 i=1,nx
     do 110 j=1,ny
     pics(i+ip,j+ip)=frame1(i,j)
110 continue
     do 111 i=1,ip
      do 111 j=1,ny
       pics(i,j)=frame1(i,j)
```

```
pics(i+nx+ip,j)=framel(i+nx-ip,j)
 111
      continue
      do 112 i=1,nx
       do 112 j=1,ip
         pics(i,j)=framel(i,j)
         pics(i,j+ny+ip)=framel(i,j+ny-ip)
112
      continue
Prediction of the Current frame is Initialized
do 240 i4=1,nx
       do 240 j4=1,ny
         recon2(i4, j4) = 0.0
240
       continue
The current frame is devided into 8*8 blocks and
C
С
   motion compensated. mcount keeps count of number
С
   of moving blocks.
mcount=0
      do 200 i=1,nx/ibs
     do 200 j=1,ny/ibs
      iact=(i-1)*ibs+1
      jact=(j-1)*ibs+1
   if (imdetect .eq. 1) then
       do 401 k=1,ibs
         do'401 l=1,ibs
           frm1msk(k,l)=framel(iact-1+k,jact-1+l)
           frm2msk(k,l)=frame2(iact-1+k,jact-1+l)
401
       continue
First the motion is detected
C
call motiondetect(frm1msk,frm2msk,ibs,indx)
            if (indx .eq. 1) then
       mcount=mcount+1
            do 410 il=1, ibs+ip*2
       do 410 j1=1,ibs+ip*2
         searg(i1,j1)=pics(i1+iact-1+ip-ip,j1+jact-1+ip-ip)
410
           continue
        do 420 i2=1,ibs
         do 420 j2=1,ibs
```

```
mask(i2,j2) = frame2(iact-1+i2,jact-1+j2)
420
           continue
if motion is detected, it is estimated and predicted
С
call matct(ip, ibs, imthd, n, nn, Num)
       motionv(i,j) = max(abs(n-7),abs(nn-7))
       mcvector(i,j)=Num
            do 430 i3=1.ibs
                do 430 j3=1,ibs
            recon2(iact-1+i3, jact-1+j3)=pics(iact+ip-1+(n-ip)-1+i
                                       jact+ip-1+(nn-ip)-1+j3)
    +
430
            continue
             else
            motionv(i,j)=0
            mcvector(i,j)=0
          do 402 k1=1,ibs
           do 402 l1=1,ibs
         recon2(iact-1+k1, jact-1+l1)=frame1(iact-1+k1, jact-1+l1)
          continue
402
           endif
    else
               do 210 i1=1, ibs+ip*2
         do 210 j1=1,ibs+ip*2
           searg(i1,j1)=pics(i1+iact-1+ip-ip,j1+jact-1+ip-ip)
210
             continue
          do 220 i2=1,ibs
           do 220 j2=1,ibs
          mask(i2,j2) = frame2(iact-1+i2,jact-1+j2)
220
          continue
    call matct(ip, ibs, imthd, n, nn, Num)
           do 230 i3=1,ibs
                do 230 j3=1,ibs
            recon2(iact-1+i3, jact-1+j3)=pics(iact+ip-1+(n-ip)-1+i
                                       jact+ip-1+(nn-ip)-1+j3)
    +
230
           continue
    endif
200 continue
```

```
MCFD signal is generated using
С
err=0.0
      do 300 i=1,nx
       do 300 j=1,ny
       err=(err+abs(frame2(i,j)-recon2(i,j)))
С
       irecon2(i,j)=recon2(i,j)
     err1(i,j) = frame2(i,j) - recon2(i,j)
 300
      continue
      if(ifld.eq.3) then
С
        call write 1frm(err1,'diff3')
С
      endif
С
С
        test
С
С
      if(ifld.eq.3) then
        call writeimgs(err1,400,512,'diff3.img')
С
      endif
С
С
      if(ifld.eq.10) then
С
С
        call writeimgs(err1,400,512,'diff10.img')
С
      endif
С
    print *,'the value of err=',err
      write(35,*) 'the value of err=',err
    if (imdetect .eq. 1) then
     print *, 'Number of blocks motion detected = ', mcount
     write(35,*)'Number of blocks motion detected =', mcount
    endif
    do 350 i = 1, nx
     do 350 j = 1, ny
          errtemp(i,j) = err1(i,j)
350
      continue
          1
С
    call bitrates(mcvector,entropy,50,64)
the coding of the MCFD signal is carried out here
С
```

```
С
```

```
С
     if(ifld.eq.10) then
     call writeint1(motionv, 50, 64, 'motion. 10')
С
С
     call writeint(err1,400,512,'diff.10')
С
     endif
С
     if(ifld.eq.24) then
С
     call writeint1(motionv,50,64,'motion.24')
С
     call writeint(err1,400,512,'diff.24')
     endif
С
С
С
     if(ifld.eq.30) then
С
     call writeint1(motionv, 50, 64, 'motion. 30')
     call writeint(err1,400,512,'diff.30')
С
     endif
С
С
```

С

```
call qmf(err1)
          print *, 'Error Signal after the vector quantization'
          write(35,*)'Error Signal after the vector quantization'
       vecmean = 0.0
    vecvar = 0.0
       do 351 i = 1, nx
       do 351 j = 1, ny
         vecmean = vecmean + (errtemp(i,j) - errl(i,j))
       vecvar = vecvar + (errtemp(i,j) - errl(i,j))**2
351
       continue
    vecmean = vecmean/(nx*ny)
    vecvar = vecvar/(nx*ny)
    vecvar = vecvar - vecmean**2
    print *, 'Variance of error signal before quant minus after
quant=',
    +
      vecvar
      write(35,*) 'Variance of error signal before quant minus af
quant
    +
       = ',vecvar
Quantized MCFD signal is added to the motion compensated
С
С
   prediction of the current frame and put into frame1 and
   this becomes the previous frame for the next current frame
C
errqnt = 0.0
    errqnt1 = 0.0
      errmean = 0.0
        do 1000 i=1,nx
```

```
do 1000 j=1,ny
    frame1(i,j)=recon2(i,j)+err1(i,j)
errqnt1 = errqnt1 + (frame2(i,j)-frame1(i,j))**2
```

```
1000
       continue
    call write out frm(ifld, frame1)
       xmers=errqnt1/(nx*ny)
       write(35,*) 'MEAN SQUARE ERROR ', xmers
       write(*,*) 'MEAN SQUARE EROR', xmers
       snr=10.0*alog10(255**2/xmers)
       write(35,*) 'S N R = ', snr
      write(36,*) 'S N R = ', snr
       fbitrate = entropy/(64.0) + entropy2/(400*512)
      write(*,*) entropy,entropy2,fbitrate,'final'
      write(38,*) ifld,fbitrate
      write(45,*) ifld,fbitrate,snr,gtotal
      write(*,*) ifld,fbitrate,snr,gtotal
C
   if all the frames are not processed go back
if (ifld .lt. mfld) goto 6000
    stop
    end
subroutine matct
C
subroutine matct(ip, ibs, imthd, n, nn, Num)
    common /AAA/ searg(24,24),mask(8,8)
    real test(13, 13)
    do 50 i=1,2*ip+1
     do 50 j=1,2*ip+1
       test(i,j)=0.0
       do 50 ii=1,ibs
         do 50 jj=1,ibs
           test(i,j)=abs(mask(ii,jj)-searg(i+ii-1,j+jj-1))+test(i
50
        continue
С
   Brute force technique
      if (imthd .eq. 1 ) then
        tmin=1.0e20
         do 100 i=1,2*ip+1
          do 100 k=1,2*ip+1
           if(test(i,k) .lt. tmin) then
```

tmin=test(i,k)

```
n=i
            nn=k
        endif
 100
         continue
    else
    call ortho(test, ip, ibs, icent, jcent)
    n=icent
    nn=icent
    endif
С
  Generates anumber between 1 & 169, The number idicates
С
  the motion information
С
C
       Num = (n-1) * (ip*2+1) + nn
       write(*,*) Num,n,nn
С
       return
    end
    subroutine ortho(test, ip, ibs, icent, jcent)
Independent Orthognal Search Technique
C
real test(ip*2+1,ip*2+1)
    icent=ip+1
    jcent=ip+1
    l=ip/2.+.5
    istep=0
             ł
10
    if ((test(icent, jcent) .lt. test(icent, jcent-1)) .and.
          (test(icent, jcent) .lt. test(icent, jcent+1))) then
    +
            icent=icent
         jcent=jcent
    else if ((test(icent, jcent-1) .lt. test(icent, jcent)) .and.
    +
          (test(icent,jcent-l) .lt. test(icent,jcent+l)))
                                                       then
            icent=icent
         jcent=jcent-1
    else if ((test(icent, jcent+1) .lt. test(icent, jcent)) .and.
    +
          (test(icent,jcent+1) .lt. test(icent,jcent-1)))
                                                       then
            icent=icent
         jcent=jcent+1
```
```
istep=istep+1
    if ((test(icent, jcent) .lt. test(icent-1, jcent)) .and.
          (test(icent,jcent) .lt. test(icent+l,jcent))) then
    +
            icent=icent
         icent=icent
    else if ((test(icent-l, jcent) .lt. test(icent, jcent)) .and.
          (test(icent-l,jcent) .lt. test(icent+l,jcent)))
                                                     then
            icent=icent-1
         jcent=jcent
    else if ((test(icent+1, jcent) .lt. test(icent, jcent)) .and.
          (test(icent+1, jcent) .lt. test(icent-1, jcent))) then
    +
            icent=icent+1
         jcent=jcent
    endif
    istep=istep+1
    if (l .ne. 1) then
       l=(1/2.0+.5)
       go to 10
       endif
    return
    end
      subroutine motiondetect(frm1msk,frm2msk,ibs,indx)
Subroutine calculates if motion is present in the
   (ibs*ibs) subblock
real frmlmsk(ibs, ibs), frm2msk(ibs, ibs)
    kount=0
    do 10 i=1, ibs
      do 10 j=1,ibs
        thrsh=abs(frmlmsk(i,j)-frm2msk(i,j))
         if (thrsh .gt. 3.0) kount=kount+1
 10 continue
    if (kount .qt. 10) then
         indx=1
    else
         indx=0
    endif
    print *,'index',indx
    return
    end
```

endif

С

C

```
59
```

C THE NEW QMF

```
С__
        subroutine qmf(image)
        integer raw, col
        parameter(raw=400,col=512)
        real image(raw,col)
        real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20)
        common /a/ coff1,coff2,coff3,coff4,ltap1,mtap1,ltap2,mtap2
     &
              ,ltap3,mtap3,ltap4,mtap4
c these are the four subband
        real llband(raw/2,col/2),lhband(raw/2,col/2)
     &
       ,hlband(raw/2,col/2),hhband(raw/2,col/2)
        common /bnds/ llband,lhband
        ,hlband,hhband
     &
С
  these are the high and low bands
        real lband(raw,col/2),hband(raw,col/2)
        common /band4/ lband,hband
c input and output images
        real inimg(raw,col),outimg(raw,col)
        common /i/ inimg
        common /o/ outimg
        real lli(raw,col/2),lhi(raw,col/2),hli(raw,col/2),
    & hhi(raw,col/2),llo(raw,col/2),lho(raw,col/2),hlo(raw,col/2)
    & hho(raw,col/2)
       common/band2/ lli,lhi,hli,
    & hhi,llo,lho,hlo,
    &
       hho
```

```
real li(raw,col/2),lo(raw,col),hi(raw,col/2),
```

```
&
        ho(raw,col)
        common /bnd1/ li,lo,hi,
     &
        ho
               limg(raw,col),himg(raw,col)
        real
        common /bnd0/ limg, himg
c motion vector
     integer motionv(50,64)
        common /motionv/ motionv
        common /ifld/ ifld
        common /hist/ hist
        integer hist(9,512)
        common /bitrat/ entropy2,entropy
С
        do 12 i=1, raw
          do 12 j=1,col
          inimg(i,j)=image(i,j)
  12
        continue
        call analysis
С
С
        if(ifld.eq.2) then
          call writeint(llband,raw/2,col/2,'band1.2')
С
          call writeint(lhband,raw/2,col/2,'band2.2')
С
С
          call writeint(hlband,raw/2,col/2,'band3.2')
С
          call writeint1(motionv,50,64,'motionv.2')
С
        elseif(ifld.eq.7) then
          call writeint(llband,raw/2,col/2,'band1.7')
С
С
          call writeint(lhband,raw/2,col/2,'band2.7')
С
          call writeint(hlband,raw/2,col/2,'band3.7')
          call writeint1(motionv, 50, 64, 'motionv.7')
С
        elseif(ifld.eq.10) then
С
          call writeint(llband,raw/2,col/2,'band1.10')
С
С
          call writeint(lhband,raw/2,col/2,'band2.10')
          call writeint(hlband,raw/2,col/2,'band3.10')
С
С
          call writeint1(motionv, 50, 64, 'motionv.10')
        elseif(ifld.eq.15) then
С
С
          call writeint(llband,raw/2,col/2,'band1.15')
          call writeint(lhband,raw/2,col/2,'band2.15')
С
          call writeint(hlband,raw/2,col/2,'band3.15')
С
          call writeint1(motionv, 50, 64, 'motionv. 15')
С
С
        elseif(ifld.eq.20) then
          call writeint(llband,raw/2,col/2,'band1.20')
С
С
          call writeint(lhband,raw/2,col/2,'band2.20')
          call writeint(hlband,raw/2,col/2,'band3.20')
С
С
          call writeint1(motionv,50,64,'motionv.20')
С
        elseif(ifld.eq.22) then
          call writeint(llband,raw/2,col/2,'band1.22')
С
```

С	<pre>call writeint(lhband,raw/2,col/2,'band2.22')</pre>
С	call writeint(hlband,raw/2,col/2,'band3.22')
С	call writeint1(motionv, 50, 64, 'motionv.22')
С	elseif(ifld.eq.24) then
С	call writeint(llband,raw/2,col/2,'band1.24')
С	call writeint(lhband,raw/2,col/2,'band2.24')
С	call writeint(hlband,raw/2,col/2,'band3.24')
С	call writeint1(motionv,50,64,'motionv.24')
С	elseif(ifld.eq.25) then
С	<pre>call writeint(llband,raw/2,col/2,'band1.25')</pre>
С	call writeint(lhband,raw/2,col/2,'band2.25')
С	call writeint(hlband,raw/2,col/2,'band3.25')
С	call writeint1(motionv,50,64,'motionv.25')
С	elseif(ifld.eq.27) then
С	call writeint(llband,raw/2,col/2,'band1.27')
С	<pre>call writeint(lhband,raw/2,col/2,'band2.27')</pre>
С	<pre>call writeint(hlband,raw/2,col/2,'band3.27')</pre>
С	call writeint1(motionv,50,64,'motionv.27')
С	elseif(ifld.eq.30) then
С	call writeint(llband,raw/2,col/2,'band1.30')
С	call writeint(lhband,raw/2,col/2,'band2.30')
С	call writeint(hlband,raw/2,col/2,'band3.30')
С	call writeint1(motionv,50,64,'motionv.30')
С	elseif(ifld.eq.35) then
С	call writeint(llband,raw/2,col/2,'band1.35')
С	call writeint(lhband,raw/2,col/2,'band2.35')
С	call writeint(hlband,raw/2,col/2,'band3.35')
С	call writeint1(motionv, 50, 64, 'motionv. 35')
С	elseif(ifld.eq.38) then
С	call writeint(llband,raw/2,col/2,'band1.38')
C	call writeint(lhband,raw/2,col/2,'band2.38')
С	call writeint(hlband,raw/2,col/2,'band3.38')
С	call writeint1(motionv,50,64,'motionv.38')
С	endif
С	
	Call segma(llband,raw/2,col/2,x1)
	call segma(lhband,raw/2,col/2,x2)
	call segma(hlband,raw/2,col/2,x3)
	call segma(h) and raw/2 col/2 x4)
	<pre>write(35,*) 'var=',x1,x2,x3,x4,(x1+x2+x3+x4)/4.0</pre>
с	

do 191 i=1,9
do 191 j=1,512
hist(i,j)=0

```
call vec_quan(llband,0)
        call vec quan(lhband,1)
        call vec_quan(hlband,2)
С
        call vec quan(hhband)
        call vbitrates(hist,entropy2,9,512)
        do 20 i=1,raw/2
        do 20 j=1,col/2
            hhband(i,j)=0.0
20
     continue
        call recon
        call m error(inimg,outimg,col,raw,output)
        write(*,*) 'error',output
write(35,*) ' mean error square of the error signal'
        write(35,*) 'between after subband and vector quantization'
        write(35,*) 'mse=',output
        call segma(llband,raw/2,col/2,x1)
        call segma(lhband,raw/2,col/2,x2)
        call segma(hlband, raw/2, col/2, x3)
        call segma(hhband, raw/2, col/2, x4)
        write(35,*) 'var=',x1,x2,x3,x4,(x1+x2+x3+x4)/4.0
        call segma(inimg,raw,col,x5)
        write(35,*) 'inimg var',x5
        call segma(outimg, raw, col, x6)
        write(35,*) 'outimg var=',x6
        do 10 i=1, raw
         do 10 j=1,col
         image(i,j)=outimg(i,j)
 10
        continue
C-----
           ------
        return
```

end

С		
c C C	SU TH	BROUTINE READ IS SUBROUTINE READ THE DATA OF FILTERS COEFFICEINTS
с		subroutine readf SUBROUTINE READF
		real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20
	æ	<pre>common /a/ coff1,coff2,coff3,coff4,ltap1,mtap1,ltap2,mtap2 ,ltap3,mtap3,ltap4,mtap4</pre>
c	10	<pre>call openf read(11,*)ltap1 write(*,*)ltap1 read(11,*)mtap1 do 10 i=ltap1,mtap1 read(11,*) coff1(i)</pre>
с	20	<pre>read(11,*)ltap2 read(11,*)mtap2 do 20 i=ltap2,mtap2</pre>
с	30	<pre>read(11,*)ltap3 read(11,*)mtap3 do 30 i=ltap3,mtap3</pre>
с	40	<pre>read(11,*)ltap4 read(11,*)mtap4 do 40 i=ltap4,mtap4</pre>
	5	format(a80) close (11)
с		RETURN END
c c		subroutine initial This subprogram initialize the main program
С		subroutine initial character*80 input_file

,

```
common /ina/ input file
        write (*,1)
     write (*,3)
        read (5,4) input file
 1
        format ('
                     1)
 3
        format (' Enter the name of the file contains the order of
        ',/,' filtes there coefficients,
     &
        input Image, and output file:')
     &
 4
        format( a80)
        return
        end
С
        subroutine openf
        character*80 input file
        common/ina/ input file
        open (11, file='.../filters.dir/in24', status='old')
        return
        end
c____
        subroutine analysis
        integer raw, col
        parameter(raw=400, col=512)
        real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20)
        common /a/ coff1,coff2,coff3,coff4,ltap1,mtap1,ltap2,mtap2
     &
               ,ltap3,mtap3,ltap4,mtap4
c these are the four subband
        real llband(raw/2,col/2),lhband(raw/2,col/2)
        , hlband(raw/2, col/2), hhband(raw/2, col/2)
     &
        common /bnds/ llband,lhband
        ,hlband,hhband
     &
c these are the high and low bands
        real lband(raw,col/2),hband(raw,col/2)
```

```
65
```

```
common /band4/ lband, hband
c input and output images
        real inimg(raw,col),outimg(raw,col)
        common /i/ inimg
        common /o/ outimg
        nx=raw
        ny=col
        call rfilter(coff1, inimg, lband, nx, ny, ltap1, mtap1)
        call rfilter(coff2, inimg, hband, nx, ny, ltap2, mtap2)
        ny=ny/2
        call cfilter(coff1,lband,llband,nx,ny,ltap1,mtap1)
        call cfilter(coff2,lband,lhband,nx,ny,ltap2,mtap2)
        call cfilter(coff1, hband, hlband, nx, ny, ltap1, mtap1)
        call cfilter(coff2, hband, hhband, nx, ny, ltap2, mtap2)
        return
        end
```

subroutine cfilter(f,a1,a2,raw,col,ltap,mtap)

```
integer col,raw,ltap,mtap,jk
        real a1(raw,col),a2(raw/2,col),f(-20:20)
        do 20 i=1,col
          do 20 j=2,raw,2
              a2(j/2,i)=0
            do 20 k=ltap,mtap
              jk=j+k
              if(jk.le.0) jk=raw+jk
              if(jk.gt.raw) jk=jk-raw
              a2(j/2,i)=a2(j/2,i)+a1(jk,i)*f(k)
  20
        continue
        return
        end
С
       subroutine ccfilter(f,a1,a2,raw,col,ltap,mtap)
        integer col, raw, ltap, mtap, jk
        real a1(raw, col), a2(raw, col), f(-20:20)
        do 20 i=1,col
          do 20 j=1,raw
            a2(j,i)=0
            do 20 k=ltap,mtap
             jk=j+k
             if(jk.le.0) jk=raw+jk
             if(jk.gt.raw) jk=jk-raw
             a2(j,i)=a2(j,i)+a1(jk,i)*f(k)
  20
        continue
        return
        end
С
       subroutine rcfilter(f,a1,a2,raw,col,ltap,mtap)
        integer col, raw, ltap, mtap, jk
        real a1(raw,col),a2(raw,col),f(-20:20)
        do 20 i=1,raw
          do 20 j=1,col
             a2(i,j)=0
            do'20 k=ltap,mtap
             jk=j+k
             if(jk.le.0) jk=col+jk
             if(jk.gt.col) jk=jk-col
             a2(i,j)=a2(i,j)+a1(i,jk)*f(k)
  20
        continue
        return
        end
С
```

```
67
```

1

```
subroutine cinter(in,out,nraw,ncol)
         integer nraw, ncol
         real in(nraw,ncol),out(nraw*2,ncol)
         do 20 j=1,ncol
           do 20 i=1,nraw
               out(2*i-1,j)=in(i,j)
               out(2*i,j)=0.0
               out(2*i,j)=in(i,j)
С
               out(2*i-1,j)=0.0
С
 20
         continue
         return
        end
С
         subroutine rinter(in,out,nraw,ncol)
         integer nraw, ncol
        real in(nraw,ncol),out(nraw,2*ncol)
        do 20 j=1,ncol
           do 20 i=1, nraw
               out(i, 2*j-1) = in(i, j)
               out(i, 2*j) = 0.0
С
               out(i,2*j)=in(i,j)
С
               out(i, 2*j-1) = 0.0
 20
        continue
        return
        end
С
        subroutine recon
        integer raw, col
        parameter(raw=400, col=512)
c input and output images
        real inimg(raw,col),outimg(raw,col)
        common /i/ inimg
        common /o/ outimg
c these are the four subband
        real llband(raw/2,col/2),lhband(raw/2,col/2)
        , hlband(raw/2, col/2), hhband(raw/2, col/2)
     &
        common /bnds/ llband, lhband
       , hlband, hhband
     &
```

```
real lli(raw,col/2),lhi(raw,col/2),hli(raw,col/2),
& hhi(raw,col/2),llo(raw,col/2),lho(raw,col/2),hlo(raw,col/2)
```

.

```
&
        hho(raw, col/2)
        common/band2/ lli,lhi,hli,
        hhi,llo,lho,hlo,
     &
     &
        hho
        real li(raw,col/2),lo(raw,col),hi(raw,col/2),
     &
        ho(raw,col)
        common /bnd1/ li,lo,hi,
     &
        ho
              limg(raw,col),himg(raw,col)
        real
        common /bnd0/ limg,himg
        real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20)
        common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
     &
              ,ltap3,mtap3,ltap4,mtap4
С
        nraw=raw
        ncol=col
        call cinter(llband,lli,raw/2,col/2)
        call ccfilter(coff4,lli,llo,raw,col/2,ltap4,mtap4)
        call cinter(lhband, lhi, raw/2, col/2)
        call ccfilter(coff3,lhi,lho,raw,col/2,ltap3,mtap3)
        call cinter(hlband,hli,raw/2,col/2)
        call ccfilter(coff4, hli, hlo, raw, col/2, ltap4, mtap4)
        call cinter(hhband, hhi, raw/2, col/2)
        call ccfilter(coff3, hhi, hho, raw, col/2, ltap3, mtap3)
        do 10 i=1,raw
          do 10 j=1,col/2
            li(i,j)=llo(i,j)+lho(i,j)
            hi(i,j)=hlo(i,j)+hho(i,j)
 10
        continue
        call rinter(li,lo,raw,col/2)
        call rcfilter(coff4,lo,limg,raw,col,ltap4,mtap4)
        call rinter(hi,ho,raw,col/2)
        call rcfilter(coff3,ho,himg,raw,col,ltap3,mtap3)
       do 20 i=1,raw
```

```
do 20 j=1,col
            outimg(i,j)=(limg(i,j)+himg(i,j))
            if(outimg(i,j).gt.255)outimg(i,j)=255
С
С
            if(outimg(i,j).lt.0)outimg(i,j)=0
   20
         continue
        return
        end
С
       subroutine segma(image,raw,col,output)
       integer raw, col
       real image(raw, col), output
С
       double precision s1,s2
       s1=0.0
       s2=0.0
       do 10 i=1, raw
          do 10 j=1,col
           sl=sl+image(i,j)
           s2=s2+(image(i,j))**2
 10
      continue
       output=s2/(raw*col)-(s1/(raw*col))**2
       return
       end
```

10

```
subroutine m_error(image1,image2,raw,col,output)
integer raw,col
real image1(raw,col),image2(raw,col),output
s=0.0
do 10 i=1,raw
    do 10 j=1,col
        s=s+(image1(i,j)-image2(i,j))**2
continue
output=s/(raw*col)
return
```

end

```
С
```

```
subroutine writeimqs(pic,nx,ny,name)
      real pic(nx,ny)
      character*1 pim(400*512)
      character*20 name
        open(1,file=name,access='direct',
    + form='unformatted', recl=nx*ny)
      do 20 i=1.nx
        do 20 j=1,ny
          ip=int(pic(i,j))+128
          if(ip.qt.255) ip=255
          if(ip.lt.0) ip=0
          if(ip.gt.128) ip=ip-256
           mm=j+(i-1)*ny
           pim(mm) = char(ip)
20
      continue
      write(1,rec=1) (pim(j),j=1,nx*ny)
      close (1)
      return
      end
```

с с

```
subroutine writeimg(pic,nx,ny,name)
      real pic(nx,ny)
      character*1 pim(400*512)
      character*20 name
        open(1,file=name,access='direct',
    + form='unformatted',recl=nx*ny)
      do 20 i=1,nx
        do 20 j=1,ny
          ip=int(pic(i,j))
          if(ip.gt.128) ip=ip-256
           mm=j+(i-1)*ny
           pim(mm) = char(ip)
20
      continue
      write(l,rec=1) (pim(j),j=1,nx*ny)
      close (1)
      return
      end
```

С С

```
subroutine writeint(pic,nx,ny,name)
       real pic(nx,ny)
       integer ipic(400,512)
       character*20 name
       open(1,file=name)
       do 5 i=1,nx
       do 5 j=1,ny
         ipic(i,j)=int(pic(i,j)+0.5)
5
       continue
       do 10 i=1,nx
        write(1,*) (ipic(i,j),j=1,ny)
10
       continue
       close (1)
       return
       end
```

С

```
subroutine writeint1(pic,nx,ny,name)
       integer pic(nx,ny)
       character*20 name
       open(1,file=name)
       do 10 i=1,nx
        write(1,*) (pic(i,j),j=1,ny)
10
       continue
       close (1)
       return
       end
```

```
subroutine vec_quan(pic,nband)
c pic: picture o be coded (200X256)
c nband: the band number (0,1,2,3)
        real pic(200,256)
        integer motionv(50,64)
```

```
real tvector(16)
       common /hist/ hist
       integer hist(9,512)
         common /motionv/ motionv
       real blv1(512,16),b1v2(512,16),b1v3(512,16)
       real b2v1(512,16),b2v2(512,16),b2v3(512,16)
       real b3v1(512,16),b3v2(512,16),b3v3(512,16)
       common /vqcodebook/ b1v1,b1v2,b1v3,b2v1,b2v2,b2v3,b3v1,b3v2
    &
                            b3v3
       nn=nband*3
       do 10 i=1,50
       do 10 j=1,64
         if(motionv(i,j).ge.5) then
           do 20 k=1,4
           do 20 l=1,4
             tvector(4*(k-1)+1) = pic((i-1)*4+k, (j-1)*4+1)
20
           continue
           mm=nn+3
         else if(motionv(i,j).ge.3) then
           do 21 k=1,4
           do 21 l=1,4
             tvector(4*(k-1)+1)=pic((i-1)*4+k,(j-1)*4+1)
21
           continue
           mm=nn+2
         else if (motionv(i,j).ge.1) then
           do 22 k=1,4
           do 22 l=1,4
             tvector(4*(k-1)+1)=pic((i-1)*4+k,(j-1)*4+1)
22
           continue
           mm=nn+1
         else
           do 33 k=1,4
           do 33 l=1,4
              pic((i-1)*4+k, (j-1)*4+1)=0.0
33
        continue
           mm=0
         endif
      if(mm.eq.1) then
           call vquantizer(tvector, b1v1, ivecnum)
           hist(1,ivecnum) = hist(1,ivecnum)+1
         else if (mm.eq.2) then
           call vquantizer(tvector, b1v2, ivecnum)
           hist(2,ivecnum) = hist(2,ivecnum)+1
         else: if (mm.eq.3) then
```

call vguantizer(tvector, b1v3, ivecnum) hist(3, ivecnum) = hist(3, ivecnum) + 1else if (mm.eq.4) then call vquantizer(tvector, b2v3, ivecnum) hist(4,ivecnum) = hist(4,ivecnum)+1 else if (mm.eq.5) then call vquantizer(tvector, b2v2, ivecnum) hist(5, ivecnum) = hist(5, ivecnum)+1 else if (mm.eq.6) then call vguantizer(tvector, b2v3, ivecnum) hist(6, ivecnum) = hist(6, ivecnum)+1 else if(mm.eq.7)then call vquantizer(tvector,b3v1,ivecnum) hist(7, ivecnum) = hist(7, ivecnum) + 1else if(mm.eq.8)then call vquantizer(tvector,b3v2,ivecnum) hist(8, ivecnum) = hist(8, ivecnum)+1 else, if (mm.eq.9) then call vquantizer(tvector,b3v3,ivecnum) hist(9,ivecnum) = hist(9,ivecnum)+1 endif if (mm.ne.0) then do 44 k=1,4 do 44 l=1,4 pic((i-1)\*4+k,(j-1)\*4+1) = tvector(4\*(k-1)+1)44 continue endif 10 continue return end С subroutine vquantizer(testv,codebook,ivecnu) С Best Matching of vector real testv(16) real codebook(512,16) rdiff = 1000000.0ivecnu = 0

1

```
do 110 m = 1,512
           adiff = 0
           do 120 n = 1,16
            adiff = adiff + (testv(n) - codebook(m,n))**2
120
           continue
           if (adiff .lt. rdiff) then
              rdiff = adiff
              ivecnu = m
           endif
110
       continue
       do 130 n = 1,16
           testv(n) = codebook(ivecnu,n)
130
       continue
       return
       end
```

```
С
```

```
subroutine readvq
```

```
real blv1(512,16),blv2(512,16),blv3(512,16)
        real b2v1(512,16),b2v2(512,16),b2v3(512,16)
        real b3v1(512,16),b3v2(512,16),b3v3(512,16)
        common /vqcodebook/ blv1,blv2,blv3,b2v1,b2v2,b2v3,b3v1,b3v2
     &
                             b3v3
     open (15,file='../quantizer.dir/b1v1.12')
        do 10 i=1,512
          read(15,*) (blv1(i,j),j=1,16)
10
     continue
        close(15)
     open (16, file='.../quantizer.dir/b1v2.34')
        do 11 i=1,512
          read(16,*) (b1v2(i,j),j=1,16)
11
     continue
        close(16)
     open (17, file='.../quantizer.dir/b1v3.56')
        do 12 i=1,512
          read(17,*) (blv3(i,j),j=1,16)
12
     continue
        close(17)
     open (18,file='../quantizer.dir/b2v1.12')
        do 13 i=1,512
          read(18,*) (b2v1(i,j),j=1,16)
13
     continue
```

```
close(18)
     open (19, file='.../quantizer.dir/b2v2.34')
         do 14 i=1,512
           read(19,*) (b2v2(i,j),j=1,16)
14
     continue
        close(19)
     open (20,file='../quantizer.dir/b2v3.56')
        do 15 i=1,512
          read(20,*) (b2v3(i,j),j=1,16)
15
     continue
        close(20)
     open (21,file='../quantizer.dir/b3v1.12')
        do 16 i=1,512
          read(21,*) (b3v1(i,j),j=1,16)
16
     continue
        close(21)
     open (22, file='.../quantizer.dir/b3v2.34')
        do 17 i=1,512
          read(22,*) (b3v2(i,j),j=1,16)
17
     continue
        close(22)
     open (23,file='../quantizer.dir/b3v3.56')
        do 18 i=1,512
          read(23,*) (b3v3(i,j),j=1,16)
18
     continue
        close(23)
        return
     end
С
       subroutine read frm(ifld,pic)
       character*1 pic(400,512)
       open(1,file='ali.100',access='direct',form=
С
     & 'unformatted', recl=512*512)
С
С
С
       open(1,file='/images/cindy40',access='direct',form=
С
     & 'unformatted', recl=512)
       open(1,file='/images/mono',access='direct',form=
     & 'unformatted', recl=512)
       open(1,file='/images/guartet',access='direct',form=
С
```

```
'unformatted', recl=512)
С
     &
С
С
       open(1,file='/images/duo',access='direct',form=
С
     & 'unformatted', recl=512)
С
        icod1 = (ifld-1)*400
        icod2 = (ifld-1) * 400 + 200
       do 10 i=1,200
        read(1,rec=icod1+i)(pic(2*i-1,j),j=1,512)
        read(1,rec=icod2+i)(pic(i*2,j),j=1,512)
10
       continue
        close(1)
        return
        end
С
     subroutine write in frm(ifld,pic)
        real pic(400,512)
        character*1 image(512*512)
        open(22,file='mono.in80',access='direct',form=
       'unformatted', recl=512*512)
     &
     do 10 i=1,400
        do 10 j=1,512
         ip=int(pic(i,j)+.5)
         if(ip.gt.255) ip=255
         if(ip.lt.0) ip=0
         if(ip.gt.127) ip=ip-256
         image((i-1)*512+j) = char(ip)
 10 continue
        do 20 i=204801,262144
          image(i) = char(003)
 20
        continue
        write(22, rec=ifld)(image(j), j=1, 512*512)
     close(22)
     return
     end
С
```

```
subroutine write_out_frm(ifld,pic)
real pic(400,512)
character*1 image(512*512)
open(21,file='mono.out80',access='direct',form=
```

```
&
       'unformatted', recl=512*512)
    do 10 i=1,400
       do 10 j=1,512
        ip=int(pic(i,j)+.5)
        if(ip.gt.255) ip=255
        if(ip.lt.0) ip=0
        if(ip.gt.127) ip=ip-256
        image((i-1)*512+j) = char(ip)
10
   continue
       do 20 i=204801,262144
         image(i) = char(100)
20
       continue
       write(21, rec=ifld)(image(j), j=1, 512*512)
    close(21)
    return
    end
```

```
С
```

```
subroutine write lfrm(pic,name)
        real pic(400, \overline{5}12)
        character*1 image(512*400)
        character*20 name
        open(22,file=name,access='direct',form=
     & 'unformatted', recl=512*400)
     do 10 i=1,400
        do 10 j=1,512
         ip=int(pic(i,j)+.5)
         if(ip.gt.255) ip=255
         if(ip.lt.0) ip=0
         if(ip.gt.127) ip=ip-256
         image((i-1)*512+j) = char(ip)
 10 continue
        write(22,rec=1)(image(j),j=1,512*400)
     close(22)
     return
     end
              1
С
```

c this subroutine calculate the entropy of each band c and find the probability of each code

.

```
subroutine vbitrates(ic,bitrate,raw,col)
         common /gtotal/ gtotal
         integer ic(raw,col),raw,col
         real entropy(9), sum(512), pr(512)
         gtotal=0
        bitrate=0
     do 10 m=1,9
         total=0
         do 20 n=1,512
          sum(n) = ic(m, n)
          total=total+sum(n)
20
     continue
С
         entropy(m) = 0.0
         do 30 n=1,512
            pr(n) = sum(n) / total
            if(pr(n).gt.0) then
              br=pr(n) *xlog2(1.0/pr(n))
              entropy(m) = entropy(m) + br
            endif
 30
        continue
        bitrate=bitrate+entropy(m)*total
        write(*,*) 'total = ',total
        gtotal=gtotal+total
 10 continue
        gtotal=gtotal/3
        write(*,*) 'gtotal = ',gtotal
write(*,*) 'ventropy = ',(entropy(i),i=1,9)
        write(*,*) 'bitrate = ', bitrate
        return
        end
```

```
С
```

c this subroutine calculate the entropy of each band c and find the probability of each code

subroutine bitrates(ic,entropy,raw,col)

integer ic(raw,col),raw,col
real entropy,sum(0:169),pr(0:169)

do 20 n=0,169

```
sum(n)=0.0
 20
        continue
С
        do 10 i=1, raw
        do 10 j=1,col
            k=ic(i,j)
            sum(k) = sum(k) + 1
 10
        continue
        entropy=0.0
        total=real(raw*col)
        do 30 n=0,169
           pr(n) = sum(n) / total
            if(pr(n).gt.0) then
              br=pr(n)*xlog2(1.0/pr(n))
              entropy=entropy+br
           endif
 30
        continue
        write(*,*) 'entropy = ',entropy
        write(*,*) 'pr = ', (pr(i), i=1, 169)
С
        return
        end
С
        function xlog2(x)
        real x
        x\log_2=\log(x)/\log(2.0)
        return
        end
С
c
```

С Simulation program for 10 band decomposition С Name: Hosam Mutlag С MAIN С character\*80 input file common /ina/ input file character\*80 inputimg common /image/ inputimg common /bitrate/ rav open(50,file='10bdpcm.res',form='formatted') do 20 rav=.2,1.3,0.1 С rav=1.0 input file='in24' inputimg='ladyp.img' call subband10(e1,or1,snr1) input file='in6' call subband10(e2,or2,snr2) input file='in8' call subband10(e3,or3,snr3) write(50,10) rav,snr1,snr2,snr3,e1,or1,e2,or2,e3,or3 20 continue do 30 rav=1.5,5.0,0.5 input file='in24' inputimg='ladyp.img' call subband10(e1,or1,snr1) input file='in6' call subband10(e2,or2,snr2) input file='in8' call subband10(e3,or3,snr3) write(50,10) rav,snr1,snr2,snr3,e1,or1,e2,or2,e3,or3 continue 30 format(f4.2,3f8.4,6f8.5) 10 close(50)

```
stop
        end
С
   This is a 10 band simulation program.
С
   it uses dpcm quantizer for the first band
С
c and laplacian quanizer for the rest. hoffman coading.
c the program does bit allocation for the 10 bands.
С
С
     subroutine subband10(xo1,xo2,snr)
        integer raw, col
        parameter(raw=256,col=256)
        real dimg(32,32), cimg(32,32)
        real p(75),pl(75)
        character*80 input file
        common /ina/ input file
        character*80 inputing
        common /image/ inputimg
        common /bitrate/ rav
        common /header/ header
        character*1 header(64)
        real x(10),output,xmean(10),entrpy(10),entrpy1(4),entrpy2(3
        real entrpy3(3)
        real xt(64), xmeant(64)
        real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20)
        common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
     &
              ,ltap3,mtap3,ltap4,mtap4
        integer iq(64),iqt(64)
        real pr(-64:64), prb(128)
        integer prnum(128)
        character*80 codbook(128)
        character*1 hc(77000)
С
        real inimg(256,256)
        real a1(128,128), a2(128,128), a3(128,128), a4(128,128)
        real b1(64,64),b2(64,64),b3(64,64),b4(64,64)
        real c1(32,32), c2(32,32), c3(32,32), c4(32,32)
```

```
82
```

```
real e1(128,128)
        real d1(64, 64)
        real qc1(32,32), qc2(32,32), qc3(32,32), qc4(32,32)
        real qb2(64,64),qb3(64,64),qb4(64,64)
        real qa2(128,128), qa3(128,128), qa4(128,128)
        integer ic1(4,32,32),icr1(4,32,32)
        integer ic2(3,64,64),icr2(3,64,64)
        integer ic3(3,128,128),icr3(3,128,128)
        real outimg(256, 256)
С
С
        call initial
        call readf
        call readimage(inimg,raw,col)
        write(*,*) 'subband analysis'
        call analysis256(inimg,a1,a2,a3,a4)
        call analysis128(a1,b1,b2,b3,b4)
        call analysis64(b1,c1,c2,c3,c4)
С
         write(*,*) 'calculating the variances and means'
        call segma(c1,32,32,x(1),xmean(1))
        call segma(c2,32,32,x(2),xmean(2))
        call segma(c3,32,32,x(3),xmean(3))
        call segma(c4,32,32,x(4),xmean(4))
        call segma(b_{2,64,64,x(5),xmean(5)})
```

```
call dpcm1(c1,dimg,32,32)
call segma(dimg,32,32,var2,rmean)
write(*,*) x(1),xmean(1)
x(1)=var2
xmean(1)=rmean
write(*,*) var2,rmean
```

call segma(b3,64,64,x(6),xmean(6))
call segma(b4,64,64,x(7),xmean(7))
call segma(a2,128,128,x(8),xmean(8))
call segma(a3,128,128,x(9),xmean(9))
call segma(a4,128,128,x(10),xmean(10))

C C

```
c convert the 10band to 64 band for bit allocation
        do 10 i=1,4
           xt(i)=x(i)
           xmeant(i)=xmean(i)
 10
    continue
     do 15 j=1,3
        do 20 i=(j*4)+1, j*4+4
           xt(i)=x(j+4)
           xmeant(i) = xmean(j+4)
     continue
 20
 15
        continue
     do 11 j=1,3
        do 21 i=(j*16)+1, j*16+16
           xt(i)=x(j+7)
           xmeant(i) = xmean(j+7)
 21
    continue
        continue
 11
С
        write(*,*) (x(i),i=1,10)
        write(*,*) (xt(i),i=1,64)
С
С
        call entropy(igt, xt, xmeant)
С
        iq(1) = iqt(1)
        iq(2)=iqt(2)
        iq(3) = iqt(3)
        iq(4) = iqt(4)
        iq(5) = iqt(5)
        iq(6)=iqt(9)
        iq(7) = iqt(13)
        iq(8) = iqt(17)
        iq(9) = iqt(33)
        iq(10) = iqt(49)
     write(*,*) (iq(i),i=1,10)
        write(*,*) (iqt(i),i=1,64)
С
С
     write(*,*) 'quantizing'
        nm=iq(1)
        call readq(nm,p,pl)
        if(nm.ne.iq(1))then
```

```
write(*,*) 'error in reading quantizer data'
stop
end if
```

	<pre>call qdpcm(c1,ic1,nm,var2,rmean,32,32,p,pl)</pre>
C C	<pre>call unifquan(c1,1,ic1,iq(1),xminx,step,4,32,32) call quantizer(c1,1,ic1,iq(1),xmean(1),x(1),4,32,32) call quantizer(c2,2,ic1,iq(2),xmean(2),x(2),4,32,32) call quantizer(c3,3,ic1,iq(3),xmean(3),x(3),4,32,32) call quantizer(c4,4,ic1,iq(4),xmean(4),x(4),4,32,32)</pre>
	<pre>call quantizer(b2,1,ic2,iq(5),xmean(5),x(5),3,64,64) call quantizer(b3,2,ic2,iq(6),xmean(6),x(6),3,64,64) call quantizer(b4,3,ic2,iq(7),xmean(7),x(7),3,64,64)</pre>
	<pre>call quantizer(a2,1,ic3,iq(8),xmean(8),x(8),3,128,128) call quantizer(a3,2,ic3,iq(9),xmean(9),x(9),3,128,128) call quantizer(a4,3,ic3,iq(10),xmean(10),x(10),3,128,128)</pre>
c	
	sum=0 length=0
	do 25 i=1,4 if(iq(i).gt.0.1) then
_	<pre>call bitrates(ic1,i,entrpy1,pr,4,32,32)</pre>
C	call convert(pr,prb,prnum,ns)
С	call sort(prb,prnum,ns)
	call hcbook(prb,codbook,ns)
-	<pre>call codehc(ic1,i,codbook,prnum,ns,hc,ll,4,32,32)</pre>
C	<pre>call decodehc(icr1,i,codbook,prnum,ns,hc,l1,4,32,32)</pre>
	<pre>do 103 i1=1,32 do 103 j1=1,32     icr=ic1(i,i1,j1)-icr1(i,i1,j1)</pre>
С	<pre>write(*,*) icr if(icr.ne.0) then     write(*,*) 'ERROR2 IN CHANNEL CODING',i,i1,j1,icr     stop endif</pre>
103	continue
	<pre>length=length+ll entrpy(i)=entrpyl(i) sum=sum+entrpyl(i)</pre>

endif

25 continue

sum = sum / 4.0С do 35 i=1,3 if(iq(i+4).gt.0.1) then call bitrates(ic2,i,entrpy2,pr,3,64,64) call convert(pr,prb,prnum,ns) call sort(prb,prnum,ns) call hcbook(prb,codbook,ns) call codehc(ic2, i, codbook, prnum, ns, hc, 11, 3, 64, 64) С call decodehc(icr2, i, codbook, prnum, ns, hc, 11, 3, 64, 64) С do 104 i1=1,64 do 104 j1=1,64 icr=ic2(i,i1,j1)-icr2(i,i1,j1) С write(\*,\*) icr if(icr.ne.0) then write(\*,\*) 'ERROR2 IN CHANNEL CODING',i,i1,j1,icr stop endif 104 continue length=length+l1 entrpy(i+4)=entrpy2(i) sum=sum+entrpy2(i) endif 35 continue sum=sum/4.0С do 45 i=1,3 if(iq(i+7).gt.0.1) then call bitrates(ic3, i, entrpy3, pr, 3, 128, 128) call convert(pr,prb,prnum,ns) call sort(prb,prnum,ns) call hcbook(prb,codbook,ns) call codehc(ic3, i, codbook, prnum, ns, hc, 11, 3, 128, 128) С call decodehc(icr3,i,codbook,prnum,ns,hc,l1,3,128,128) С

86

```
do 105 i1=1,128
         do 105 j1=1,128
           icr=ic3(i,i1,j1)-icr3(i,i1,j1)
С
         write(*,*) icr
         if(icr.ne.0) then
           write(*,*) 'ERROR2 IN CHANNEL CODING',i,i1,j1,icr
           stop
         endif
 105
         continue
         length=length+ll
          entrpy(i+7)=entrpy3(i)
          sum=sum+entrpy3(i)
         endif
 45
    continue
        write(*,*) 'entropy=',sum/4.0
        xol=sum/4.0
        xo2=real(length)/(256.0**2)
С
     write(*,*) 'DEquantizing'
        call dequantizer(cimg,1,icr1,iq(1),rmean,var2,4,32,32)
        call dpcdecoder(cimg,qc1,32,32)
        call m_error(qc1,c1,32,32,cerror)
        write(*,*) 'c1 error',cerror
С
        call unifdequan(qc1,1,icr1,iq(1),xminx,step,4,32,32)
С
        call dequantizer(qc1,1,icr1,iq(1),xmean(1),x(1),4,32,32)
        call dequantizer(qc2,2,icr1,iq(2),xmean(2),x(2),4,32,32)
        call dequantizer(qc3,3,icr1,iq(3),xmean(3),x(3),4,32,32)
        call dequantizer(qc4,4,icr1,iq(4),xmean(4),x(4),4,32,32)
        call dequantizer(qb2,1,icr2,iq(5),xmean(5),x(5),3,64,64)
        call dequantizer(qb3,2,icr2,iq(6),xmean(6),x(6),3,64,64)
        call dequantizer(qb4,3,icr2,iq(7),xmean(7),x(7),3,64,64)
        call dequantizer(qa2,1,icr3,iq(8),xmean(8),x(8),3,128,128)
        call dequantizer(qa3,2,icr3,iq(9),xmean(9),x(9),3,128,128)
        call dequantizer(qa4,3,icr3,iq(10),xmean(10),x(10),3,128,12
С
    write(*,*) 'synthesis'
```

```
call synthesis64(qc1,qc2,qc3,qc4,d1)
call synthesis128(d1,qb2,qb3,qb4,e1)
```

call synthesis256(e1,qa2,qa3,qa4,outimg)

```
С
        do 131 i=1,256
         do 131 j=1,256
           if(outimg(i,j).gt.255) then
          outimg(i,j) = 255
          else if(outimg(i,j).lt.0) then
          outimg(i,j)=0
          endif
 131
        continue
        call m error(inimg,outimg,col,raw,output)
        call writeimg(outimg,256,256,'out.img')
С
С
        write(*,*) 'error',output
        snr=10*alog10(255**2/output)
        write(*,*) 'SNR =',snr
С
        return
        end
С
С
     SUBROUTINE READ
С
     THIS SUBROUTINE READ THE DATA OF FILTERS COEFFICEINTS
С
        subroutine readf
         SUBROUTINE READF
С
        real coff1(-20:20),coff2(-20:20),coff3(-20:20),coff4(-20:20)
        common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
     &
               ,ltap3,mtap3,ltap4,mtap4
        call openf
        write(*,*) 'reading filter coefficeints'
С
        read(11,*)ltap1
         write(*,*)ltap1
С
        read(11,*)mtap1
        do 10 i=ltap1, mtap1
   10
              read(11,*) coff1(i)
С
        read(11,*)ltap2
        read(11,*)mtap2
        do 20 i=ltap2,mtap2
```

```
20
              read(11,*) coff2(i)
С
        read(11,*)ltap3
        read(11,*)mtap3
        do 30 i=ltap3,mtap3
   30
              read(11,*) coff3(i)
С
        read(11,*)ltap4
        read(11,*)mtap4
        do 40 i=ltap4, mtap4
   40
              read(11,*) coff4(i)
С
        close (11)
        RETURN
        END
С
        subroutine initial
С
        This subprogram initialize the main program
С
С
        subroutine initial
        character*80 input file, inputimg
        real rav
        common /ina/ input file
        common /image/ inputimg
        common /bitrate/ rav
        write(*,*) 'initialization'
С
        write (*,1)
     write (*,3)
        read (5,4) input_file
        write(*,*) ' Enter the name of the input image'
        read(*,*) inputimg
        write(*,*) ' Enter the bit rate'
        read(*,*) rav
1
        format ('
                   1)
3
        format (' Enter the name of the file contains the order of
     & ',/,' filtes there coefficients,')
        format( a80)
4
        return
        end
```

```
subroutine openf
```

C\_

```
character*80 input_file
common/ina/ input_file
open (11,file='../filters.dir/'//input_file,status='old')
return
end
```

c\_

```
subroutine analysis256(inimg,llband,lhband,hlband,hhband)
        integer raw, col
        parameter(raw=256,col=256)
        real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20)
        common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
     &
               ,ltap3,mtap3,ltap4,mtap4
c these are the four subband
        real llband(raw/2,col/2),lhband(raw/2,col/2)
     &
        , hlband (raw/2, col/2), hhband (raw/2, col/2)
  these are the high and low bands
С
        real lband(raw,col/2),hband(raw,col/2)
c input and output images
        real inimg(raw,col)
        nx=raw
        ny=col
        call rfilter(coff1, inimg, lband, nx, ny, ltap1, mtap1)
        call rfilter(coff2, inimg, hband, nx, ny, ltap2, mtap2)
        ny=ny/2
        call cfilter(coff1,lband,llband,nx,ny,ltap1,mtap1)
        call cfilter(coff2,lband,lhband,nx,ny,ltap2,mtap2)
```

```
call cfilter(coff1,hband,hlband,nx,ny,ltap1,mtap1)
call cfilter(coff2,hband,hhband,nx,ny,ltap2,mtap2)
return
end
```

c these are the four subband real llband(raw/2,col/2),lhband(raw/2,col/2) & ,hlband(raw/2,col/2),hhband(raw/2,col/2)

c input and output images real inimg(raw,col)

С

nx=raw ny=col call rfilter(coff1, inimg, lband, nx, ny, ltap1, mtap1) call rfilter(coff2, inimg, hband, nx, ny, ltap2, mtap2) ny=ny/2 call cfilter(coff1, lband, llband, nx, ny, ltap1, mtap1) call cfilter(coff2, lband, lhband, nx, ny, ltap2, mtap2)

```
call cfilter(coff1,hband,hlband,nx,ny,ltap1,mtap1)
call cfilter(coff2,hband,hhband,nx,ny,ltap2,mtap2)
return
end
```

```
subroutine analysis64 (inimg, llband, lhband, hlband, hhband)
        integer raw, col
        parameter(raw=64, col=64)
        real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20)
        common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
     &
               ,ltap3,mtap3,ltap4,mtap4
c these are the four subband
        real llband(raw/2,col/2),lhband(raw/2,col/2)
       , hlband (raw/2, col/2), hhband (raw/2, col/2)
     &
c these are the high and low bands
        real lband(raw, col/2), hband(raw, col/2)
c input and output images
        real inimg(raw,col)
        nx=raw
        ny=col
        call rfilter(coff1, inimg, lband, nx, ny, ltap1, mtap1)
        call rfilter(coff2, inimg, hband, nx, ny, ltap2, mtap2)
        ny=ny/2
        call cfilter(coff1,lband,llband,nx,ny,ltap1,mtap1)
        call cfilter(coff2,lband,lhband,nx,ny,ltap2,mtap2)
```

```
call cfilter(coff1, hband, hlband, nx, ny, ltap1, mtap1)
```

```
call cfilter(coff2, hband, hhband, nx, ny, ltap2, mtap2)
        return
        end
С
        subroutine rfilter(f,a1,a2,raw,col,ltap,mtap)
        integer col,raw,ltap,mtap
        real a1(raw,col),a2(raw,col/2),f(-20:20)
        do 20 i=1,raw
          do 20 j=2,col,2
             a2(i,j/2)=0
            do 20 k=ltap,mtap
             ik=i+k
             if(jk.le.0) jk=col+jk
             if(jk.gt.col) jk=jk-col
             a2(i,j/2)=a2(i,j/2)+a1(i,jk)*f(k)
  20
        continue
        return
        end
С
       subroutine cfilter(f,a1,a2,raw,col,ltap,mtap)
        integer col,raw,ltap,mtap,jk
        real a1(raw,col),a2(raw/2,col),f(-20:20)
        do 20 i=1,col
          do 20 j=2, raw, 2
             a2(j/2,i)=0
            do 20 k=ltap,mtap
             jk=j+k
             if(jk.le.0) jk=raw+jk
             if(jk.gt.raw) jk=jk-raw
             a2(j/2,i)=a2(j/2,i)+a1(jk,i)*f(k)
  20
        continue
        return
        end
С
       subroutine ccfilter(f,a1,a2,raw,col,ltap,mtap)
        integer col, raw, ltap, mtap, jk
        real a1(raw,col),a2(raw,col),f(-20:20)
```

```
do 20 i=1,col
    do 20 j=1,raw
    a2(j,i)=0
```

```
do 20 k=ltap,mtap
              jk=j+k
              if(jk.le.0) jk=raw+jk
              if(jk.gt.raw) jk=jk-raw
              a2(j,i)=a2(j,i)+a1(jk,i)*f(k)
        continue
  20
        return
        end
С
       subroutine rcfilter(f,a1,a2,raw,col,ltap,mtap)
         integer col,raw,ltap,mtap,jk
        real a1(raw, col), a2(raw, col), f(-20:20)
        do 20 i=1, raw
          do 20 j=1,col
              a2(i,j)=0
             do 20 k=ltap,mtap
              jk=j+k
              if(jk.le.0) jk=col+jk
              if(jk.gt.col) jk=jk-col
              a2(i,j)=a2(i,j)+a1(i,jk)*f(k)
  20
        continue
        return
        end
С
        subroutine cinter(in,out,nraw,ncol)
        integer nraw, ncol
        real in(nraw,ncol),out(nraw*2,ncol)
        do 20 j=1,ncol
          do 20 i=1,nraw
               out(2*i-1,j)=in(i,j)
               out(2*i,j)=0.0
               out(2*i,j)=in(i,j)
С
С
               out(2*i-1,j)=0.0
 20
        continue
        return
        end
С
        subroutine rinter(in,out,nraw,ncol)
        integer nraw, ncol
        real in(nraw,ncol),out(nraw,2*ncol)
        do 20 j=1,ncol
          do 20 i=1, nraw
               out(i, 2*j-1) = in(i, j)
               out(i, 2*j) = 0.0
               out(i,2*j)=in(i,j)
С
               out(i,2*j-1)=0.0
С
 20
        continue
```
```
return
end
```

```
subroutine synthesis256(llband,lhband,hlband,hhband,outimg)
        integer raw, col
        parameter(raw=256, col=256)
c input and output images
        real outimg(raw,col)
c these are the four subband
        real llband(raw/2,col/2),lhband(raw/2,col/2)
        , hlband (raw/2, col/2), hhband (raw/2, col/2)
     &
        real lli(raw,col/2),lhi(raw,col/2),hli(raw,col/2),
        hhi(raw,col/2),llo(raw,col/2),lho(raw,col/2),hlo(raw,col/2)
     &
     & hho(raw,col/2)
        real li(raw,col/2),lo(raw,col),hi(raw,col/2),
     & ho(raw,col)
        real
              limg(raw,col),himg(raw,col)
        real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20)
        common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
     &
               ,ltap3,mtap3,ltap4,mtap4
с_
```

```
nraw=raw
ncol=col
call cinter(llband,lli,raw/2,col/2)
call ccfilter(coff4,lli,llo,raw,col/2,ltap4,mtap4)
call cinter(lhband,lhi,raw/2,col/2)
call ccfilter(coff3,lhi,lho,raw,col/2,ltap3,mtap3)
call cinter(hlband,hli,raw/2,col/2)
call ccfilter(coff4,hli,hlo,raw,col/2,ltap4,mtap4)
call cinter(hhband,hhi,raw/2,col/2)
call cinter(hhband,hhi,raw/2,col/2)
call ccfilter(coff3,hhi,hho,raw,col/2,ltap3,mtap3)
```

```
do 10 i=1,raw
          do 10 j=1, col/2
            li(i,j)=llo(i,j)+lho(i,j)
            hi(i,j)=hlo(i,j)+hho(i,j)
  10
        continue
        call rinter(li,lo,raw,col/2)
        call rcfilter(coff4,lo,limg,raw,col,ltap4,mtap4)
        call rinter(hi,ho,raw,col/2)
        call rcfilter(coff3,ho,himg,raw,col,ltap3,mtap3)
        do 20 i=1, raw
          do 20 j=1,col
            outimg(i,j)=1*(limg(i,j)+himg(i,j))
            if(outimg(i,j).gt.255)outimg(i,j)=255
С
С
            if(outimg(i,j).lt.0)outimg(i,j)=0
   20
         continue
        return
        end
С
        subroutine synthesis128(llband,lhband,hlband,hhband,outimg)
        integer raw, col
        parameter(raw=128, col=128)
c input and output images
        real outimg(raw,col)
c these are the four subband
        real llband(raw/2,col/2),lhband(raw/2,col/2)
     &
        , hlband (raw/2, col/2), hhband (raw/2, col/2)
        real lli(raw,col/2),lhi(raw,col/2),hli(raw,col/2),
     & hhi(raw,col/2),llo(raw,col/2),lho(raw,col/2),hlo(raw,col/2)
       hho(raw, col/2)
     &
```

```
real li(raw,col/2),lo(raw,col),hi(raw,col/2),
& ho(raw,col)
```

```
real limg(raw,col),himg(raw,col)
```

real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20)

```
common /a/ coff1,coff2,coff3,coff4,ltap1,mtap1,ltap2,mtap2
,ltap3,mtap3,ltap4,mtap4
```

С

&

```
nraw=raw
        ncol=col
        call cinter(llband,lli,raw/2,col/2)
        call ccfilter(coff4,lli,llo,raw,col/2,ltap4,mtap4)
        call cinter(lhband, lhi, raw/2, col/2)
        call ccfilter(coff3,lhi,lho,raw,col/2,ltap3,mtap3)
        call cinter(hlband,hli,raw/2,col/2)
        call ccfilter(coff4,hli,hlo,raw,col/2,ltap4,mtap4)
        call cinter(hhband, hhi, raw/2, col/2)
        call ccfilter(coff3, hhi, hho, raw, col/2, ltap3, mtap3)
С
        do 10 i=1, raw
          do 10 j=1,col/2
            li(i,j)=llo(i,j)+lho(i,j)
            hi(i,j)=hlo(i,j)+hho(i,j)
        continue
  10
        call rinter(li,lo,raw,col/2)
        call rcfilter(coff4,lo,limg,raw,col,ltap4,mtap4)
        call rinter(hi,ho,raw,col/2)
        call rcfilter(coff3,ho,himg,raw,col,ltap3,mtap3)
        do 20 i=1, raw
          do 20 j=1,col
            outimg(i,j)=1*(limg(i,j)+himg(i,j))
            if(outimg(i,j).gt.255)outimg(i,j)=255
С
            if(outimg(i,j).lt.0)outimg(i,j)=0
С
         continue
   20
        return
        end
С
```

```
subroutine synthesis64(llband,lhband,hlband,hhband,outimg)
         integer raw, col
        parameter(raw=64, col=64)
c input and output images
        real outimg(raw, col)
c these are the four subband
        real llband(raw/2,col/2),lhband(raw/2,col/2)
     &
        , hlband (raw/2, col/2), hhband (raw/2, col/2)
        real lli(raw,col/2),lhi(raw,col/2),hli(raw,col/2),
     & hhi(raw,col/2),llo(raw,col/2),lho(raw,col/2),hlo(raw,col/2)
     & hho(raw,col/2)
        real li(raw,col/2),lo(raw,col),hi(raw,col/2),
     & ho(raw,col)
        real
              limg(raw,col),himg(raw,col)
        real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20)
        common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
     &
               ,ltap3,mtap3,ltap4,mtap4
С
        nraw=raw
        ncol=col
        call cinter(llband,lli,raw/2,col/2)
        call ccfilter(coff4,lli,llo,raw,col/2,ltap4,mtap4)
        call cinter(lhband,lhi,raw/2,col/2)
        call ccfilter(coff3,lhi,lho,raw,col/2,ltap3,mtap3)
        call cinter(hlband,hli,raw/2,col/2)
        call ccfilter(coff4,hli,hlo,raw,col/2,ltap4,mtap4)
        call cinter(hhband, hhi, raw/2, col/2)
        call ccfilter(coff3, hhi, hho, raw, col/2, ltap3, mtap3)
С
```

```
do 10 i=1,raw
```

```
do 10 j=1,col/2
            li(i,j)=llo(i,j)+lho(i,j)
            hi(i,j)=hlo(i,j)+hho(i,j)
  10
        continue
        call rinter(li,lo,raw,col/2)
        call rcfilter(coff4,lo,limg,raw,col,ltap4,mtap4)
        call rinter(hi,ho,raw,col/2)
        call rcfilter(coff3, ho, himg, raw, col, ltap3, mtap3)
        do 20 i=1,raw
          do 20 j=1,col
            outimg(i,j)=l*(limg(i,j)+himg(i,j))
            if(outimg(i,j).gt.255)outimg(i,j)=255
С
С
            if(outimg(i,j).lt.0)outimg(i,j)=0
   20
         continue
        return
        end
```

```
subroutine readimage(pic,nx,ny)
       real pic(nx,ny)
       character*1 pim(65600), header(64)
       integer ilady(256,256)
       character*80 inputimg
       common /image/ inputimg
         open(1,file='/images/'//inputimg,access='direct',
     + form='unformatted', recl=65600)
         write(*,*) 'reading the inputimage','/images/'//inputimg
         read(1,rec=1) (pim(j),j=1,65600)
         read(1,rec=1) (header(j), j=1, 64)
         close (1)
         do 1000 i=1,nx
          k=(i-1)*ny
         l=k+65
         11 = 1 + 255
         mm=0
         do 1001 m=1,11
         mm=mm+1
         ilady(i,mm)=ichar(pim(m))
         if(ilady(i,mm).lt.0) ilady(i,mm)=256+ilady(i,mm)
         pic(i,mm)=float(ilady(i,mm))
1001
         continue
```

```
1000 continue
return
end
```

```
subroutine segma(image,raw,col,output,rmean)
     integer raw, col
     real image(raw,col),output
     double precision s1,s2
     s1=0.0
     s2=0.0
     do 10 i=1, raw
        do 10 j=1,col
         sl=sl+image(i,j)
         s2=s2+(image(i,j))**2
10
     continue
     rmean= s1/(raw*col)
     output=s2/(raw*col)-(rmean)**2
     return
     end
```

```
С
```

С

```
subroutine m_error(image1,image2,raw,col,output)
integer raw,col
real image1(raw,col),image2(raw,col),output
s=0.0
do 10 i=1,raw
    do 10 j=1,col
        s=s+(image1(i,j)-image2(i,j))**2
10 continue
    output=s/(raw*col)
    return
end
```

subroutine write(pic,nx,ny)
real pic(nx,ny)

```
character*1 pim(65600), header(64)
      common /header/ header
        open(1,file='image.img',access='direct',
    + form='unformatted', recl=65600)
      do 10 i=1,64
      pim(i)=header(i)
10
      continue
      do 20 i=1,nx
        do 20 j=1,ny
          ip=int(pic(i,j))
          if(ip.gt.128) ip=ip-256
           mm = 64 + j + (i - 1) * ny
           pim(mm) = char(ip)
20
      continue
      write(1,rec=1) (pim(j),j=1,65600)
      close (1)
      return
      end
```

```
С
```

```
subroutine writeimg(pic,nx,ny,name)
      real pic(nx,ny)
      character*1 pim(65600)
      character*20 name
        open(1,file=name,access='direct',
    + form='unformatted',recl=nx*ny)
      do 20 i=1,nx
        do 20 j=1,ny
          ip=int(pic(i,j))
          if(ip.gt.128) ip=ip-256
           mm=j+(i-1)*ny
           pim(mm) = char(ip)
20
      continue
      write(l,rec=1) (pim(j),j=1,nx*ny)
      close (1)
      return
      end
             í
       function xlog2(x)
```

```
real x
xlog2=alog(x)/alog(2.0)
return
end
```

;

```
с
с
```

```
c quantizer subroutine
c for odd number of levels (3 to 35),0, 64 and 128 levels
        subroutine quantizer(a,nb,c,n,mean,var2,level,raw,col)
c a:real image
c c:integer quantized image
c n:quantizer order
c nb:band number
c mean: image mean
c var: image variance
c level, raw, col: output array dimension
        integer raw,col,level
        real a(raw,col),p(75),pl(75),mean,var
        integer c(level,raw,col)
С
        if(n.eq.0)then
          do 4 i=1, raw
          do 4 j=1,col
            c(nb,i,j)=0
 4
          continue
        go to 44
        endif
С
        m=n
        call readq(n,p,pl)
        if(n.ne.m)then
        write(*,*) 'error in reading quantizer data'
        stop
        end if
         nn = n/2
        var=sqrt(var2)
        do 10 i=1,raw
          do 10 j=1,col
              xx=(a(i,j)-mean)
            do 20 k=nn,1,-1
                px=p(k) *var
              if(xx.gt.px)then
                c(nb,i,j)=k
                go to 100
              else if(xx.lt.(-px))then
                c(nb,i,j) = -k
                go to 100
              endif
20
        continue
        c(nb,i,j)=0
100
        continue
        continue
10
```

÷

```
44
        continue
        return
        end
С
С
   dequantizer subroutine
С
c for odd number of levels (3 to 35),0, 64 and 128 levels
        subroutine dequantizer(b,nb,ic,n,mean,var2,level,raw,col)
c b:real quantized image
c ic:integer quantized image 3D
c n:quantizer order
c nb:band number
c mean: image mean
c var: image variance
        integer raw, col, level
        real b(raw,col),p(75),pl(75),mean,var
        integer ic(level,raw,col)
С
        if(n.eq.0)then
          do 4 i=1,raw
          do 4 j=1,col
            b(i,j)=mean
4
          continue
        go to 44
        endif
С
        m=n
        call readq(n,p,pl)
        if(n.ne.m)then
        write(*,*) 'error in reading quantizer data'
        stop
        end if
        var=sqrt(var2)
        do 10 i=1,raw
          do 10 j=1,col
           k=ic(nb,i,j)
           if(k.gt.0)then
             b(i,j)=pl(k)*var+mean
           else if(k.lt.0)then
             k≐−k
          b(i,j)=-pl(k) *var+mean
           else if(k.eq.0)then
             b(i,j)=mean
        endif
10
        continue
44
        continue
```

```
103
```

```
return
end
```

```
С
С
c quantizer subroutine
c for odd number of levels (3 to 35),0, 64 and 128 levels
        subroutine quant(a,b,c,n,mean,var,raw,col)
c a:real image
c b:real quantized image
c c:integer quantized image
c n:quantizer order
c mean: image mean
c var: image variance
        integer raw, col
        real a(raw,col),b(raw,col),p(75),pl(75),mean,var
        integer c(raw,col)
С
        if(n.eq.0)then
          do 4 i=1,raw
          do 4 j=1,col
            b(i,j)=mean
            c(i,j)=0
 4
          continue
        qo to 44
        endif
С
        m=n
        call readq(n,p,pl)
        if(n.ne.m)then
        write(*,*) 'error in reading quantizer data'
        stop
        end if
         nn = n/2
        var=sqrt(var)
        do 10 i=1, raw
          do 10 j=1,col
              xx=(a(i,j)-mean)
            do 20 k=nn,1,-1
                px=p(k) *var
              if(xx.gt.px)then
                b(i,j)=(pl(k)*var)+mean
                c(i,j)=k
                go to 100
              else if(xx.lt.(-px))then
                b(i,j)=(-pl(k)*var)+mean
                c(i,j) = -k
                go to 100
```

```
20
        continue
        b(i,j)=0+mean
        c(i,j)=0
 100
        continue
 10
        continue
 44
        continue
        return
        end
С
c read subroutine which
c read quantizer data
        subroutine readq(n,p,pl)
        real p(75),pl(75)
       open(32,file='.../quantizer.dir/quant2.dat'
              ,access='direct',form='unformatted',recl=1000)
     &
        nn=(n)/2
        if(n.le.35.and.n.ge.3)then
           number=n/2+1
        else if(n.eq.64) then
           number=20
        else if(n.eq.128) then
           number=21
        else
           write(*,*) 'the order of the filter is rong'
           STOP
        endif
        read(32, rec=number) n, (p(i), i=1, nn), (pl(i), i=1, nn)
        close(32)
        return
        end
С
С
С
c This subroutine allocate the bitrate for each subband
c by finding the suitable quantizers
c Number of band is 64
С
        subroutine entropy(iq,var,mean)
        parameter(nn=64)
        real var(nn)
        real mean(nn),br(nn)
        integer iq(nn)
```

endif

double precision xx

```
logical ql(nn)
         common /bitrate/ temprav
        prav=temprav
С
С
  check for negative br and cancell the corespondant band
С
        do 10 i=1,nn
           ql(i)=.true.
 10
     continue
 100
        continue
        call gain(var, ql, xx, nn, ncount)
        rav=prav*nn/real(ncount)
        do 30 i=nn,1,-1
            if(ql(i)) then
              x1=var(i)/xx
              br(i) = rav + 0.5 \times x \log_2(x1)
              if(br(i).lt.0) then
                br(i)=0
                ql(i)=.false.
                iq(i)=0
                go to 100
              endif
            endif
        continue
 30
С
c optimize the bit rates
С
 200
        continue
С
        write(*,*) 'rav=',rav
        call gain(var,ql,xx,nn,ncount)
        write(*,*) xx
С
        do 40 i=nn,1,-1
          if(ql(i)) then
             xl=var(i)/xx
             br(i) = rav + 0.5 \times log2(x1)
             in=int(2.0**br(i))
             iq(i) = (in+1)/2*2-1
             if(iq(i).lt.3) then
                 iq(i)=0
             else if(iq(i).ge.128) then
                 iq(i)=128
             else if(iq(i).ge.64) then
                 iq(i)=64
             else if(iq(i).ge.35) then
                 iq(i) = 35
             endif
             ql(i)=.false.
             if(ncount.gt.2)then
                if(iq(i).ge.3) then
```

```
xiq=real(iq(i))
                  xr = xloq2(xiq)
                else
                  xr=0
                endif
                 xc=real(ncount)
                 xrav=rav
                 rav=(xrav*xc-xr)/(xc-1.0)
С
         write(55,*) br(i),iq(i),xr,var(i),rav
              go to 200
            endif
          endif
        continue
 40
c calculate the actuall bit rate
        sum=0.0
        do 50 i=1,nn
           write(55,*) i,iq(i),var(i),mean(i)
С
           if(iq(i).gt.1) then
             br(i)=xlog2(real(iq(i)))
           else
             br(i)=0
           endif
           sum=sum+br(i)
 50
        continue
        rav=sum/nn
С
        write(55,*) 'rav =',rav
        return
        end
С
С
c This subroutine calculates the gain
        subroutine gain(var,ql,xx,nband,ncount)
        real var(nband)
        logical gl(nband)
        double precision xx
        xx=1.0
        count=1.0
        do 10 i=1, nband
          if(ql(i))then
           xx=xx*var(i)
           count=count+1
          endif
 10
        continue
        xx=xx**(1.0/count)
        ncount=int(count+.5)
```

return end

```
С
c this subroutine calculate the entropy of each band
c and find the probability of each code
        subroutine bitrates(ic,ln,entropy,pr,level,raw,col)
        integer ic(level,raw,col),ln,level,raw,col
        real entropy(level), sum(-64:64), pr(-64:64)
        do 20 n=-64,64
         sum(n)=0.0
 20
        continue
        do 10 i=1, raw
        do 10 j=1,col
           k=ic(ln,i,j)
           sum(k) = sum(k) + 1
 10
        continue
        entropy(ln)=0.0
        total=real(raw*col)
        do 30 n = -64, 64
           pr(n)=sum(n)/total
           if(pr(n).qt.0) then
             br=pr(n) *xlog2(1.0/pr(n))
             entropy(ln)=entropy(ln)+br
           endif
 30
        continue
        write(*,*) 'entropy = ',entropy
С
        return
        end
С
```

subroutine convert(pr,prb,prnum,ns)

```
real pr(-64:64),prb(128)
integer prnum(128),ns
n=0
do 10 i=-64,64
    if(pr(i).gt.0.0) then
        n=n+1
        prb(n)=pr(i)
        prnum(n)=i
```

```
endif
 10 continue
     ns=n
        return
     end
С
С
c Program to sort the probability table
С
        subroutine sort(pr,prnum,ns)
        real pr(128), temppr
        integer prnum(128),tempprn
С
c sorting algorithm
        do 10 k=ns-1,1,-1
          do 20 i=1,k
           if(pr(i).lt.pr(i+1)) then
             temppr=pr(i)
             pr(i) = pr(i+1)
             pr(i+1)=temppr
             tempprn=prnum(i)
             prnum(i)=prnum(i+1)
             prnum(i+1)=tempprn
           endif
 20
          continue
        continue
 10
        return
        end
С
С
c subroutine to find the huffman codebook
С
     subroutine hcbook(prb,codbook,ns)
     real pr(128*129/2),prb(128)
        integer pointer(128*129/2)
     character*1 code(128*129/2)
    character*80 codbook(128),tempc,codebook
        logical flag
        do 19 k=1,128
19
           codbook(k) = '
        do 50 k=1,ns
           pr(k) = prb(k)
```

```
pointer(k)=0
           code(k) = 'x'
 50 continue
        do 51 k=ns+1,128*129/2
          pr(k)=0.0
          pointer(k)=0
          code(k) = 'x'
 51
        continue
          code(ns*(ns+1)/2-1)='1'
          code(ns*(ns+1)/2-2)='0'
        last=0
        do 20 j=1,ns-2
          nn=ns-j
          xx=pr(last+nn)+pr(last+nn+1)
          code(last+nn)='0'
          code(last+nn+1)='1'
          flag=.false.
          do 10 i=last+1,last+nn
             if(flag) then
              pr(i+nn+1)=pr(i-1)
              pointer(i-1)=i+nn+1
            else
               if(xx.gt.pr(i))then
               flag=.true.
                pr(i+nn+1) = xx
                pointer(last+nn)=i+nn+1
                pointer(last+nn+1)=i+nn+1
              else
                pr(i+nn+1)=pr(i)
                pointer(i)=i+nn+1
              endif
            endif
 10
          continue
          last=last+nn+1
 20
        continue
        do 11 i=1,ns
           k=i
          codbook(i)='
          tempc=''
          do 12 j=1,ns-1
          if(code(k).eq.'0'.or.code(k).eq.'1')then
            write(*,*) code(k)
С
            codebook=code(k)//tempc
            write(*,*) codebook
С
            tempc=codebook
```

L

۰.

```
endif
С
           write(*,*) 'k=',k
           k=pointer(k)
           write(*,*) 'k=',k
С
 12
         continue
         codbook(i)=codebook//'
                                    1
 11
        continue
        write(*,*) (pr(i),i=1,14)
write(*,*) (pointer(i),i=1,14)
С
С
        write(*,*) (code(i),i=1,14)
С
С
        write(51,*) (codbook(i),i=1,5)
     return
     end
C_
c subroutine to code an image to hoffman
        subroutine codehc(image,level,codbook,prnum,ns,hc,ll,lev
                            ,raw,col)
     &
        integer lev, raw, col
        integer image(lev,raw,col),prnum(128),ns
        character*80 codbook(128), ccode
        character*1 hc(77000)
        do 5 i=1,77000
             hc(i) = ' '
 5
     continue
        ncount=0
        do 10 i=1, raw
        do 10 j=1,col
          icode=image(level,i,j)
          do 20 k=1,ns
              if(prnum(k).eq.icode)then
                ncode=k
                qo to 21
              endif
20
          continue
          write(*,*) 'ERROR IN READING CODBOOK'
          stop
21
          continue
```

```
111
```

```
length=index(codbook(ncode),' ')-1
          ccode=codbook(ncode)
          do 30 k=1,length
            ncount=ncount+1
            hc(ncount) = ccode(k:k)
 30
          continue
 10
        continue
        ll=ncount
        do 50 i1=1,77000
        if(hc(i1).ne.'0'.and.hc(i1).ne.'1') go to 60
 50 continue
 60
        112 = i1 - 1
        if(ll2.ne.ll) then
         write(*,*) ' length of hc is not correct',112
        endif
        return
        end
С
c program to decode the hoffman code
С
        subroutine decodehc(image,level,codbook,prnum,ns,hc,ll,
     &
                             lev,raw,col)
        integer lev, col, raw
        integer image(lev,raw,col),icode(128*128),prnum(128),ns
        character*80 codbook(128),ccode
        character*1 hc(77000), code1, code2
        logical flag(128)
        npexel=0
        ncount=0
        do 5 i=1,ns
          flag(i)=.true.
          if(flag(i)) then
            write(76,*) 'true'
С
       endif
5
        continue
        do 10 ii=1,11
          code1=hc(ii)
```

.

	ncount=ncount+1
	do 20 k=1,ns
	if(flag(k)) then
С	write(75,*) 'true'
	ccode=codbook(k)
	code2=ccode(ncount:ncount)
	if(code2.eg.code1) then
	length=index(ccode ( ()-1
C	write(74 *) length
C	if(longth og ngount) go to 30
	olco
	$e_{1Se}$
	IIdy(K)Idise.
	enull
20	continuo
20	continue
	ao to 10
	90 20 10
30	ncount=0
	do 40 $i=1.ns$
	flag(i)=.true.
40	
	npexel=npexel+1
с	write(73,*) npexel
-	icode(npexel)=prnum(k)
10	continue
	do 50 i=1,raw
	do 50 j=1,col
	<pre>image(level,i,j)=icode((i-1)*col+j)</pre>
С	<pre>write(72,*) image(level,i,j)</pre>
50	continue
	return
	end
-	
C	······································

```
subroutine unifquan(pic,ll,image,order,min,step,level,raw,co
integer raw,col
integer image(level,raw,col),order
real pic(raw,col)
real min,max
```

```
c assume the minimum value and the maximum
    min =pic(1,1)
```

```
\max = pic(1,1)
    check if all pixels values are between min and max
С
       do 10 i=1, raw
         do 10 j=1,col
           if(pic(i,j).gt.max) then
             max = pic(i,j)
           else if(pic(i,j).lt.min)then
             \min = pic(i,j)
           end if
       continue
 10
    rescale the image and convert it into image array
С
         step = (max-min)/real(order)
         ixx=order/2
         do 20 i=1,raw
           do 20 j=1,col
             image(ll,i,j) =int((pic(i,j)-min)/step+.5)-ixx
 20
         continue
        return
        end
С
       subroutine unifdequan(pic,ll,image,order,min,step,level,raw,
       integer raw, col
       integer image(level,raw,col),order
       real pic(raw,col)
       real step, min
    rescale the image and convert it into pic array
С
         ixx=order/2
         do 20 i=1, raw
           do 20 j=1,col
            pic(i,j)=real(image(ll,i,j)+ixx)*step+min
         continue
 20
        return
        end
С
```

ł

```
subroutine dpcdecoder(e,img,raw,col)
        parameter(ii=32,jj=32)
        integer raw, col
        real e(raw,col),img(raw,col)
        real temp(0:ii+1,0:jj+1)
        do 10 i=0,raw+1
        do 10 j=0,col+1
            temp(i,j)=0
    continue
 10
        do 20 i=1, raw
        do 20 j=1,col
           temp(i,j)=e(i,j)+0.25*(temp(i,j-1)+temp(i-1,j-1)+
     &
                     temp(i-1,j)+temp(i-1,j+1))
           img(i,j)=temp(i,j)
 20
     continue
        return
     end
C
С
c program for noise less dpcm coding algorithm
С
          subroutine qdpcm(u,ie,iq,var,mean,raw,col,p,pl)
        parameter(ii=32,jj=32)
        integer raw, col
     real u(raw,col),ud(0:ii+1,0:jj+1),udb
        real e,ed
        real p(75),pl(75)
        integer ie(4, raw, col)
        real mean, var
        do 10 i=0, raw+1
        do 10 j=0,raw+1
         ud(i,j)=0.0
 10 continue
          a1=0.25
        a2 = 0.25
        a3=0.25
        a4=0.25
     do 20 m=1,raw
     do 20 n=1,col
         udb=a1*ud(m-1,n)+a2*ud(m,n-1)+a3*ud(m-1,n-1)+a4*ud(m-1,n+1
         e=u(m,n)-udb
         call quantizer1(e,ed,ie(1,m,n),iq,mean,var,p,pl)
```

```
ud(m,n) = udb+ed
 20 continue
     return
     end
C_
c
c quantizer subroutine
c for odd number of levels (3 to 35),0, 64 and 128 levels
        subroutine quantizer1(a,b,c,n,mean,var2,p,pl)
c a:real image
c b: real quantized image
c c:integer quantized
c n:quantizer order
c mean: image mean
c var: image variance
С
        real a,p(75),pl(75),mean,var
         integer c
С
         if(n.eq.0)then
             c=0
        ao to 44
        endif
С
         nn = n/2
        var=sqrt(var2)
               xx=(a-mean)
             do 20 k=nn,1,-1
                 px=p(k) *var
               if(xx.gt.px)then
                 b=(pl(k)*var)+mean
                 c=k
                 go to 100
               else if(xx.lt.(-px))then
                 b=(-pl(k)*var)+mean
                 c = -k
                   go to 100
              endif
 20
        continue
        b=mean
        c=0
 100
        continue
 44
        continue
        return
        end
С
С
```

,

```
c program for dpcm coding algorithm
     subroutine dpcm1(x,y,raw,col)
        parameter(ii=32,jj=32)
        integer raw, col
     real x(raw,col),temp(0:ii+1,0:ii+1),y(raw,col)
        write(*,*) 'dpcm1'
        do 10 i=0, ii+1
        do 10 j=0,jj+1
           if(i.eq.0.or.j.eq.0.or.i.gt.ii.or.j.gt.jj) then
              temp(i,j)=0
           else
              temp(i,j)=x(i,j)
           endif
    continue
10
     do 20 i=1,raw
     do 20 j=1,col
       y(i,j) = temp(i,j) - 0.25 * (temp(i-1,j) + temp(i,j-1))
                  +temp(i-1,j-1)+temp(i-1,j+1))
     &
     continue
20
     return
     end
C_
               ÷
```

,

## Bibliography

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