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A comparative study of image coding techniques : filter banks vs. discrete cosine transform

Hosam Fawzi Mutlag New Jersey Institute of Technology

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Abstract

Title of Thesis: A Comparative Study of Image Coding Techniques: Filter Banks vs. Discrete Cosine Transform.

Hosam Mutlag, Master of Science, 1991

Thesis directed by: Dr. Ali N. Akansu

Subband coding of still image frames using Binomial PR-QMF has been presented in this thesis. Simulation results have shown that the performance of subband coding using a low complexity coder is practically the same as the performance of the industry standard (8×8) DCT based image coder.

A low bit rate adaptive video coding technique is also introduced in this thesis. The redundancy within adjacent video frames is exploited by motion compensated interframe prediction. The Motion Compensated Frame Difference (MCFD) signals are filtered by employing Binomial PR-QMF structure into four subbands. Then, the subbands are quantized using an efficient Motion Based Adaptive Vector Quantization (MBAVQ) algorithm. Here, the adaptation scheme is based on block motion vectors rather than local signal energy which was used in earlier works of several researchers. The new technique results in a reduction in bit rate by nearly (40%) due to the drop of the extra bits used for local variances. Moreover, for the video test sequances considered, MBAVQ method gives superior SNR results over local Variance Based Adaptive Vector Quantiztion (VBAVQ) scheme especially for high motion frames.

 $\frac{1}{2} \int$ A Comparative Study of Image Coding Techniques: Filter Banks vs. Discrete Cosine Transform

> by Hosam Fawzi Mutlag

Thesis submitted to the Faculty of the Graduate School of the New Jersey Institute of Technology in partial fulfillment of the requirements for the degree of Master of Science in Electrical Engineering

APPROVAL SHEET

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Acknowledgement

I would like to express my gratitude to Dr. Ali Akansu for his valuable contribution, advices, patience and understanding. Furthermore, I appreciate his support and encouragement during the entire research period.

I am greateful to Dr. Yeheskil Bar-Ness and Dr. Zoran Siveski for their effort and time in reviewing this work.

Thanks and regards to my parents for their moral support and love.

Last but not least, I would like to thank the members of the Center for Communication and Signal Processing Research at New Jersey Institute of Technology and to all my friends for their help which made the period of this work a pleasent one.

To my parents

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Chapter 1 Introduction

Typical video has spatial resolution of approximately 512×512 pixels per frame. At 8 bits per pixel per color channel and 30 frames per second, this translates into a rate of nearly 180×10^6 bits/s. The large channel capacity and memory requirements for digital image transmission and storage makes it desirable to consider data compression techniques. Image data compression is concerned with minimizing the number of bits required to represent an image with an acceptable visual quality. Compression can be achieved by transforming the signal, projecting it on a basis of functions, and then encoding the transform coefficients. These transforms vary from the conventional block transforms to ideal subband filter banks. The principle of subband coding has recently been successfully applied to data compression of both still images and video. In subband coding, the signal to be coded is decomposed into narrow band pass signals (subbands). Each subband is then accordingly subsampled and encoded with a bit rate matched to the signal statistics in that subband.

Subband coding was first introduced by Crochiere, et. al.[1] in speech coding and then extended to multidimensional signals by Vetterli[2]. Then, this concept was applied to the coding of images [3] [4]. Perfect Reconstruction Quadrature Mirror Filters (PR-QMF) have been proposed as structures suitable for subband coding [5] [6]. These filters employing a tree decomposition structure provide a basis for a multiresolution signal representation [7] [8]. Recently, discrete wavelet transforms have been proposed as a new approach for multi-resolution signal decomposition. More recently, Binomial QMF-Wavelet Transform has been proposed for multiresolution signal decomposition [10] [11].

In this thesis, we have studied the performance of Binomial QMF in subband coding with comparison to the Discrete Cosine Transform (DCT). Simulation programs were developed for a complete subband coding system (codec) consisting of the following four distinct parts:

- An analysis filter bank splitting the input signal into subbands.
- An encoder which consists of:
	- 1. Bit allocation algorithm.
	- 2. Quantizers.
	- 3. Huffman encoder.
- A decoder whose purpose is to produce an approximation to the original subband signals.
- The synthesis filter bank that combines the decoded subband signals, to reconstruct the received signal.

A similar DCT based codec is also developed for comparison purposes.

We also propose a new Motion Based Adaptive Vector Quantization (MBAVQ) approach for subband video coding. In video signals, interframe images have significant frame to frame redundancy. For that reason, in video coding techniques, video frames are motion compensated using an efficient search algorithm to remove temporal redundancies. The resulting prediction error signals, Motion Compensated Frame Difference (MCFD), are coded using an efficient coding scheme. In our video coding scheme, the MCFD signals are divided into four subbands. Least significant band (highest frequency), is neglected and the other bands are vector quantized adaptively.

This thesis is organized as follows. Chapter 2 describes the Binomial-QMF which have been used. In the following chapter, Discrete Cosine Transform is discussed. In chapter 4, optimum bit allocation and gain of transform coding over PCM is explained. Quantization is reviewed in chapter 5. The next chapter discusses source entropy and Huffman coding. In chapter 7, Motion Based Adaptive Vector Quantization (MBA-VQ) video coding scheme is introduced. Also, a comprehensive set of simulation results is presented in this chapter. Finally, we conclude the thesis with a discussion of the results and future research.

Chapter 2 PR_QMF

2.1 Theoretical **Derivations**

The QMF bank is a multirate digital filter bank. There are decimators in the system which down-sample a signal sequence, and there are interpolators which perform upsampling. The input-output relation for a two-fold decimator can be written in the transform domain as

$$
Y(e^{jw}) = \frac{1}{2} \left[X(e^{jw/2}) + X(-e^{jw/2}) \right]
$$
 (2.1)

where $Y(e^{jw})$ has a period of 2π . The effect of compression in time domain is an expansion (or stretching) in frequency domain. The transform domain relation for a two-fold interpolator is

$$
Y(e^{jw}) = X(e^{jw2})
$$

\n
$$
Y(z) = X(z^2)
$$
\n(2.2)

The interpolator causes compression in the frequency domain.

Let us consider Figure (2.1) in which a two-channel QMF system is shown. Based on relations (2.1) and (2.2), it is possible to express $\hat{X}(z)$ in Figure 1 as

$$
\hat{X}(z) = \frac{1}{2} \left[H_0(z) F_0(z) + H_1(z) F_1(z) \right] X(z) \n+ \frac{1}{2} \left[H_0(-z) F_0(z) + H_1(-z) F_1(z) \right] X(-z)
$$
\n(2.3)

The reconstructed signal can in general be subject to three types of distortions:

- Aliasing distortion.
- Phase distortion.
- Amplitude distortion.

The second term in (2.3) represents the effects of aliasing and imaging. This term can be made to drop simply by choosing the synthesis filters to be

$$
F_0(z) = -H_1(-z) \tag{2.4}
$$

$$
F_1(z) = H_0(-z) \tag{2.5}
$$

The QMF bank becomes a linear and time-invariant system with transfer function

$$
T(z) = \frac{\hat{X}(z)}{X(z)} = \frac{1}{2} \left[H_0(z) H_1(-z) - H_1(z) H_0(-z) \right]
$$
 (2.6)

If $|T(e^{jw})|$ = constant for all w, then there is no amplitude distortion. Also, if $T(z)$ is a linear-phase FIR function, then $\arg[T(e^{jw})] = kw$, and there is no phase distortion. This means $T(z)$ is a delay, i.e., $T(z) = Cz^{-n_0}$, so that the reconstructed signal $\hat{x}(n)$ is a delayed version of $x(n)$.

Smith and Barnwell [6] have shown that one can simultaneously eliminate both amplitude and phase distortions by choosing $(2N-1)$ odd and

$$
H_1(z) = z^{(2N-1)}H_0(-z^{-1})
$$
\n(2.7)

Therefore,

$$
T(z) = \frac{1}{2} z^{-(2N-1)} \left[H_0(z) H_0(z^{-1}) - H_0(-z) H_0(-z^{-1}) \right] = C z^{-(2N-1)} \tag{2.8}
$$

The perfect reconstruction requirement of a two-band QMF reduces to[5]

$$
Q(z) = H(z)H(z^{-1}) + H(-z)H(-z^{-1}) = C
$$

= $R(z) + R(-z)$ (2.9)

where $H(z)$ is a low-pass filter of length 2N.

The PR requirement, Eq.(2.9), can be readily recast in an alternate, time domain form [10]. First, one notes that $R(z)$ is a spectral density function and hence is representable by a finite series of the form

$$
R(z) = \gamma_{2N-1} z^{2N-1} + \gamma_{2N-2} z^{2N-2} + \ldots + \gamma_0 z^0 + \ldots + \gamma_{2N-1} z^{-(2N-1)} \tag{2.10}
$$

Then

$$
R(-z) = -\gamma_{2N-1}z^{2N-1} + \gamma_{2N-2}z^{2N-2} - \ldots + \gamma_0z^0 - \gamma_1z^{-1} \ldots - \gamma_{2N-1}z^{-(2N-1)} \tag{2.11}
$$

Therefore $Q(z)$ consists only of even powers of *z*. To force $Q(z) = C$ which is a constant, it suffices to make all even indexed coefficients in $R(z)$ equal to zero, except for n=0. However, the γ_n coefficients in *R(z)* are simply the samples of the autocorrelation $\rho(n)$ given by

$$
\rho(n) = \sum_{k=0}^{2N-1} h(k)h(k+n) = \rho(-n)
$$
\n
$$
\stackrel{\text{def}}{=} h(n) \odot h(n)
$$
\n(2.12)

where \odot indicates a correlation operation. This follows from the z-transform relationships

$$
R(z) = H(z)H(z^{-1}) \longleftrightarrow h(n) * h(-n) = \rho(n) \tag{2.13}
$$

where $\rho(n)$ is the convolution of $h(n)$ with $h(-n)$, or equivalently, the time autocorrelation, Eq.(2.12). Hence, we need to set $\rho(n) = 0$ for n even, and $n \neq 0$. Therefore.

$$
\rho(2n) = \sum_{k=0}^{2N-1} h(k)h(k+2n) = 0, \qquad n \neq 0 \tag{2.14}
$$

If the normalization is imposed,

$$
\sum_{k=0}^{2N-1} |h(k)|^2 = 1\tag{2.15}
$$

one obtains the PR requirement in time

$$
\sum_{k=0}^{2N-1} h(k)h(k+2n) = \delta_n \tag{2.16}
$$

2.2 The Binomial Family

The binomial family of orthogonal sequences [13] [14] is generated by successive differencing of the binomial sequence, which is defined on the finite interval $[0, N]$ by

$$
x_0(k) = \begin{cases} \begin{pmatrix} N \\ k \end{pmatrix} = \frac{N!}{(N-k)!k!} & 0 \le k \le N \\ 0 & \text{otherwise} \end{cases}
$$
 (2.17)

The other members of the binomial family are obtained from

$$
x_r(k) = \nabla^r \left(\begin{array}{c} N-r \\ k \end{array} \right) \quad r = 0, 1, \dots, N \tag{2.18}
$$

where

$$
\nabla f(n) = f(n) - f(n-1)
$$

this is the backward difference operator. Taking successive differences yields Bonomial function family of length $(N+1)$

$$
x_r(k) = {N \choose k} \sum_{\nu=0}^r (-2)^{\nu} {r \choose \nu} \frac{k^{(\nu)}}{N^{(\nu)}} \qquad (2.19)
$$

where $k^{(\nu)}$ is the forward factorial function, a polynomial in *k* of degree ν

$$
k^{(\nu)} = \begin{cases} k(k-1)\dots(k-\nu+1) & \nu \ge 1\\ 1 & \nu = 1 \end{cases}
$$
 (2.20)

This family of binomially-weighted polynomials has a number of properties. Taking *z* transform, we obtain

$$
X_0(z) = (1 + z^{-1})^N
$$

\n
$$
X_r(z) = (1 - z^{-1})^r (1 + z^{-1})^{N-r}
$$
\n(2.21)

The binomial matrix X is the $(N + 1) \times (N + 1)$ matrix

$$
X=[x_r(k)]
$$

where $x_r(k)$ is the entry in the r^{th} row and k^{th} column. The salient property of this matrix is that the rows are orthogonal to the columns,

$$
\sum_{k=0}^{N} x_r(k)x_k(s) = (2)^N \delta_{r-s}
$$
\n(2.22)

or

$$
X^2 = (2)^N I \tag{2.23}
$$

Additionally, the Bionomial filters are linear phase quadrature mirror filters. From eq. (2.21), we see that

$$
X_r(-z) = X_{N-r}(z) \tag{2.24}
$$

which implies

$$
(-1)^{k} x_{r}(k) = x_{N-r}(k) \quad r = 0, 1, \cdots, N \tag{2.25}
$$

Also,

$$
z^{-N}X_r(z^{-1}) = (-1)^r X_r(z)
$$
\n(2.26)

implies

$$
x_r(N-k) = (-1)^r x_r(k) \tag{2.27}
$$

Equations (2.25) and (2.27) demonstrate the symmetry and asymmetry of the rows and columns of the Bionomial matrix X . From equation (2.25) , we can infer that the complementary filters X_r and $X_{N-r}(z)$ have magnitude responses which are mirror images about $\omega=\pi/2$

$$
|X_r(\exp j(\frac{\pi}{2} - \omega))| = |X_{N-r}(\exp j(\frac{\pi}{2} + \omega))|
$$
 (2.28)

Moreover, the cross-correlation of the sequence $x_r(n)$, and $x_s(n)$ is defined as

$$
\rho_{rs} = x_r(n) * x_s(-n) = \sum_{k=0}^{N} x_r(k) x_s(n+k) \rightleftharpoons R_{rs}(z) \tag{2.29}
$$

and

$$
R_{rs}(z) = X_r(z^{-1})X_s(z)
$$
\n(2.30)

Now for any real crosscorrelation,

$$
\rho_{rs}(-n) = \rho(n) \qquad \forall s, r \tag{2.31}
$$

Also, the quadrature mirror property of $Eq.(2.25)$ implies that

$$
\rho_{rs}(n) = -\rho_{rs}(n) \qquad (s-r) \text{ is odd}
$$

$$
\rho_{rs}(n) = \rho_{rs}(n) \qquad (s-r) \text{ is even} \qquad (2.32)
$$

These properties are subsequently used in driving the perfect reconstruction Bionomial QMF in the next section.

2.3 The Binomial QMF

Now, it is a straight forward matter to impose PR condition of Eq.(2.16) on the binomial family [10]. First, the half-band filter is

$$
h(n) = \sum_{r=0}^{\frac{N-1}{2}} \theta_r x_r(n)
$$

or in the *z* transform

$$
H(z) = \sum_{r=0}^{\frac{N-1}{2}} \theta_r (1 + z^{-1})^{N-r} (1 - z^{-1})^r = (1 + z^{-1})^{(N+1)/2} F(z)
$$
(2.33)

where $F(z)$ is FIR filter of order $(N-1)/2$. For convenience, take $\theta_0=1$, and later impose the normalization of Eq.(2.15). Substituting (2.33) into (2.12) gives

$$
\rho(n) = \left(\sum_{r=0}^{\frac{N-1}{2}} \theta_r x_r(n)\right) \odot \sum_{s=0}^{\frac{N-1}{2}} \theta_s x_s(n)
$$

\n
$$
= \sum_{r=0}^{\frac{N-1}{2}} \sum_{s=0}^{\frac{N-1}{2}} \theta_r \theta_s [x_r(n) \odot x_s(n)]
$$

\n
$$
= \sum_{r=0}^{\frac{N-1}{2}} \sum_{s=0}^{\frac{N-1}{2}} \theta_r \theta_s \rho_{rs}(n)
$$

\n
$$
= \sum_{r=0}^{\frac{N-1}{2}} \theta_r^2 \rho_{rr}(n) + \sum_{r=0, r \neq s}^{\frac{N-1}{2}} \sum_{s=0}^{\frac{N-1}{2}} \theta_r \theta_s \rho_{rs}(n)
$$
(2.34)

Eq.(2.32) implies that the second summation in Eq.(2.34) has only terms where the indices differ by an even integer. Therefore the autocorrelation for the binomial halfband low-pass filter is

$$
\rho(n) = \sum_{n=0}^{\frac{N-1}{2}} \theta_r^2 \rho_{rr}(n) + 2 \sum_{l=1}^{\frac{N-3}{2}} \sum_{\nu=0}^{\frac{N-1}{2}-2l} \theta_\nu \theta_{\nu+2l} \rho_{\nu,\nu+2l}(n) \tag{2.35}
$$

Finally, the PR requirement is

$$
\rho(n) = 0, \qquad n = 2, 4, \dots, N - 1 \tag{2.36}
$$

This condition gives a set of $\frac{N-1}{2}$ nonlinear algebraic equations, in the $\frac{N-1}{2}$ unknowns $\theta_1, \theta_2 \ldots, \theta_{\frac{N-1}{2}}$. From these, we can obtain the corresponding Binomial PR-QMFs.

These filters can be implemented using either the purely FIR structure, or the pole-zero cancellation configuration. The latter is shown in Fig.2 for $N = 5$. Wherein both low-pass and high-pass filters are simultaneously realized. Coefficient θ_0 can be taken equal to unity, leaving only θ_1 and θ_2 as tap weights. These are the only multiplications needed when using the Binomial network as the half-band QMF rather than the six $h(n)$ weights in a transversal structure.

2.4 M-Band Tree Decomposition

After a given signal $x(t)$ is sampled at f_s to give a signal $X(n)$, and split it into two signals $X_L(n)$ and $X_H(n)$, with the reduction of the sampling rate to $f_s/2$, the decomposition can be extended to more than two subbands by processing these signals in the same manner as the initial signal $X(n)$. Four signals are thus obtained with a reduction of the sampling rate to $f_s/4$. The decomposition/reconstruction can be generalized by repeating n times the previously described splitting using n-stage tree decomposition. In figure (2.3), a four subband decomposition is shown. The signal $X(n)$ is split into four signals X_{LL} , X_{LH} , X_{HL} , and X_{HH} .

2.5 Two Dimensional Separable Case

The extention of the QMF decomposition/reconstruction concept to two dimentional case is relatively straight forward in separable filter case. The conditions required for splitting two dimensional signals into more than two bands, with alias-free recon-

Figure 2.1: Two Channel QMF Bank.

Figure 2.2: Low-pass and high-pass QMF filters from Binomial Network.

Figure 2.3: Four-band tree decomposition for one dimensional signal $X(n)$

struction of the original signal, is the separability feature of the filters, that is

$$
h(n_1, n_2) = h_1(n_1)h_2(n_2)
$$
\n(2.37)

These separability characteristics of the filter provide an alternative method of implementation for 2D-QMF banks. The analysis and synthesis are done as shown in figure (2.4) and figure (2.5) where the structure consists of one-dimensional filters. The computation is performed first along one axis (rows) and then along the other axis (columns). It can be shown that application of this filter structure will permit an alias-free reconstruction of the input signal at the receiver. The detailed analysis/synthesis and perfect reconstruction are given in [3]. Figure (2.4) shows four-band tree decomposition of a two-dimensional signal $X(n, m)$ while figure (2.5) shows the construction tree.

Figure 2.4: Four-band tree decomposition of a two-dimensional signal $X(n, m)$

Figure 2.5: Four-band reconstruction of a two-dimensional signal $X(n, m)$

Chapter 3 Discrete Cosine Transform

The term image transform usually refers to a class of unitary matrices used for representing images. Just as one-dimensional signal can be represented by an orthogonal series of basis functions, an image can also be expanded in terms of a discrete set of basis arrays called basis images. These basis images can be generated by unitary matrices. To make transform coding practical, a given image is divided into small rectangular blocks, and each block is transform coded independently. For an $N \times M$ image divided into NM/pq blocks, each of size $p \times q$, the main storage requirements for implementing the transform are reduced by afactor of *MN/pq.* The computational load is reduced by a factor of $\log_2 MN / \log_2 pq$ for fast transform requiring $\propto N \log_2$ operations to transform an $N \times 1$ vector. For 256 \times 256 images divided into 8 \times 8 blocks, these factors are 1024 and 2.66, respectively. Although the operation count is not greatly reduced, the complexity of the hardware for implementing small size transform is reduced significantly. However, smaller block sizes yield lower compression. The $N \times N$ discrete cosine transform (DCT) matrix $C = \{c(k, n)\}\$, is defined $as[16]$

$$
c(k,n) = \begin{cases} \frac{1}{\sqrt{N}} & k = 0, 0 \le n \le N-1\\ \sqrt{\frac{2}{N}} \cos \frac{\pi (2n+1)}{2N} & 1 \le k \le N-1, 0 \le n \le N-1 \end{cases}
$$
(3.1)

The two-dimensional discrete cosine transform of a two-dimensional sequence

 $\{u(m,n), 0 \leq m \leq N-1, 0 \leq n \leq N-1\}$, is obtained by

$$
v(k,l) = \sum_{m=0}^{N-1} \sum_{n=0}^{N-1} c(k,m)u(m,n)c(l,n)
$$
\n(3.2)

The inverse transformation is given by

$$
u(m,n) = \sum_{k=0}^{N-1} \sum_{l=0}^{N-1} c(k,m)v(k,l)c(l,n)
$$
\n(3.3)

or in vector notation

$$
V = CUC^T \tag{3.4}
$$

and

$$
U = C^T V C \tag{3.5}
$$

The cosine transform is real and orthogonal, that is,

$$
C = C^* \rightarrow C^{-1} = C^T \tag{3.6}
$$

To implement one-dimensional DCT of size *N* using fast transform algorithms, it requires *0(N* log *N)* operations, where one operation is a real multiplication and a real addition. For two-dimensional DCT of size $N \times N$, the transformation is equivalent to 2N one-dimensional DCT. Hence, the total number of operation is $O(N^2 \log_2 N)$.

Chapter 4

Optimum Bit Allocation and Gain of Transform Coding

4.1 Bit Allocation

Bit allocation is a major concern in any coding scheme where a given quota of bits must be efficiently distributed among a number of different quantizers. Encoding in subbands offers several advantages. By appropriately allocating the bits in different bands, the number of quantizer levels and hence reconstruction error variance can be separateley controlled in each band and the shape of the overall reconstruction error spectrum can be controlled as a function of frequency. Using this approach, the noise spectrum can be shaped according to the subjective noise perception of the human eye. Moreover, the error in coding a subband is confined to that subband.

The overall MSE incurred in an orthonormal, equal bandwidth, subband coding scheme with *K* subbands is given by

$$
D = \sum_{i=1}^{K} D_i(r_i), \tag{4.1}
$$

where $D_i(r_i)$ is the distortion rate performance of the encoder operating on the ith subband at r_i bits/sample. The average bit rate is given by

$$
R = \frac{1}{K} \sum_{i=1}^{K} r_i
$$
\n(4.2)

In equation (4.2) we have assumed that all subbands have the same number of pixels.

If this is not the case, the r_i 's should be multiplied by appropriate weighting coefficients to account for the variability in the number of pixels in any subband. The bit allocation we have used is based on the Lagrange multiplier technique in which it is assumed that all bands have the same pdf type and the distortion rate performance of the quantizers are given by

$$
D_i(r_i) = \epsilon_i^2 \sigma_i^2 2^{-2r_i}
$$
\n(4.3)

where, σ_i^2 is the variance of the *i*th subband and ϵ_i^2 is the quantization correction factor for that band. If we assume equal band distortions $D_i = D_j$, the following bit allocation is obtained [16]

$$
r_{i} = r_{ave} + \frac{1}{2} \log_{2} \left[\frac{\sigma_{i}^{2}}{\left[\prod_{i=1}^{K} \sigma_{i}^{2}\right]^{\frac{1}{K}}}\right]
$$
(4.4)

where r_{ave} is the average bit rate.

 ${r_i}$ are not restricted to be non-negative here. In practice, they are truncated to zero if they become negative. A negative bit allocation result implies that if that band is completely discarded its reconstruction error contribution is still less than the corresponding distortion for the given rate. Equation (4.4) holds only for regular tree structures of subband coding. For irregular tree structure, with N_1 bands in the first level of the tree and only band p is decomposed further into N_2 bands in the second level of the tree, the corresponding optimum bit allocation expressions are found as $[18]$

$$
r_{k1} = r_{ave} + \frac{1}{2} \log_2 \frac{\sigma_{k1}^2}{\left[\prod_{k1=1, k1 \neq p}^{N_1} \sigma_{k1}^2\right]^{\frac{1}{N_1}} \left[\prod_{k2=1}^{N_2} \sigma_{pk2}^2\right]^{\frac{1}{N_1 N_2}}}
$$
(4.5)

and,

$$
r_{pk2} = r_{ave} + \frac{1}{2} \log_2 \frac{\sigma_{pk2}^2}{\left[\prod_{k1=1, k1 \neq p}^{N_1} \sigma_{k1}^2\right]^{\frac{1}{N_1}} \left[\prod_{k2=1}^{N_2} \sigma_{pk2}^2\right]^{\frac{1}{N_1 N_2}}}
$$
(4.6)

4.2 Gain of Transform Coding

An N band orthonormal transform implies the variance preservation condition:

$$
\sigma_x^2 = \frac{1}{N} \sum_{k=0}^{N-1} \sigma_k^2
$$
\n(4.7)

where σ_x^2 is the input signal variance with zero mean and σ_k^2 are the band variances. All orthonormal, variance preserving, signal decomposition techniques can be evaluated by employing the energy compaction criterion, namely the gain of transform coding over pulse code modulation G_{TC} [16]. If we assume that all the bands and the input signal have the same *pdf* type, the distortion ratio of PCM over transform coding at the same bit rate can be easily obtained as

$$
G_{TC} = \frac{D_{PCM}}{D_{TC}} = \frac{\sigma_x^2}{\left[\prod_{k=1}^N \sigma_k^2\right]^{\frac{1}{N}}}
$$
(4.8)

This equation holds for regular binary tree structures of subband technique. For irregular tree case equation (4.8) can be modified properly. Assuming an N_1 band in the first level of the tree, and only band p is decomposed further into N_2 bands in the second level of the tree, the corresponding gain is [18]

$$
G_{TC} = \frac{\sigma_x^2}{\left(\left[\prod_{k=1, k_1 \neq p}^{N_1} \sigma_{k_1}^2 \right] \left[\prod_{k=1}^{N_2} \sigma_{pk_2}^2 \right]^{\frac{1}{N_2}} \right)^{\frac{1}{N_1}}} \tag{4.9}
$$

Chapter 5 Quantization

5.1 Introduction

The subsequent step is quantization. Quantization is the most common form of data compression and is fundamental to any digitization scheme or data compression system. Quantization causes entropy reduction and hence it is an irreversible operation. A quantizer maps a continuous variable u into a discrete \dot{u} , which takes values from a finite set $\{r_1, \dots, r_L\}$ of numbers. This mapping is generally a stair case function and the quantization rule is as follows:

- Define $\{t_k, k = 1, \ldots, L+1\}$ as a set of increasing transition or decision levels with t_1 and t_{L+1} as the minimum and maximum values, respectively, of u.
- If u lies in interval $[t_k, t_{k+1})$, then it is mapped to r_k , the kth reconstruction level.

We will consider only zero memory quantizers, which operate on one input sample at a time, and the output value depends only on that input. Each quantizer is operating at a different rate by using different quantization tables. The rate is usually measured by

$$
R_k = \log_2 N_k \tag{5.1}
$$

where R_k and N_k are the rate and the size of the table for the kth quantizer respectively. In zero memory quantizer (scalar quantization), the quantization is performed such that quantization noise is minimized, with a given number of quantization levels.

5.2 The Optimum Mean Square or Lloyd-Max Quantizer

This quantizer minimizes the mean square error for a given number of quantization levels. Let u be a real scalar random variable with a continuous probability density function $p_u(u)$. It is desired to find the decision level t_k and the reconstruction levels r_k for an *L* level quantizer such that the mean square error [17]

$$
\mathcal{E} = E[(u - \dot{u})^2] = \int_{t_1}^{t_{L+1}} (u - \dot{u})^2 p_u(u) du \qquad (5.2)
$$

is minimized. Rewriting this as

$$
\mathcal{E} = \sum_{i=1}^{L} \int_{t_i}^{t_{i+1}} (u - r_i)^2 p_u(u) du \qquad (5.3)
$$

By differentiating with respect to t_k and r_k and equating the result to zero, we get

$$
\frac{\partial \mathcal{E}}{\partial t_k} = (t_k - r_{k-1})^2 p_u(t_k) - (t_k - r_k)^2 p_u(t_k) = 0 \tag{5.4}
$$

$$
\frac{\partial \mathcal{E}}{\partial r_k} = 2 \int_{t_k}^{t_{k+1}} (u - r_k) p_u(u) du = 0 \quad 1 \le k \le L \tag{5.5}
$$

Using the fact that $t_{k-1} \leq t_k$, this gives

$$
t_k = \frac{(r_k + r_{k-1})}{2} \tag{5.6}
$$

$$
r_k = \frac{\int_{t_k}^{t_{k+1}} u p_u(u) du}{\int_{t_k}^{t_{k+1}} p_u(u) du} = E[u|u \in \mathcal{J}_k]
$$
(5.7)

where \mathcal{J}_k is the kth interval $[t_k, tk + 1)$. These results state that the optimum transition levels lie halfway between the optimum reconstruction levels, which, in turn, lie at the center of mass of the probability density in between the transition levels. Both (5.6) and (5.7) are nonlinear equations that have to be solved simultaneously. In practice, these equations can be solved by an iterative scheme such as the Newton method.

A commonly used probability densities for quantization of image-related data is the Laplacian densities, which is defined as,

$$
p_u = \frac{\alpha}{2} \exp(-\alpha |u - \mu|)
$$
 (5.8)

where μ denotes the mean of u . The variance is given by

$$
\sigma^2 = \frac{2}{\alpha} \tag{5.9}
$$

5.2.1 Properties of the Optimum Mean Square Quantizer

This quantizer has several interesting properties.

• The quantizer output is an unbiased estimate of the input, that is,

$$
E[\dot{u}] = E[u] \tag{5.10}
$$

• The quantization error is orthogonal to the quantizer output, that is,

$$
E[(u - \dot{u})\dot{u}] = 0. \t(5.11)
$$

• The variance of the quantizer output is reduced by the factor $1 - f(B)$, where *f(B)* denotes the mean square distortion of the B-bit quantizer for unity variance inputs, that is,

$$
\sigma_u^2 = [1 - f(B)]\sigma_u^2 \tag{5.12}
$$

• It is sufficient to design mean square quantizers for zero mean unity variance distributions.

5.3 Vector Quantization

When Shannon [20] proved that signals from an information source could be coded at a bit rate no greater than the entropy of the source, he showed that this would be
achieved, not by coding individual samples but by collecting samples into groups and encoding the groups. This is the basis of vector quantization.

A vector quantization scheme involves an encoder, a decoder and a codebook. The codebook is a lookup table with a k-bit address and 2^k entries. Each entry in the table is a vector of samples. Both the encoder and the decoder have copies of the codebook. The encoder looks at each incoming vector of samples \bar{x} and selects the code-word $\bar{y} = \bar{C}(i)$ which is the best match to \bar{x} . Formally, let $d(\bar{x}, \bar{y})$ be the distortion measure. The Euclidean distance $(\bar{x} - \bar{y})^T(\bar{x} - \bar{y})$ is frequently used as the measure but any positive definite quadratic form in $(\bar{x} - \bar{y})$ will serve. The encoder transmits that vector index *i* for which $d(\bar{x}, C(i))$ is minimum. The decoder receives *i* and presents $C(i)$ as its output. Clearly the success of this technique depends on having a well chosen codebook. This codebook is, in general, not complete: i.e., the number of vectors in the codebook is finite while the number of possible vectors is usually, for all practical purposes, infinite. The distortion we must live with is that which arises from assigning an incoming vector to a codebook entry that does not quite match. The quantization process is reduced to a simple search and comparison procedure. The major difficulty in implementing vector quantizers is the time required to search a very large codebook. There are several algorithms for designing a codebook. We employed LBG algorithm which is summarized in the next subsection.

Lloyd's Algorithm

A procedure for designing codebooks was developed by Lloyd [21]. This algorithm was designed as a clustering technique for use in pattern recognition and related fields. It was extended to vector quantization by Linde, Buzo, and Gray [22] and therefore called LBG algorithm. The algorithm consists of an iterative technique for refining an initial codebook. The algorithm is as follows:

- 1. Encode a selection of test data with the current codebook and measure the average distortion. If the distortion is small enough the algorithm terminates.
- 2. For each address *i* in the codebook, find the centroid of all the input vectors which were mapped into *i* and make this centriod the new $\overline{C}(i)$. Go to step 1.

The test data used in this process are refered to as the training set. Each word in the codebook is used to represent a cluster of possible input vectors from this set.

Chapter 6

Source Entropy and Huffman Coding

6.1 Entropy

The entropy of a source with *L* symbols is defined as the average information generated by the source, i.e., Entropy

$$
H = -\sum_{k=1}^{L} p_k \log_2 p_k \quad \text{bits/symbol} \tag{6.1}
$$

where p_k is the probability of having k th symbol. For a digital image considered as a source of independent pixels, its entropy can be estimated from its histogram. For the given *L* levels, the entropy of a source is maximized for uniform distributions, i.e., $p_k = 1/L, k = 1, \cdots, L$. In that case

$$
\max H = -\sum_{k=1}^{L} \frac{1}{L} \log_2 \frac{1}{L} = \log_2 L \quad \text{bits}
$$
 (6.2)

The entropy of a source gives the lower bound on the number of bits required to encode its output. Obviously, this definition assumes a stationary source.

6.2 Entropy Coding: The Huffman Coding Algorithm

The quantizer output is generally coded by a fixed-length binary code word having *B* bits. If the quantized pixels are not uniformaly distributed, then their entropy will be

less than *B,* and there exists a code that uses less than *B* bits per pixel. In entropy, the goal is to encode a block of *M* pixels containing *MB* bits with probabilities p_i , $i = 0, 1, \dots, L - 1, L = 2^{MB}$, by $(-\log_2 p_i)$ bits, so that the average bit rate is

$$
\sum_{i} p_i(-\log_2 p_i) = H \tag{6.3}
$$

This gives a variabe-length code for each block, where highly probable blocks (or symbols) are represented by small length codes, and vice versa. If $(-\log_2 p_i)$ is not an integer, the achieved rate exceeds *H* but approaches it asymtotically with increasing block size. For a given block size, a technique called *Huffman coding* is the most commonly used fixed to variable length encoding method [17].

6.2.1 The Huffman Coding Algorithm

- 1. Arrange the symbol probabilities *pi* in a decreasing order and consider them as leaf nodes of a tree.
- 2. While there is more than one node:
	- Merge the two nodes with smallest probability to form a new node whose probability is the sum of the two merged nodes.
	- Arbitrarily assign 1 and 0 to each pair of branches merging into a node.
- 3. Read sequentially from the root node to the leaf node where the symbol is located.

The preceding algorithm gives the *Huffman code book* for any given set of probabilities. Coding and decoding is done simply by looking up values in a table.

Chapter 7 Experimental Studies

Since we are emphasizing the low complexity of the Binomial filter banks introduced in[10], it would be attractive to apply these filters in complete coder structures characterized by low overall complexity. Both still image coders and video coders are implemented and tested.

7.1 Subband Coding of Still Images

In a still image coder, the original image constitutes the input to the analysis filter bank. The subband signals are represented in a bit efficient manner using optimum bit allocation expressions given in equations (4.5) and (4.6). A problem with these two equations is that they may give negative values. Moreover, the r_{kj} may not match any of the existing quantizers bit rate. Actually, we need 2^{rk_j} to be an integer equales to the number of levels in one of the used quantizers.

To solve these problems, subbands with negative r_{kj} are truncated to zero and dropped from our future calculations. Then the average bit rate r_{ave} is recalculated for the remaining bands. After that, a quickly converging iterative algorithm has been developed to find the integer number of levels corresponding to each of the remaining bands.

• First, we start with the band which has the highest frequency. The bit rate corresponding to this band is calculted using (4.5), (4.6) and the new average bit rate r_{ave} recalculated previously.

- Then, choose the quantizer with the number of levels equals or less than $2^{r_{k_j}}$.
- Excessive bit rate resulting from this integer truncation is used to reevaluate average bit rate r_{ave} for the remaining bands and this band is dropped from our next calculations.
- This new average bit rate r_{ave} is used in equations (4.5) and (4.6) to calculate the number of levels for the next lower band by repeating the same preceding procedure.

We should notice that by assigning the excessive bits, resulting from truncation, to lower bands, we are putting more emphasis on low frequencies. This feature enhances the image quality because of the mechanism of the human visual system which is more sensitive to low frequencies. For different average bit rates and the "Lady" test image, this adaptive bit allocation results are shown in Table (7.1). These numbers reflect the iterative adjustment to remove negative and noninteger values.

Next step is to quantize each band using the corresponding number of levels found using the Bit Allocation Algorithm. The bands are quantized using normalized Laplacian quantizers [17] with the following numbers of levels: $3, 5, 7, \cdots, 29, 31, 33, 35, 64, 128$. In this thesis, we have considered two methods of encoding the lowest frequency band; differential quantization such as DPCM and PCM quantization. In our experiments, we observed that the pixel to pixel correlation of upper band signals is very low. Consequently for these bands, PCM quantizer schemes are preferred to differential coding methods. We have therefore applied differential coding to the lowest band of the input signal in our experiments. In the DPCM encoder used here, a linear predictor constructs its predicted pixel as the weighted summation of previous pixels. Thus:

$$
\bar{u}(m,n) = \frac{1}{4} [u(m-1,n) + u(m,n-1) + u(m-1,n-1) + u(m-1,n+1)] \quad (7.1)
$$

The prediction error is then quantized by a symmetric non-uniform quantizer.

After quantization, the bands are entropy coded using Huffman coding scheme. To reduce complexity, only first order Huffman coding is used to code each quantized band.

Now, we can transmit the coded image over the channel. Here, we assume error free channel. It is suitable to use progressive transmission scheme for subband coding. The main objective of progressive transmission is to allow the receiver to recognize a picture as quickly as possible at minimum cost, by sending a low resolution level picture first. Further details of the picture are obtained by sequentially receiving the remaining encoded bands. At the receiver, the bands are decoded and reconstructed using the synthesizing filter bank.

A similar 8×8 two-dimensional DCT based coder is also developed for comparison purposes.

The peak-to-peak signal to noise ratio is used as the objective performance criterion and defined as

$$
SNR_{pp}(dB) = 10 \log_{10} \left[\frac{255^2}{mse} \right] \tag{7.2}
$$

where mse is the mean square coding error.

7.2 Subband Video Coding with Motion Based Adaptive Vector Quantization

In our low complexity video coder, we employ block matching for motion compensation using Brute-force method. The block size is set to 8×8 and the maximum displacement, both horizontally and vertically, is set to ± 6 pixels. The motion compensated frame difference (MCFD) signal is encoded with subband coding. The MCFD signal is split into four subbands, namely LL - HL - LH - HH bands, using 4-tap separable Binomial PR-QMF filters. The subbands are vector quantized adaptively after discarding the highest frequency band HIT.

Each 4×4 block in the subbands corresponds to a motion vector of a block with size 8×8 in the original resolution. From experiments, we have found that there is a relation between motion vector magnitude and prediction error variance of the corresponding block. In general, large magnitude motion vectors represent high variance blocks in MCFD while blocks with no motion have small variances. Using this relation, vector quantizers (codebook) were classified depending on the magnitude of block motion. The magnitude of block motion \hat{m} is defined as follow:

$$
\hat{m} = \max(|i|, |j|) \tag{7.3}
$$

where *i* and *j* are the horizontal and vertical motion respectively. We classified motion into three groups as follow:

- Group 1: $\hat{m} = 1$ or 2.
- Group 2: $\hat{m} = 3$ or 4.
- Group 3: $\hat{m} = 5$ or 6.

Using blocks corresponding to these groups for each of the three bands, codebooks were generated using LBG algorithm. As a result of this, we got 9 codebooks each of 512 length. This technique were compared to the work done by Kadur in his thesis [23][19]. Kadur used local variances of each block in subbands to classify them. This means that we need to transmit extra bits for local variances for each 4×4 block. The total bit rate for this technique can be written as

$$
B = B_M + B_{SB} + B_G \tag{7.4}
$$

where

- B_M = average bits/pixel for motion information.
- B_{SB} = average bits/pixel for subband signal.

• B_G = average bits/pixel for the transmission of local variances.

Using our MBAVQ scheme, *BG* is not needed. This means a significant reduction in bit rate (40%) in general.

Figure 7.1: SNR versus entropy for 7 subband decomposition using Laplacian quantizers for all the subbands

Figure 7.2: SNR versus entropy for 10 subband decomposition using Laplacian quantizers for all the subbands

Figure 7.3: SNR versus entropy for 7 subband decomposition using DPCM for the low frequency band and Laplacian quantizers for the remaining subbands

Figure 7.4: SNR versus entropy for 7 subband decomposition using different quantization schemes for the lowest frequency subband

Figure 7.5: SNR versus entropy for DCT, 7 and 10 subband decomposition using Laplacian quantization scheme for the lowest frequency subband

Figure 7.6: SNR versus entropy for DCT and 10 subband decomposition using DPCM quantization scheme for the lowest frequency subband

Figure 7.7: Variance vs Motion for MCFD frames 10, 24 and 30 of CINDY

Figure 7.8: SNR vs frame index of MBAVQ and VBAVQ for CINDY

Figure 7.9: Entropy vs frame index of MBAVQ and VBAVQ for CINDY

Figure 7.10: Energy Compaction Gain, G_{TC} vs frame index for (8×8) DCT and 4-tap Binomial QMF subband structure for MCFD of CINDY

Figure 7.11: SNR vs frame index of MBAVQ and VBAVQ for MONO

Figure 7.12: Entropy vs frame index of MBAVQ and VBAVQ for MONO

Figure 7.13: Motion vs frame index of MBAVQ and VBAVQ for MONO

Figure 7.14: SNR vs frame index of MBAVQ and VBAVQ for DUO

Figure 7.15: Entropy vs frame index of MBAVQ and VBAVQ for DUO

Figure 7.16: Motion vs frame index of MBAVQ and VBAVQ for DUO

	number of levels in the quantizer						
bit rate \parallel band 1 band 2 band 3 band 4 band 5 band 6 band 7							
$\parallel 0.75$	64						

Table 7.1: Results from bit allocation algorithm

Chapter 8 Discussions and Future Research

Subband coding of digital images using Binomial PR-QMF has been presented for still images. Simulation results has shown that the performance of subband coding using a low complexity coder is practically the same as the performance of the industry standard (8×8) DCT based codecs.

For video signals an efficient Motion Based Adaptive Vector Quantization (MBAVQ) subband coding method has been introduced. This new approach is compared with the Variance Based Adaptive Vector Quantization of subband video [19]. The new technique resulted in a reduction in bit rate by nearly (40%) due to the drop of the extra bits used for local variances. Moreover, this method gave superior SNR results especially for high motion frames. This clearly demonstrates that subband coding with MBAVQ should be considered as an atractive and powerful method for video coding. The modelling of MCFD signal and its relation with motion compensation techniques are open problems for future research.

APPENDIX A

```
\mathbf{C}SOURCE CODE FOR THE PROPOSED ADAPTIVE SUBBAND VIDEO CODING WI
\mathbf{C}MOTION COMPENSATION using MBAVQ
\mathbf CnxRow size of the picture
\mathbf{C}ny
             Column size of the picture
C
     framel Previous Frame
     frame2 Current Frame
C.
             Search frame from the previous frame
     pics
C.
             Prediction of the current frame with motion compensati
\mathbf Crecon
\mathbf{C}ibs
             Block size (8 is used)
             Assumed maximum displacement (Max of 8)
\mathbf Cip
C.
     frm2msk ibs*ibs size mask of the current frame to be
               motion compensated
\mathbf Cfrmlmsk ibs*ibs size mask of the previous frame in the
\mathbf Csame geometrical position (used for motion detection
C
             Motion Compensated Frame Difference Signal
\mathbf{C}errl
\mathbf{C}Search Region (ibs+ip)*(ibs+ip)searg
     mask
             Same as frm2msk
\mathbf Cparameter(nx=400,ny=512)
       character*80 input file
        common /ina/ input file
        common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
     s.
                 , 1tap3, mtap3, 1tap4, mtap4
        integer motionv(50,64), ifld
        common /motionv/ motionv
        common /ifld/ ifld
       real blv1(512,16),blv2(512,16),blv3(512,16)
       real b2v1(512,16), b2v2(512,16), b2v3(512,16)
       real b3v1(512,16), b3v2(512,16), b3v3(512,16)
       common /vqcodebook/ blv1, blv2, blv3, b2v1, b2v2, b2v3, b3v1, b3v2
     £.
                            b3v3
       integer movector(50, 64)common /bitrat/ entropy2, entropy<br>common /gtotal/ gtotal
C.
         real framel(nx, ny), frame2(nx, ny), pics(416, 528),
       recon2(nx, ny), frm1msk(8,8), frm2msk(8,8), er1(nx, ny),
    ++errtemp(nx, ny)integer ifrm1(nx, ny), ifrm2(nx, ny)
       character*1 pim(nx,ny)
```
 $character*1$ pim1(nx, ny)

common $/AA/$ searg $(24, 24)$, mask $(8, 8)$

```
\mathbf{C}Codebooks are read
C****************************
```

```
\mathbf{C}
```
C.

call readf

 $write(*,*)$ $(coff1(i), i=ltap1, mtap1)$ $\mathbf C$

```
call readvq
```

```
c mfld: final field to be read
c ifld: starting field number
     mfld = 80ifld = 1call read frm(ifld, pim)
\mathbf{C}Frame One is read into framel array
do 10 i=1, nxdo 10 j=1, nyifrm1(i,j)=ichar(pim(i,j))if(ifrm1(i,j).lt.0) ifrm1(i,j)=256+ifrm1(i,j)frame1(i,j)=float(ifrm1(i,j))10<sup>°</sup>continue
   call write_in_frm(ifld,frame1)
   call write out frm(ifld, framel)
     write(35,*) 'Original Image'
The loop to process mfld number of frames begins here
\mathbf{C}6000
     ifld = ifld+1write(*,*) 'Frame Number = ', ifld
       write(35,*) 'Frame Number =', ifld
     call read frm(ifld, pim1)
Current frame is read into frame2
\mathbf{C}
```

```
do 11 i=1, nxdo 11 j=1, nyifrm2(i,j)=ichar(pim1(i,j))if (if rm2(i, j).lt.0) if rm2(i, j)=256+ifrm2(i, j)frame2(i, j) = float(ifrm2(i, j))11continue
    call write in frm(ifld, frame2)
      if(ifld.eq.11) then
\mathbf Ccall write 1frm(frame2,'frame11')
\mathbf{C}\mathbf{C}endif
Auto-Correlation, Mean, Variance are calculated in
\mathbf{C}the subroutine autocor
\mathbf{C}c ip: displacement
      ip=6the mask block size
\mathbf Cibs:
      ibs = 8imthd: Enter 1 for Brute-force method and 2 for Orthogonal src
\mathbf{C}inthd=1imdetect: Enter 1 if motion-detection is required'
\mathbf Cimdetect=1
Search Array pics is Initialized
do 100 i=1, nx+2*ipdo 100 j=1, ny+2*ippics(i, j) = 0.0100 continue
Search Array is generated from the previous frame.
\mathbf{C}\mathbf{C}Borders are filled with first (or last) ip
  rows (or clums) of the previous frame
\mathbf{C}do 110 i=1,nx
     do 110 j=1, nypics(i+ip, j+ip) = frame1(i, j)110 continue
      do 111 i=1, ip
       do 111 j=1, nypics(i, j) = frame1(i, j)
```

```
pics(i+nx+ip, j)=framel(i+nx-ip, j)111
      continue
      do 112 i=1, nxdo 112 j=1,ippics(i,j)=framel(i,j)pics(i, j+ny+ip) = frame1(i, j+ny-ip)112
      continue
Prediction of the Current frame is Initialized
do 240 i4=1, nxdo 240 j4=1, nyrecon2(i4, j4)=0.0240
        continue
The current frame is devided into 8*8 blocks and
\mathbf{C}\mathbf{C}motion compensated. moount keeps count of number
\mathbf{C}of moving blocks.
mcount=0do 200 i=1, nx/ibsdo 200 j=1, ny/ibsiact=(i-1)*ibs+1\texttt{jact} = (\texttt{j}-1) * \texttt{ibs} + 1if (imdetect .eq. 1) then
        do 401 k=1, ibs
          do' 401 l=1, ibs
            frm1msk(k,1)=frame1(iact-1+k,jact-1+1)\texttt{frm2msk}(k,l)=\texttt{frame2}(iact-1+k,jact-1+l)401
        continue
First the motion is detected
\mathbf{C}call motiondetect(frm1msk, frm2msk, ibs, indx)
             if (indx .eq. 1) then
        mcount=mcount+1do 410 i1=1, ibs+ip*2
        do 410 j1=1, ibs+ip*2
          searg(i1, i1) = pics(i1 + iact - 1 + ip - ip, i1 + jact - 1 + ip - ip)410
            continue
         do 420 i2=1, ibs
          do 420 j2=1, ibs
```

```
mask(i2, j2)=frame2(iact-1+i2, jact-1+j2)420
           continue
if motion is detected, it is estimated and predicted
\mathbf{C}call matct(ip, ibs, imthd, n, nn, Num)
       motionv(i, j) = max(abs(n-7), abs(nn-7))mcvector(i, j)=Num
            do 430 13=1, ibs
                do 430 13=1, ibs
             recon2(iact-1+i3,jact-1+j3)=pics(iact+ip-1+(n-ip)-1+i)jact+ip-1+(nn-ip)-1+j3)430
            continue
              else
             motionv(i, j) = 0mcvector(i, j) = 0do 402 k1=1, ibs
           do 402 11=1, ibs
         recon2(iact-1+k1,iact-1+11)=frame1(iact-1+k1,iact-1+11)continue
402
           endif
    else
               do 210 i1=1, ibs+ip*2
         do 210 j1=1, ibs+ip*2
            searg(i1,j1)=pics(i1+iact-1+ip-ip,j1+jact-1+ip-ip)
210
              continue
           do' 220 i2=1, ibsdo 220 i2=1, ibs
           mask(i2, j2) = frame2(iact-1+i2, jact-1+j2)220
           continue
    call matct(ip, ibs, imthd, n, nn, Num)
           do 230 i3=1, ibs
                do 230 j3=1, ibs
             recon2(iact-1+i3,jact-1+j3)=pics(iact+ip-1+(n-ip)-1+i)jact+ip-1+(nn-ip)-1+j3)+230
           continue
    endif
200 continue
```

```
MCFD signal is generated using
\mathbf{C}err=0.0do 300 i=1, nx
        do 300 j=1, nyerr = (err + abs(frame2(i,j) - recon2(i,j)))\mathbf C\text{irecon2}(i,j) = \text{recon2}(i,j)err1(i,j) = frame2(i,j) - recon2(i,j)300
       continue
       if(iflde.g.3) then\mathbf Ccall write 1frm(err1,'diff3')
\mathbf Cendif
\mathbf C\mathbf Ctest
\mathbf C\mathbf Cif(ifld.eq.3) then
         call writeimgs(err1,400,512,'diff3.img')
\mathbf Cendif
\overline{c}C
       if(ifld.eq.10) then
\mathbf C\overline{c}call writeimgs (err1, 400, 512, 'diff10. img')
\mathbf cendif
C.
    print *,'the value of err=',err
       write(35,*) 'the value of err=', err
    if (imdetect .eq. 1) then
      print *, 'Number of blocks motion detected = ', mcount
      write(35,*)'Number of blocks motion detected =', mcount
    endif
    do 350 i = 1, nxdo 350 j = 1, nyerrtemp(i, j) = err1(i, j)350
       continue
            \pm\mathbf Ccall bitrates (mcvector, entropy, 50, 64)
the coding of the MCFD signal is carried out here
\mathbf{C}
```

```
\mathbf{C}
```

```
\mathbf{C}if(ifld.eq.10) then
       call writeint1(motionv, 50, 64, 'motion. 10')
\mathbf C\mathbf{C}call writeint(err1,400,512,'diff.10')
\overline{c}endif
\mathbf{C}if(ifld.eq.24) then\mathbf ccall writeint1(motionv, 50, 64, 'motion. 24')
\mathbf{C}call writeint(err1,400,512,'diff.24')
      endif
\mathbf{C}C.
\mathbf{C}if(ifld.eq.30) then
\mathbf{C}call writeint1(motionv, 50, 64, 'motion. 30')
      call writeint(err1,400,512,'diff.30')
\mathbf{C}endif
\mathbf{C}C.
```
 $\mathbf C$

```
call qmf(err1)print *, 'Error Signal after the vector quantization'
            write(35,*)'Error Signal after the vector quantization'
        vecmean = 0.0\text{vecvar} = 0.0do 351 i = 1, nxdo 351 j = 1, nyvecmean = vecmean + (errtemp(i,j) - err1(i,j))
        \text{vecvar} = \text{vecvar} + (\text{errtemp}(i, j) - \text{err1}(i, j)) **2
 351
        continue
    \text{vecmean} = \text{vecmean} / (\text{nx*ny})\text{vecvar} = \text{vecvar}/(\text{nx*ny})\text{vecvar} = \text{vecvar} - \text{vecmean} * 2print *, 'Variance of error signal before quant minus after
quant=\prime,
    +vecvar
        write(35,*) 'Variance of error signal before quant minus af
quant
     += ', vecvar
Quantized MCFD signal is added to the motion compensated
\mathbf{C}\mathbf{C}prediction of the current frame and put into framel and
    this becomes the previous frame for the next current frame
\mathbf{C}error = 0.0errqnt1 = 0.0errmean = 0.0do 1000 i=1, nx
```

```
do 1000 i=1.nyframe1(i,j) = recon2(i,j) + err1(i,j)errqnt1 = errqnt1 + (frame2(i,j)-frame1(i,j))**2
```

```
1000
       continue
    call write out frm(ifld, frame1)
       xmerserrqrt1/(nx*ny)write(35,*) 'MEAN SQUARE ERROR ', xmers
       write(*,*) 'MEAN SQUARE EROR', xmers
       snr=10.0*alog10(255**2/xmers)write(35,*) 'S N R = ', snr
       write(36,*) 'S N R = \prime, snr
       fbitrate = entropy/(64.0) + entropy2/(400*512)write(*,*) entropy, entropy2, fbitrate, 'final'
       write(38, *) ifld, fbitrate
       write(45,*) ifld, fbitrate, snr, gtotal
       write(*, *) ifld, fbitrate, snr, gtotal
\mathbf{C}if all the frames are not processed go back
if (ifld .lt. mfld) goto 6000
    stop
    end
subroutine matct
\mathbf{C}subroutine match(ip, ibs, imthd, n, nn, Num)common /AA/ searg(24, 24), mask(8, 8)real test(13, 13)do 50 i=1,2*ip+1
      do 50 j=1, 2*ip+1test(i,j)=0.0do 50 i=1, ibsdo 50 jj=1, ibs
           test(i,j) = abs(maxk(ii,jj) - searg(i+ii-1,j+jj-1)) + test(i50
        continue
\mathbf{C}Brute force technique
       if (imthd .eq. 1) then
        tmin=1.0e20do 100 i=1, 2*ip+1do 100 k=1,2*ip+1
            if(test(i, k) .lt. tmin) then
```
 $tmin=test(i,k)$

```
n = inn=kendif
 100
          continue
     else
     call ortho(test, ip, ibs, icent, jcent)
    n = icent
    nn = icent
     endif
\overline{c}Generates anumber between 1 & 169, The number idicates
\mathbf{C}the motion information
\mathbf{C}\overline{C}Num = (n-1) * (ip*2+1) + nnwrite(*, *) Num, n, nn
\mathbf Creturn
    end
    subroutine ortho(test, ip, ibs, icent, jcent)
Independent Orthognal Search Technique
\mathbf{C}real test(ip*2+1, ip*2+1)icent=ip+1jcent=ip+11 = ip/2. + .5istep=0
              ÷
10
    if ((test(icent,jcent) .lt. test(icent,jcent-1)) .and.
           (test(icent,jcent) .1t. test(icent,jcent+1))) then
    +icent=icent
         jcent=jcent
    else if ((test(icent,jcent-1) .lt. test(icent,jcent)) .and.
    +(test(icent, jcent-1).lt. test(icent, jcent+1))then
            icent=icent
         jcent=jcent-1
    else if ((test(icent,jcent+1) .1t. test(icent,jcent)) .and.
    +(test(icent, jcent+1) .1t. test(icent, jcent-1))then
            icent=icent
         jcent=jcent+1
```
```
istep=istep+1
    if ((test(icent,jcent) .1t. test(icent-1,jcent)) .and.
          (test(icent,jcent) .1t. test(icent+1,jcent))) then
    +icent=icent
         icent=icent
    else if ((test(icent-1,jcent) .1t. test(icent,jcent)) .and.
          (test(icent-1, jcent).lt. test(icent+1, jcent))then
            icent=icent-1
         jcent=jcent
    else if ((test(icent+1,jcent) .1t. test(icent,jcent)) .and.
          (test(icent+1,jcent) .1t. test(icent-1,jcent))) then
    +icent=icent+1
         jcent=jcent
    endif
    istep=istep+1
    if (1 \tcdot ne. 1) then
       l = (1/2.0+.5)go to 10
       endif
    return
    end
      subroutine motiondetect (frmlmsk, frm2msk, ibs, indx)
Subroutine calculates if motion is present in the
   (ibs*ibs) subblock
real frmlmsk(ibs, ibs), frm2msk(ibs, ibs)
    kount=0
    do 10 i=1, ibs
      do 10 j=1, ibs
        thrsh=abs(frm1msk(i,j)-frm2msk(i,j))if (thrsh .gt. 3.0) kount=kount+1
 10 continue
    if (kount .gt. 10) then
         indx=1else
         indx=0endif
    print *,'index', indx
    return
    end
```
endif

 \mathbf{C}

 $\mathbf C$

 \mathbf{C}

C************************** from here

C THE NEW OMF

```
\mathbf{c}_{\perp}subroutine qmf(image)
         integer raw, col
         parameter(raw=400,col=512)
         real image(raw, col)
         real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20
         common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
      \boldsymbol{\delta}, 1tap3, mtap3, 1tap4, mtap4
c these are the four subband
         real llband(raw/2,col/2),lhband(raw/2,col/2)
      & in
         , hlband (\text{raw}/2, \text{col}/2), hhband (\text{raw}/2, \text{col}/2)common /bnds/ llband, lhband
         , hlband, hhband
      & in
\mathbf{C}these are the high and low bands
         real lband(raw, col/2), hband(raw, col/2)
         common /band4/ lband, hband
c input and output images
         real inimg(raw, col), outimg(raw, col)
         common /i/ inimg
         common /o/ outimq
         real lli(raw,col/2), lhi(raw,col/2), hli(raw,col/2),
     & hhi(raw, col/2), llo(raw, col/2), lho(raw, col/2), hlo(raw, col/2)& hho(\text{raw}, \text{col}/2)common/band2/ lli, lhi, hli,
     & hhi, llo, lho, hlo,
     \mathbf{x}hho
```

```
real li(raw, col/2), lo(raw, col), hi(raw, col/2),
```

```
&.
          ho(raw,col)common /bnd1/ li, lo, hi,
      s.
          ho
                 limg(raw,col), himg(raw,col)
          real
          common /bnd0/ limg, himg
c motion vector
      integer motionv(50,64)
          common /motionv/ motionv
          common /ifld/ ifld
          common /hist/ hist
          integer hist(9,512)
          common /bitrat/ entropy2, entropy
C.
          do 12 i=1, raw
            do 12 j=1, col\text{inimg}(\text{i},\text{j}) = \text{image}(\text{i},\text{j})12continue
          call analysis
\mathbf C\mathbf Cif(ifld.eq.2) thencall writeint(llband, raw/2, col/2, 'band1.2')
\mathbf Ccall writeint(lhband, raw/2, col/2, 'band2.2')
\overline{c}\mathbf Ccall writeint(hlband, raw/2, col/2, 'band3.2')
            call writeint1(motionv, 50, 64, 'motionv. 2')
C
C.
         elseif(ifld.eq.7) then
            call writeint(llband, raw/2, col/2, 'band1.7')
Ċ.
Ċ.
            call writtent(lhband, raw/2, col/2, 'band2.7')C.
            call writeint(hlband, raw/2, col/2, 'band3.7')
            call writeint1(motionv, 50, 64, 'motionv. 7')
\overline{c}elseif(ifld.eq.10) then
Ċ
            call writeint(llband, raw/2, col/2, 'band1.10')
\mathbf CĊ
            call writeint(lhband, raw/2, col/2, 'band2.10')
            call writeint(hlband, raw/2, col/2, 'band3.10')
C
\mathbf Ccall writeint1(motionv, 50, 64, 'motionv. 10')
         elseif(ifld.eq.15) then
\mathbf C\mathbf Ccall writeint(llband, raw/2, col/2, 'band1.15')
            call writeint(lhband, raw/2, col/2, 'band2.15')
\mathbf Ccall writeint(hlband, raw/2, col/2, 'band3.15')
\overline{c}call writeint1(motionv, 50, 64, 'motionv. 15')
\mathbf C\mathbf Celseif(ifld.eq.20) then
            call writeint(llband, raw/2, col/2, 'band1.20')
C
C
            call writtent(lhband, raw/2, col/2, 'band2.20')call writeint(hlband, raw/2, col/2, 'band3.20')
C
\mathbf Ccall writeint1(motionv,50,64,'motionv.20')
\mathbf Celseif(ifld.eq.22) then
            call writeint(llband, raw/2, col/2, 'band1.22')
C
```


do 191 i=1,9
do 191 j=1,512
hist(i,j)=0

 \mathbf{c} and \mathbf{c}

```
call vec_quan(llband, 0)
         call vec quan(lhband, 1)
         call vec_quan(hlband,2)
\mathbf Ccall vec quan(hhband)
         call vbitrates(hist, entropy2, 9, 512)
         do 20 i=1, raw/2do 20 j=1, col/2
            hhband(i,j)=0.020
      continue
         call recon
         call m error(inimg, outimg, col, raw, output)
         write(*,*) 'error', output<br>write(35,*) ' mean error square of the error signal'
         write (35,*) 'between after subband and vector quantization'
         write (35,*) 'mse=', output
         call segma (llband, raw/2, col/2, x1)
         call segma (lhband, raw/2, col/2, x2)
         call segma(hlband, raw/2, col/2, x3)
         call segma (hhband, raw/2, col/2, x4)
         write(35,*) 'var=',x1,x2,x3,x4,(x1+x2+x3+x4)/4.0
         call segma(inimg, raw, col, x5)
         write(35,*) 'inimg var', x5call segma (outimg, raw, col, x6)
         write (35,*) 'outimg var=', x6
         do 10 i=1, raw
          do 10 j=1, col
          image(i, j) = outputimq(i, j)10continue
------------------
        return
```
end

 $\bar{1}$

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```
common /ina/ input file
          write (*,1)write (*,3)read (5,4) input file
 \mathbf{1}format ('
                         \prime)
 3
          format (' Enter the name of the file contains the order of
          ',/,' filtes there coefficients,
      \mathbf{\hat{x}}input Image, and output file:')
      \boldsymbol{\mathbf{\hat{x}}}\overline{4}format (a80)
          return
          end
\mathbf{C}subroutine openf
          character*80 input file
          common/ina/ input file
          open (11, file='../filters.dir/in24', status='old')
          return
          end
C_{\underline{\hspace{2mm}}\phantom{\beta\lambda}}subroutine analysis
          integer raw, col
          parameter(raw=400,col=512)
          real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20
          common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
      \delta, 1tap3, mtap3, 1tap4, mtap4
c these are the four subband
          real llband(raw/2, col/2), lhband(raw/2, col/2), hlband (\text{raw}/2, \text{col}/2), hhband (\text{raw}/2, \text{col}/2)\mathbf{\delta}common /bnds/ llband, lhband
          , hlband, hhband
      &
c these are the high and low bands
          real lband(raw, col/2), hband(raw, col/2)
```

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65
```

```
common /band4/ lband, hband
c input and output images
        real inimg(raw, col), outimg(raw, col)
        common /i/ inimg
        common /o/ outimq
        nx=raw
        ny = colcall rfilter(coff1,inimq,lband,nx,ny,ltap1,mtap1)
        call rfilter(coff2, inimq, hband, nx, ny, ltap2, mtap2)
        ny=ny/2call cfilter(coff1, lband, llband, nx, ny, ltap1, mtap1)
        call cfilter(coff2, lband, lhband, nx, ny, ltap2, mtap2)
        call cfilter(coff1, hband, hlband, nx, ny, ltap1, mtap1)
        call cfilter(coff2, hband, hhband, nx, ny, ltap2, mtap2)
        return
        end
```
C.

```
subroutine rfilter(f,al,a2,raw,col,ltap,mtap)
      integer col, raw, ltap, mtap
      real a1(raw, col), a2(raw, col/2), f(-20:20)do 20 i=1, raw
        do 20 j=2, col, 2
           a2(i,j/2)=0do 20 k=1tap, mtap
           ik=1+kif(jk.le.0) jk=col+jkif(jk.get,col)jk=jk=cola2(i,j/2) = a2(i,j/2) + a1(i,jk) * f(k)20
      continue
      return
      end
```
subroutine cfilter(f,a1,a2,raw,col,ltap,mtap)

```
integer col, raw, ltap, mtap, jk
         real al(raw, col), a2(raw/2, col), f(-20:20)do 20 i=1, col
           do 20 j=2, raw, 2a2(j/2,i)=0do 20 k=1tap, mtap
              jk=j+kif(jk.le.0) jk=raw+jkif(jk.get raw) jk=jk-rawa2(j/2,i) = a2(j/2,i) + a1(jk,i) * f(k)20
         continue
         return
         end
\mathbf Csubroutine ccfilter(f,a1,a2,raw,col,ltap,mtap)
         integer col, raw, ltap, mtap, jk
         real al(raw, col), a2(raw, col), f(-20:20)do 20 i=1, col
           do 20 j=1, raw
             a2(j,i)=0do 20 k=ltap, mtap
              jk=j+kif(jk.le.0) jk=raw+jkif(jk.gt.raw) jk=jk-raw
              a2(j,i)=a2(j,i)+a1(jk,i)*f(k)20continue
         return
         end
\mathbf{C}subroutine rcfilter(f,a1,a2,raw,col,ltap,mtap)
         integer col, raw, ltap, mtap, jk
         real a1(raw,col), a2(raw,col), f(-20:20)do 20 i=1, raw
           do 20 j=1, col
              a2(i,j)=0d^2 20 k=1tap, mtap
              jk=j+k
              if(jk.le.0) jk=col+jkif(jk.get,col) jk=jk-cola2(i,j) = a2(i,j) + a1(i,jk) * f(k)20
        continue
        return
        end
\mathbf{C}
```

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```
 $\bar{\rm t}$

```
subroutine cinter(in, out, nraw, ncol)
          integer nraw, ncol
          real in(nraw, ncol), out(nraw*2, ncol)
          do 20 j=1, ncol
            do 20 i=1, nraw
                 out(2*i-1,j)=in(i,j)out(2*i,j)=0.0out (2\star i, \overline{j}) = in(i, j)\mathbf Cout(2*i-1,i)=0.0\mathbf C20
          continue
          return
          end
\mathbf Csubroutine rinter(in, out, nraw, ncol)
          integer nraw, ncol
          real in(nraw, ncol), out(nraw, 2*ncol)
          do 20 j=1, ncol
            do 20 i=1, nraw
                 out(i, 2 * j - 1) = in(i, j)out(i, 2 * j) = 0.0\mathbf Cout(i, 2 * j) = in(i, j)\overline{\mathbf{C}}out(i, 2 * j - 1) = 0.020
         continue
         return
         end
C
         subroutine recon
         integer raw, col
         parameter(raw=400,col=512)
c input and output images
         real inimg(raw, col), outimg(raw, col)
         common /i/ inimg
         common /o/ outimq
c these are the four subband
         real llband(raw/2, col/2), lhband(raw/2, col/2), hlband (raw/2, col/2), hhband (raw/2, col/2)
      &
         common /bnds/ llband, lhband
        , hlband, hhband
      &
```

```
real lli(raw,col/2), lhi(raw,col/2), hli(raw,col/2),
& hhi(raw, col/2), llo(raw, col/2), lho(raw, col/2), hlo(raw, col/2)
```
 $\overline{1}$

```
\deltahho(\text{raw}, \text{col}/2)common/band2/ lli, lhi, hli,
         hhi, llo, lho, hlo,
     &
     \deltahho
         real li(raw,col/2), lo(raw,col), hi(raw,col/2),
     &
        ho(raw, col)common /bnd1/ li, lo, hi,
     &.
        ho
                limg(raw, col), himg(raw, col)
         real
         common /bnd0/ limg, himq
         real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20
         common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
     \mathbf{\hat{x}}, 1tap3, mtap3, 1tap4, mtap4
\mathbf{C}nraw=raw
        ncol=col
        call cinter(llband, lli, raw/2, col/2)
        call ccfilter(coff4, lli, llo, raw, col/2, ltap4, mtap4)
        call cinter(lhband, lhi, raw/2, col/2)
        call ccfilter(coff3, lhi, lho, raw, col/2, ltap3, mtap3)
        call cinter(hlband, hli, raw/2, col/2)
        call ccfilter(coff4,hli,hlo,raw,col/2,ltap4,mtap4)
        call cinter(hhband, hhi, raw/2, col/2)
        call ccfilter(coff3,hhi,hho,raw,col/2,ltap3,mtap3)
        do 10 i=1, raw
           do 10 j=1, col/2
             li(i,j)=lio(i,j)+lho(i,j)hi(i,j) = hlo(i,j) + hho(i,j)10continue
        call rinter(li, lo, raw, col/2)call rcfilter(coff4, lo, limg, raw, col, ltap4, mtap4)
        call rinter(hi, ho, raw, col/2)
        call rcfilter(coff3, ho, himq, raw, col, ltap3, mtap3)
        do 20 i=1, raw
```
 $\mathbf C$

```
do 20 j=1, col
              outimg(i, j) = (limg(i, j) + himg(i, j))if(outimg(i,j).gt.255) outimg(i,j)=255
\mathbf C\mathbf{C}if(outimg(i,j).lt.0)outimg(i,j)=0
   20
          continue
         return
         end
\mathbf{C}subroutine segma (image, raw, col, output)
        integer raw, col
        real image(raw, col), output
\overline{c}double precision s1, s2
        s1=0.0s2=0.0do 10 i=1, raw
           do 10 j=1, col
            sl=sl+image(i,j)s2 = s2 + (image(i,j)) **210<sup>°</sup>continue
        output=s2/(raw*col)-(s1/(raw*col))**2return
        end
```
 $\mathbf C$

 $10[°]$

```
subroutine m_error(image1,image2,raw,col,output)
integer raw, col
real image1(raw, col), image2(raw, col), output
s = 0.0do 10 i=1, raw
   do 10 i=1, col
    s=s+(image1(i,j)-image2(i,j))**2continue
output=s/(raw*col)
return
```
end

```
\mathbf{C}
```

```
subroutine writeimqs(pic, nx, ny, name)
      real pic(nx, ny)character*1 \text{pim}(400*512)character*20 name
        open(1,file=name,access='direct',
    + form='unformatted', recl=nx*ny)
      do 20 i=1.nxdo 20 j=1, nyip=int(pic(i,j))+128if(ip.get.255) ip=255if(ip.lt.0) ip=0
          if(ip.get.128) ip=ip-256mm=j+(i-1)*nypim(mm) = char(ip)20
      continue
      write(1, rec=1) (pim(j), j=1, nx*ny)close (1)return
      end
```
 \mathbf{C} \mathbf{C}

```
subroutine writeimg(pic, nx, ny, name)
      real pic(nx, ny)character*1 pim(400*512)
      character*20 name
        open(1,file=name,access='direct',
    + form='unformatted', recl=nx*ny)
      do 20 i=1, nxdo 20 j=1, ny
          ip=int(pic(i,j))if(ip.get.128) ip=ip-256mm=j+(i-1)*nypim(mm) = char(ip)20
      continue
      write(1, rec=1) (pim(j), j=1, nx*ny)close (1)return
      end
```
 \mathbf{C} \mathbf{c}

```
subroutine writeint(pic, nx, ny, name)
       real pic(nx, ny)
       integer ipic(400,512)
       character*20 name
       open(1,file=name)
       do 5 i=1, nxdo 5 j=1, nyipic(i,j)=int(pic(i,j)+0.5)5
       continue
       do 10 i=1, nxwrite(1,*) (ipic(i,j), j=1, ny)
10
       continue
       close (1)return
       end
```
 \mathbf{C}

 $\mathbf c$

```
subroutine writeint1(pic, nx, ny, name)
       integer pic(nx, ny)
       character*20 name
       open(1, file=name)
       do 10 i=1, nxwrite(1,*) (pic(i,j), j=1, ny)
10
       continue
       close (1)return
       end
```

```
subroutine vec_quan(pic, nband)
c pic: picture o be coded (200X256)
c nband: the band number (0,1,2,3)real pic(200,256)
        integer motionv(50,64)
```

```
real tvector(16)common /hist/ hist
        integer hist(9,512)
          common /motionv/ motionv
       real blv1(512,16), blv2(512,16), blv3(512,16)
       real b2v1(512,16), b2v2(512,16), b2v3(512,16)
       real b3v1(512,16), b3v2(512,16), b3v3(512,16)
       common /vqcodebook/ blv1, blv2, blv3, b2v1, b2v2, b2v3, b3v1, b3v2
    \boldsymbol{\delta}b3v3nn = nband*3
       do 10 i=1,50do 10 j=1,64if(motionv(i, j).ge.5) then
            do 20 k=1,4
            do 20 l=1,4tvector(4*(k-1)+1) = \text{pic}((i-1)*4+k, (i-1)*4+1)20
            continue
            mm=nn+3else if(motionv(i,j).ge.3) thendo 21 k=1, 4do 21 \ 1=1,4tvector(4*(k-1)+1)=pic((i-1)*4+k,(j-1)*4+1)
21
            continue
            mm=nn+2else if (motionv(i,j).ge.1) then
            do 22 k=1, 4do 22 \; l=1,4tvector(4*(k-1)+1)=pic((i-1)*4+k,(j-1)*4+1)
22continue
            mm = nn + 1else
            do 33 k=1, 4do 33 l=1, 4pic((i-1)*4+k,(j-1)*4+1)=0.033
        continue
           mm=0endif
      if(mm.eq.1) thencall vquantizer (tvector, b1v1, ivecnum)
           hist(1, ivecnum) = hist(1, ivecnum) + 1else if(mm.eq.2)then
            call vquantizer (tvector, blv2, ivecnum)
           hist(2, ivecnum) = hist(2, ivecnum) + 1else if(mm.eq.3) then
```
call vquantizer (tvector, b1v3, ivecnum) hist(3, ivecnum) = hist(3, ivecnum) +1 else if(mm.eq.4)then call vquantizer (tvector, b2v3, ivecnum) $hist(4, ivecnum) = hist(4, ivecnum) + 1$ else $if(mm.eq.5)$ then call vquantizer (tvector, b2v2, ivecnum) $hist(5, ivecnum) = hist(5, ivecnum) + 1$ else if (mm.eq.6) then call vquantizer (tvector, b2v3, ivecnum) $hist(6, ivecnum) = hist(6, ivecnum) + 1$ else if(mm.eq.7)then call vquantizer (tvector, b3v1, ivecnum) hist(7, ivecnum) = hist(7, ivecnum)+1 else if(mm.eq.8)then call vquantizer (tvector, b3v2, ivecnum) hist(8, ivecnum) = hist(8, ivecnum) +1 else, if (mm.eq.9) then call vquantizer (tvector, b3v3, ivecnum) $hist(9, ivecnum) = hist(9, ivecnum) + 1$ endif if (mm.ne.0) then do $44 k=1,4$ do $44 \text{ } 1=1,4$ $pic((i-1)*4+k,(j-1)*4+1) = tvector(4*(k-1)+1)$ 44 continue endif 10 continue return end \mathbf{C} subroutine vquantizer (testv, codebook, ivecnu) c Best Matching of vector real $testv(16)$ real codebook(512,16) $rdiff = 1000000.0$ ivecnu = 0

 \mathbf{r}

```
do 110 m = 1,512adiff = 0do 120 n = 1,16\text{adiff} = \text{adiff} + (\text{testv}(n) - \text{codebook}(m, n)) **2120
             continue
             if (adiff .1t. rdiff) then
                rdiff = \text{adiff}ivecnu = m
             endif
110
        continue
        do 130 n = 1,16testv(n) = codebook(ivecnu,n)130
        continue
        return
        end
```

```
\mathbf{C}
```

```
subroutine readvq
```

```
real blv1(512,16),blv2(512,16),blv3(512,16)
        real b2v1(512,16), b2v2(512,16), b2v3(512,16)
        real b3v1(512,16), b3v2(512,16), b3v3(512,16)
        common /vqcodebook/ blv1, blv2, blv3, b2v1, b2v2, b2v3, b3v1, b3v2
     \deltab3v3open (15, file='../quantizer.dir/b1v1.12')
        do 10 i=1,512read(15,*) (b1vl(i,j), j=1,16)10continue
        close(15)open (16, file='../quantizer.dir/blv2.34')
        do 11 i=1,512read(16,*) (b1v2(i,j), j=1, 16)11continue
        close(16)open (17, file='../quantizer.dir/b1v3.56')
        do 12 i=1,512read(17,*) (b1v3(i,j), j=1, 16)12continue
        close(17)open (18, file='../quantizer.dir/b2v1.12')
        do 13 i=1,512read(18,*) (b2v1(i,j), j=1, 16)13continue
```

```
close(18)open (19, file='../quantizer.dir/b2v2.34')
         do 14 i=1,512
           read(19,*) (b2v2(i,j), j=1, 16)14
      continue
         close(19)open (20, file='../quantizer.dir/b2v3.56')
         do 15 i=1,512
           read(20,*) (b2v3(i,j), j=1, 16)15
      continue
         close(20)open (21, file='../quantizer.dir/b3v1.12')
         do 16 i=1,512
           read(21,*) (b3v1(i,j), j=1, 16)16
      continue
         close(21)open (22, file='../quantizer.dir/b3v2.34')
         do 17 i=1,512
           read(22,*) (b3v2(i,j), j=1, 16)17continue
         close(22)open (23, file='../quantizer.dir/b3v3.56')
         do 18 i=1,512read(23,*) (b3v3(i,j), j=1, 16)18
     continue
         close(23)return
     end
\mathbf{C}subroutine read frm(ifld, pic)
       character*1 pic(400, 512)open(1,file='ali.100',access='direct',form=
C.
     & 'unformatted', recl=512*512)
\mathbf C\mathbf{C}\mathbf{C}open(1,file='/images/cindy40',access='direct',form=
\mathbf C& 'unformatted', recl=512)
       open(1,file='/images/mono',access='direct',form=
     & 'unformatted', recl=512)
       open(1,file='/images/quartet',access='direct',form=
\mathbf C
```

```
'unformatted', recl=512)
\mathtt{C}\delta\mathbf{C}\mathbf{C}open(1,file='/images/duo',access='direct',form=
\mathbf{C}& 'unformatted', recl=512)
\mathbf Cicod1 = (if1d-1)*400icod2 = (ifld-1)*400 + 200do 10 i=1,200read(1, rec=icod1+i) (pic(2*i-1,j), j=1,512)read(1, rec=icod2+i)(pic(i*2, j), j=1, 512)
10
        continue
         close(1)return
         end
\mathbf Csubroutine write in frm(ifld, pic)
         real pic(400, 512)character*1 image(512*512)
         open(22, file='mono.in80', access='direct', form=
        'unformatted', recl=512*512)
     &.
     do 10 i=1,400do 10 j=1,512ip=int(pic(i,j)+.5)if(ip.get.255) ip=255if(ip.1t.0) ip=0
          if(ip.get.127) ip=ip-256image((i-1)*512+j) = char(ip)10 continue
         do 20 i=204801,262144
           image(i) = char(003)20
         continue
         write(22, rec=ifld) (image(j), j=1, 512*512)close(22)return
     end
\mathbf{C}
```

```
subroutine write out frm(ifld, pic)
   real pic(400,512)
   character*1 image(512*512)
   open(21, file='mono.out80', access='direct', form=
```

```
\mathbf{\hat{x}}'unformatted', recl=512*512)
    do 10 i=1,400do 10 i=1,512ip=int(pic(i,j)+.5)if(ip.get.255) ip=255if (ip.lt.0) ip=0if(ip.get.127) ip=ip-256image((i-1)*512+j) = char(ip)10 continue
       do 20 i=204801,262144
         image(i) = char(100)20
       continue
       write(21, rec=ifld) (image(j), j=1, 512*512)close(21)return
    end
```

```
\mathbf{C}
```

```
subroutine write 1frm(pic, name)
         real pic(400, 512)character*1 image(512*400)character*20 name
         open(22, file=name, access='direct', form=
      & 'unformatted', recl=512*400)
      do 10 i=1,400
         do 10 j=1,512
          ip=int(pic(i,j)+.5)if(ip.gt.255) ip=255
          if(ip.lt.0) ip=0if(ip.get.127) ip=ip-256image((i-1)*512+j) = char(ip)10 continue
         write (22, \text{rec}=1) (image(j), j=1, 512*400)
     close(22)return
     end
               \sim 1 ^{\circ}\mathbf{c} and \mathbf{c}
```
c this subroutine calculate the entropy of each band c and find the probability of each code

 \mathbf{v}

```
subroutine vbitrates(ic, bitrate, raw, col)
         common /gtotal/ gtotal
         integer ic(raw, col), raw, col
         real entropy(9), sum(512), pr(512)qtotal=0bitrate=0
     do 10 m=1,9
         total=0
         do 20 n=1,512
          sum(n) = ic(m, n)total = total + sum(n)20
     continue
\mathbf{C}entropy(m)=0.0do 30 n=1,512
            pr(n) = sum(n) / totalif(pr(n).gt.0) thenbr=pr(n)*xlog2(1.0/pr(n))entropy(m) = entropy(m) + brendif
 30<sup>°</sup>continue
         bitrate=bitrate+entropy(m)*total
         write(*,*) 'total = ', total
         gtotal=gtotal+total
 10 continue
         qtotal=qtotal/3
        write(*,*) 'gtotal = ',gtotal<br>write(*,*) 'ventropy = ',(entropy(i),i=1,9)
        write(*,*) 'bitrate' = ', bitrate'
         return
         end
```

```
\mathbf{C} and \mathbf{C}
```
c this subroutine calculate the entropy of each band c and find the probability of each code

subroutine bitrates(ic, entropy, raw, col)

integer ic(raw, col), raw, col real entropy, sum $(0:169)$, pr $(0:169)$

do 20 n=0,169

```
sum(n)=0.020continue
\mathbf Cdo 10 i=1, raw
         do 10 j=1, colk=ic(i,j)sum(k) = sum(k) + 110
         continue
         entropy=0.0total=real(raw*col)
         do 30 n=0,169
            pr(n) = sum(n) / totalif(pr(n).gt.0) thenbr=pr(n)*xlog2(1.0/pr(n))entropy=entropy+br
            endif
 30
         continue
         write(*,*) 'entropy = ', entropy
         write (*,*) 'pr = ', (pr(i), i=1, 169)
\mathbf Creturn
         end
\mathbf{C}function xlog2(x)
         real x
         xlog2 = alog(x)/alog(2.0)return
         end
\mathbf C\mathbf{c}
```
 $\mathbf C$ Simulation program for 10 band decomposition \mathbf{C} Name: Hosam Mutlaq \mathbf{C} **MAIN** \mathbf{C} character*80 input file common /ina/ input file character*80 inputing common /image/ inputimg common /bitrate/ rav open(50, file='10bdpcm.res', form='formatted') do 20 rav=.2, 1.3, 0.1 $\mathbf C$ rav= 1.0 input file='in24' inputimg='ladyp.img' call subband10(e1, or1, snr1) input file='in6' call subband10(e2, or2, snr2) input file='in8' call subband10(e3, or3, snr3) $write(50, 10)$ rav, snr1, snr2, snr3, e1, or1, e2, or2, e3, or3 20 continue do 30 rav=1.5,5.0,0.5 input file='in24' inputimg='ladyp.img' call subband10(e1, or1, snr1) input file='in6' call subband10(e2, or2, snr2) input file='in8' call subband10(e3, or3, snr3) write(50,10) rav, snr1, snr2, snr3, e1, or1, e2, or2, e3, or3 continue 30 format(f4.2,3f8.4,6f8.5) 10 $close(50)$

```
stop
         end
\mathbf CThis is a 10 band simulation program.
\mathbf{C}it uses dpcm quantizer for the first band
\mathbf{c}c and laplacian quanizer for the rest. hoffman coading.
c the program does bit allocation for the 10 bands.
\mathbf{C}\mathbf{C}subroutine subband10(xo1,xo2,snr)
         integer raw, col
         parameter(raw=256,col=256)
         real \dim(g(32,32), \dim(g(32,32))real p(75), p1(75)character*80 input file
         common /ina/ input file
         character*80 inputimq
         common /image/ inputimg
         common /bitrate/ rav
         common /header/ header
         character*1 header(64)
         real x(10), output, xmean(10), entrpy(10), entrpy1(4), entrpy2(3)
         real entrpy3(3)
         real xt(64), xmeant(64)real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20
         common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
     \mathbf{\hat{x}}, ltap3, mtap3, ltap4, mtap4
         integer iq(64), iqt(64)real pr(-64:64), prb(128)integer prnum(128)
         character*80 codbook(128)
         character*1 hc(77000)
\mathbf{C}real \text{ining}(256, 256)real a1(128,128), a2(128,128), a3(128,128), a4(128,128)
        real b1(64, 64), b2(64, 64), b3(64, 64), b4(64, 64)real c1(32,32), c2(32,32), c3(32,32), c4(32,32)
```

```
82
```

```
real el(128,128)
         real d1(64, 64)real qc1(32,32), qc2(32,32), qc3(32,32), qc4(32,32)
         real q b2(64, 64), q b3(64, 64), q b4(64, 64)real qa2(128,128), qa3(128,128), qa4(128,128)
         integer ic1(4,32,32), icr1(4,32,32)
         integer ic2(3,64,64), icr2(3,64,64)
         integer ic3(3,128,128), icr3(3,128,128)
         real outimg(256, 256)\mathbf C\overline{C}call initial
         call readf
         call readimage(inimg, raw, col)
         write(*,*) 'subband analysis'
         call analysis256(inimg, al, a2, a3, a4)
         call analysis128(a1,b1,b2,b3,b4)
         call analysis64(bl, cl, c2, c3, c4)
\mathbf{C}write(*,*) 'calculating the variances and means'
         call seqma(cl, 32, 32, x(1), xmean(1))call segma (c2, 32, 32, x(2), xmean(2))
         call segma(c3,32,32,x(3), xmean(3))
         call segma (c4, 32, 32, x(4), xmean(4))
         call segma(b2,64,64,x(5), xmean(5))
```

```
call segma(b3, 64, 64, x(6), xmean(6))
call segma (b4, 64, 64, x(7), xmean(7))
call ségma(a2,128,128, x(8), xmean(8))
call segma(a3,128,128,x(9), xmean(9))
call segma(a4,128,128, x(10), xmean(10))
call dpcm1(c1,dimg, 32, 32)
call segma (dimg, 32, 32, var2, rmean)
write(*, *) x(1), xmean(1)x(1) = var2xmean(1) = rmeanwrite(*, *) var2, rmean
```

```
\mathbf CC
```
 $\mathbf C$

```
c convert the 10band to 64 band for bit allocation
         do 10 i=1,4xt(i)=x(i)xmean(i) = xmean(i)10<sup>°</sup>continue
     do 15 j=1,3do 20 i = (j*4)+1, j*4+4xt(i)=x(j+4)xmean(i) = xmean(i+4)continue
 20<sup>°</sup>15
         continue
     do 11 j=1,3do 21 i=(j*16)+1,j*16+16xt(i)=x(j+7)xmean(i) = xmean(i+7)21continue
         continue
 11\overline{C}write(*,*) (x(i), i=1, 10)write(*,*) (xt(i), i=1, 64)C.
\mathbf{C}call entropy(iqt, xt, xmeant)
\mathbf Ciq(1)=iqt(1)iq(2)=iqt(2)iq(3)=iqt(3)iq(4)=iqt(4)iq(5) = iqt(5)iq(6)=iqt(9)iq(7) = iqt(13)iq(8) = iqt(17)iq(9) = iqt(33)iq(10) = iqt(49)write(*,*) (iq(i), i=1, 10)write (*, *) (iqt(i), i=1,64)
\mathbf{C}\mathbf{C}write(*, *) 'quantizing'
         nm = iq(1)call readq(nm, p, pl)if(nm.ne.iq(1)) then
```

```
write (*, *) 'error in reading quantizer data'
  stop
end if
```


endif

25 continue

 $sum = sum/4.0$ \mathbf{c} do 35 i=1.3 $if(iq(i+4).gt.0.1) then$ call bitrates(ic2, i, entrpy2, pr, 3, 64, 64) call convert(pr, prb, prnum, ns) call sort(prb, prnum, ns) call hcbook(prb, codbook, ns) call codehc(ic2, i, codbook, prnum, ns, hc, 11, 3, 64, 64) $\mathbf C$ call $decodehc(icr2, i, codbook, prnum, ns, hc, ll, 3, 64, 64)$ \ddot{c} do 104 i1=1,64 do 104 $j1=1,64$ $icr = ic2(i, i1, j1) - icr2(i, i1, j1)$ $\mathbf C$ write $(*,*)$ icr if(icr.ne.0) then write(*,*) 'ERROR2 IN CHANNEL CODING', i, i1, j1, icr stop endif 104 continue length=length+11 $entropy(i+4) = entry2(i)$ $sum = sum + entry2(i)$ endif $35₁$ continue $sum = sum/4.0$ $\mathbf{C}_{_}$ do 45 i=1,3 $if(iq(i+7).qt.0.1) then$ call bitrates(ic3, i, entrpy3, pr, 3, 128, 128) call convert(pr, prb, prnum, ns) call sort (prb, prnum, ns) call hcbook(prb, codbook, ns) call codehc(ic3, i, codbook, prnum, ns, hc, 11, 3, 128, 128) $\mathbf C$ call decodehc(icr3, i, codbook, prnum, ns, hc, 11, 3, 128, 128) C

```
do 105 i1=1,128
          do 105 j1=1,128
            icr = ic3(i, i1, j1) - icr3(i, i1, j1)\overline{c}write(*, *) icr
          if(icr.ne.0) then
            write(*,*) 'ERROR2 IN CHANNEL CODING', i, i1, j1, icr
            stop
          endif
 105
          continue
          length=length+11
           entropy(i+7) = entry3(i)sum = sum + entry3(i)endif
 45
    continue
         write(*, *) 'entropy=', sum/4.0
         xol=sum/4.0
         xo2 = real (length) / (256.0**2)\mathbf{C}write(*, *) 'DEquantizing'
         call dequantizer(cimg, 1, icr1, iq(1), rmean, var2, 4, 32, 32)
         call dpcdecoder(cimq, qc1, 32, 32)
         call m_error(qc1,c1,32,32,cerror)
         write(*, *) 'cl error', cerror
\mathbf Ccall unifdequan(qcl, 1, icrl, iq(1), xminx, step, 4, 32, 32)
\overline{c}call dequantizer(qc1,1,icr1,iq(1),xmean(1),x(1),4,32,32)
         call dequantizer(qc2, 2, icr1, iq(2), xmean(2), x(2), 4, 32, 32)
         call dequantizer(qc3, 3, icr1, iq(3), xmean(3), x(3), 4, 32, 32)
         call dequantizer (qc4, 4, icr1, iq(4), xmean(4), x(4), 4, 32, 32)call dequantizer(qb2,1,icr2,iq(5),xmean(5),x(5),3,64,64)
         call dequantizer(qb3,2,icr2,iq(6),xmean(6),x(6),3,64,64)
         call dequantizer(qb4,3,icr2,iq(7),xmean(7),x(7),3,64,64)
         call dequantizer(qa2,1,icr3,iq(8),xmean(8),x(8),3,128,128)
         call dequantizer(qa3,2,icr3,iq(9),xmean(9),x(9),3,128,128)
         call dequantizer(qa4,3,icr3,iq(10), xmean(10), x(10), 3, 128, 12
\mathbf{C}write(*, *) 'synthesis'
```

```
call synthesis64(qc1,qc2,qc3,qc4,d1)
call synthesis128(d1,qb2,qb3,qb4,e1)
```
call synthesis256(e1, qa2, qa3, qa4, outimg)

```
\mathbf{C}do 131 i=1,256
           do 131 j=1,256if(outimg(i, j).gt.255) then
            outimg(i, j) =255
            else if(outimg(i,j).lt.0) thenoutimg(i,j)=0endif
 131
         continue
         call m error(inimg, outimg, col, raw, output)
         call writeimg(outimg, 256, 256, 'out.img')
\mathbf C\mathbf Cwrite(*, *) 'error', output
         snr=10*alog10(255**2/output)write(*,*) 'SNR =', snr
\mathbf{C}return
         end
\mathbf{C}\mathbf CSUBROUTINE READ
\mathbf{C}THIS SUBROUTINE READ THE DATA OF FILTERS COEFFICEINTS
\overline{C}subroutine readf
          SUBROUTINE READF
C
         real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20
         common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
      \delta, ltap3, mtap3, ltap4, mtap4
         call openf
         write(*,*) 'reading filter coefficeints'
C.
         read(11, *) ltap1
          write(*, *)ltapl
\mathbf{C}read(11, *) mtap1
         do 10 i=ltap1, mtap1
   10read(11,*) coff1(i)\mathbf Cread(11, *) ltap2
         read(11,*)mtap2do 20 i=ltap2, mtap2
```

```
20
                 read(11,*) coff2(i)\overline{c}read(11, *)ltap3
         read(11,*)mtap3do 30 i=ltap3, mtap3
   30
                 \text{read}(11,*) \text{coff3}(i)c
         read(11, *)ltap4
         read(11, *) mtap4
         do 40 i=1tap4, mtap4
   40
                 read(11, \star) coff4(i)
\mathbf{C}close (11)RETURN
         END
\mathbf Csubroutine initial
\mathbf CThis subprogram initialize the main program
\mathbf C\mathbf Csubroutine initial
         character*80 input file, inputimq
         real rav
         common /ina/ input file
         common /image/ inputimg
         common /bitrate/ rav
         write(*, *) 'initialization'
\mathbf Cwrite (*,1)write (*, 3)read (5,4) input_file
         write(*,*) ' Enter the name of the input image'
         read (*, *) inputimg
         write(*,*) \prime Enter the bit rate'
         read(*,*) rav
 \mathbf{1}format ('
                       \left\langle \right\rangle3
         format (' Enter the name of the file contains the order of
      \delta ',/,' filtes there coefficients,')
         format( a80)
 \overline{4}return
         end
```

```
subroutine openf
```
 C_{\perp}

```
character*80 input file
common/ina/ input Fileopen (11, file='../filters.dir/'//input_file, status='old')
return
end
```
 \mathbf{C}

```
subroutine analysis256(inimg, llband, lhband, hlband, hhband)
         integer raw, col
         parameter(raw=256,col=256)
         real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20
         common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
      \boldsymbol{\delta}, 1tap3, mtap3, 1tap4, mtap4
c these are the four subband
         real llband(raw/2,col/2),lhband(raw/2,col/2)
     &
         , hlband (\text{raw}/2, \text{col}/2), hhband (\text{raw}/2, \text{col}/2)these are the high and low bands
\mathbf{C}real lband(raw, col/2), hband(raw, col/2)
c input and output images
         real inimg(raw, col)
         nx=raw
         ny = colcall rfilter(coff1,inimg,lband,nx,ny,ltap1,mtap1)
         call rfilter(coff2, inimg, hband, nx, ny, ltap2, mtap2)
         ny = ny/2call cfilter(coff1, lband, llband, nx, ny, ltap1, mtap1)
         call cfilter(coff2, lband, lhband, nx, ny, ltap2, mtap2)
```

```
call cfilter(coff1, hband, hlband, nx, ny, ltap1, mtap1)
call cfilter(coff2, hband, hhband, nx, ny, ltap2, mtap2)
return
end
```
subroutine analysis128(inimg, llband, lhband, hlband, hhband) integer raw.col parameter(raw=128,col=128) real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20 common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2 δ , ltap3, mtap3, ltap4, mtap4

c these are the four subband real $llband(raw/2, col/2)$, $lhband(raw/2, col/2)$, hlband $(\text{raw}/2, \text{col}/2)$, hhband $(\text{raw}/2, \text{col}/2)$ &

these are the high and low bands \mathbf{C} real lband(raw, col/2), hband(raw, col/2)

c input and output images real inimq(raw, col)

 \mathbf{C}

 $nx = raw$ $ny = col$ call rfilter(coff1, inimg, lband, nx, ny, ltap1, mtap1) call rfilter(coff2, inimg, hband, nx, ny, ltap2, mtap2) $ny = ny/2$ call cfilter(coff1, lband, llband, nx, ny, ltap1, mtap1) call cfilter(coff2, lband, lhband, nx, ny, ltap2, mtap2)

```
call cfilter(coff1, hband, hlband, nx, ny, ltap1, mtap1)
call cfilter(coff2, hband, hhband, nx, ny, ltap2, mtap2)
return
end
```
 \mathbf{C}

```
subroutine analysis64(inimg, llband, lhband, hlband, hhband)
         integer raw, col
         parameter (raw=64, col=64)
         real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20
         common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
     \delta, ltap3, mtap3, ltap4, mtap4
c these are the four subband
         real llband(raw/2, col/2), lhband (raw/2, col/2), hlband (\text{raw}/2, \text{col}/2), hhband (\text{raw}/2, \text{col}/2)&
c these are the high and low bands
         real lband(raw,col/2), hband(raw, col/2)
c input and output images
        real inimg(raw, col)
        nx = rawny=col
        call rfilter(coff1,inimg,lband,nx,ny,ltap1,mtap1)
        call rfilter(coff2,inimg,hband,nx,ny,ltap2,mtap2)
        ny = ny/2call cfilter(coff1, lband, llband, nx, ny, ltap1, mtap1)
        call cfilter(coff2, lband, lhband, nx, ny, ltap2, mtap2)
```

```
92
```
call cfilter(coff1, hband, hlband, nx, ny, ltap1, mtap1)

```
call cfilter(coff2, hband, hhband, nx, ny, ltap2, mtap2)
         return
         end
C.
         subroutine rfilter(f,a1,a2,raw,col,ltap,mtap)
         integer col, raw, ltap, mtap
         real al(raw, col), a2(raw, col/2), f(-20:20)do 20 i=1, raw
           do 20 j=2, col, 2
              a2(i,j/2)=0do 20 k=ltap, mtap
              ik=1+kif(jk.le.0) jk=col+jkif(jk.get,col) jk=jk=cola2(i,j/2) = a2(i,j/2) + a1(i,jk) * f(k)20<sub>o</sub>continue
         return
         end
\mathbf Csubroutine cfilter(f,a1,a2,raw,col,ltap,mtap)
         integer col, raw, ltap, mtap, jk
        real al(raw, col), a2(raw/2, col), f(-20:20)do 20 i=1, col
           do 20 j=2, raw, 2a2(j/2,i)=0do 20 k=1tap, mtap
              ik=1+kif(jk.le.0) jk=raw+jkif(jk.getraw) jk=jk-rawa2(j/2,i) = a2(j/2,i) + a1(jk,i) * f(k)20continue
        return
        end
C.
       subroutine ccfilter(f,a1,a2,raw,col,ltap,mtap)
        integer col, raw, ltap, mtap, jk
        real a1(raw,col), a2(raw,col), f(-20:20)
```

```
do 20 i=1, col
  do 20 j=1, rawa2(j,i)=0
```

```
do 20 k=1tap, mtap
                jk=j+kif(jk.le.0) jk=raw+jkif(jk.getraw) jk=jk-rawa2(j,i)=a2(j,i)+a1(jk,i)*f(k)continue
  20
         return
         end
\mathbf{C}subroutine rcfilter(f,a1,a2,raw,col,ltap,mtap)
          integer col, raw, ltap, mtap, jk
         real al(raw,col), a2(raw,col), f(-20:20)do 20 i=1, raw
            do 20 j=1, col
               a2(i,j)=0do 20 k=1tap, mtap
                \dot{\uparrow}k=\dot{\uparrow}+k
                if(jk.le.0) jk=col+jkif(jk.get,col)jk=jk=cola2(i,j) = a2(i,j) + a1(i,jk) * f(k)20
         continue
         return
         end
\mathbf{C}subroutine cinter(in, out, nraw, ncol)
         integer nraw, ncol
         real in(nraw, ncol), out(nraw*2, ncol)
         do 20 j=1, ncol
            do 20 i=1, nraw
                 out(2*i-1,j)=in(i,j)out (2 * i, j) = 0.0out(2*i,j)=in(i,j)\mathbf C\mathbf Cout(2*i-1,j)=0.020continue
         return
         end
\mathbf{C}subroutine rinter(in, out, nraw, ncol)
         integer nraw, ncol
         real in(nraw, ncol), out(nraw, 2*ncol)
         do 20 j=1, ncol
            do 20 i=1, nraw
                 out(i, 2 * j - 1) = in(i, j)out(i, 2*1) = 0.0\mathbf Cout(i, 2 * j) = in(i, j)out(i, 2+j-1)=0.0\overline{c}continue
 20
```
```
end
        subroutine synthesis256(llband, lhband, hlband, hhband, outimg)
        integer raw, col
        parameter(raw=256,col=256)
c input and output images
        real outimg(raw, col)
```

```
c these are the four subband
          real llband(raw/2,col/2), lhband(raw/2,col/2)
          , hlband (\text{raw}/2, \text{col}/2), hhband (\text{raw}/2, \text{col}/2)&
```
return

 \mathbf{C}

```
real lli(raw,col/2), chi(raw,col/2), hli(raw,col/2),
   hhi(raw,col/2),llo(raw,col/2),lho(raw,col/2),hlo(raw,col/2)
&
& hho(raw, col/2)real li(raw, col/2), lo(raw, col), hi(raw, col/2),
\& ho(raw, col)
   real
          \lim q(raw, col), himq(raw, col)
   real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20
   common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
\mathbf{\hat{x}}, ltap3, mtap3, ltap4, mtap4
```

```
\overline{c}
```

```
nraw=raw
ncol=col
call cinter(llband, lli, raw/2, col/2)
call ccfilter(coff4, lli, llo, raw, col/2, ltap4, mtap4)
call cinter(lhband, lhi, raw/2, col/2)
call ccfilter(coff3, lhi, lho, raw, col/2, ltap3, mtap3)
call cinter(hlband, hli, raw/2, col/2)
call ccfilter(coff4,hli,hlo,raw,col/2,ltap4,mtap4)
call cinter(hhband, hhi, raw/2, col/2)
call ccfilter(coff3,hhi,hho,raw,col/2,ltap3,mtap3)
```
 \overline{C}

```
do 10 i=1, raw
            do 10 j=1, col/2
              li(i,j)=llo(i,j)+lho(i,j)hi(i, j) = hlo(i, j) + hho(i, j)10
         continue
         call rinter(li, lo, raw, col/2)
         call rcfilter(coff4,lo,limg,raw,col,ltap4,mtap4)
         call rinter(hi, ho, raw, col/2)
         call rcfilter(coff3, ho, himg, raw, col, ltap3, mtap3)
         do 20 i=1, raw
           do 20 j=1, col
              outimg(i,j)=1*(ling(i,j)+himg(i,j))if(outimg(i,j).gt.255)outimg(i,j)=255
\mathbf C\mathbf{C}if(outimg(i,j).lt.0)outimg(i,j)=0
   20
          continue
         return
         end
\mathbf{C}subroutine synthesis128(11band, 1hband, h1band, hhband, outimg)
         integer raw, col
         parameter (raw=128, col=128)
c input and output images
         real outimg(raw.col)
c these are the four subband
         real llband(raw/2, col/2), lhband(raw/2, col/2)\delta, hlband (\text{raw}/2, \text{col}/2), hhband (\text{raw}/2, \text{col}/2)real lli(raw,col/2), lhi(raw,col/2), hli(raw,col/2),
     & hhi(raw, col/2), llo(raw, col/2), lho(raw, col/2), hlo(raw, col/2)hho(\text{raw}, \text{col}/2)&
```

```
real li(raw, col/2), lo(raw, col), hi(raw, col/2),
   ho(raw, col)\delta
```

```
limg(raw,col), himg(raw,col)
real
```
real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20

```
common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
       , ltap3, mtap3, ltap4, mtap4
```
 \mathbf{C}

 δ

```
nraw=raw
         ncol=col
         call cinter(llband, lli, raw/2, col/2)
         call ccfilter(coff4, lli, llo, raw, col/2, ltap4, mtap4)
         call cinter(lhband, lhi, raw/2, col/2)
         call ccfilter(coff3, lhi, lho, raw, col/2, ltap3, mtap3)
         call cinter(hlband, hli, raw/2, col/2)call ccfilter(coff4,hli,hlo,raw,col/2,ltap4,mtap4)
         call cinter(hhband, hhi, raw/2, col/2)
         call ccfilter(coff3, hhi, hho, raw, col/2, ltap3, mtap3)
\mathbf Cdo 10 i=1, raw
           do 10 j=1, col/2
             li(i,j)=lio(i,j)+lho(i,j)hi(i,j) = hlo(i,j) + hho(i,j)continue
  10
         call rinter(li, lo, raw, col/2)call rcfilter(coff4, lo, limg, raw, col, ltap4, mtap4)
         call rinter(hi, ho, raw, col/2)call rcfilter(coff3, ho, himg, raw, col, ltap3, mtap3)
         do 20 i=1, raw
           do 20 j=1, col
             outimg(i, j) = 1*(ling(i, j) + himg(i, j))if(outimg(i,j).gt.255)outimg(i,j)=255
\mathbf{C}if(outimg(i,j).lt.0)outimg(i,j)=0
\mathbf Ccontinue
   20
         return
         end
C.
```

```
subroutine synthesis64(llband, lhband, hlband, hhband, outimg)
         integer raw, col
         parameter(raw=64, col=64)
c input and output images
         real outimg(raw, col)
c these are the four subband
         real 11band(raw/2,col/2), 1hband(raw/2,col/2)
      \delta, hlband (\text{raw}/2, \text{col}/2), hhband (\text{raw}/2, \text{col}/2)real lli(raw,col/2), lhi(raw,col/2), hli(raw,col/2),
      & hhi(raw, col/2), llo(raw, col/2), lho(raw, col/2), hlo(raw, col/2)& hho(raw, col/2)real li(raw,col/2), lo(raw,col), hi(raw,col/2),
      \& ho(raw, col)
         real
                limg(raw,col), himg(raw,col)
         real coff1(-20:20), coff2(-20:20), coff3(-20:20), coff4(-20:20
         common /a/ coff1, coff2, coff3, coff4, ltap1, mtap1, ltap2, mtap2
      \mathbf{\hat{x}}, ltap3, mtap3, ltap4, mtap4
\mathbf{C}nraw=raw
         ncol=col
         call cinter(llband, lli, raw/2, col/2)
         call ccfilter(coff4, lli, llo, raw, col/2, ltap4, mtap4)
         call cinter(lhband, lhi, raw/2, col/2)
         call ccfilter(coff3, lhi, lho, raw, col/2, ltap3, mtap3)
         call cinter(hlband, hli, raw/2, col/2)
         call ccfilter(coff4,hli,hlo,raw,col/2,ltap4,mtap4)
         call cinter(hhband, hhi, raw/2, col/2)
         call ccfilter(coff3,hhi,hho,raw,col/2,ltap3,mtap3)
\overline{C}
```

```
do 10 i=1, raw
```

```
do 10 j=1, col/2
              li(i,j)=lio(i,j)+lho(i,j)hi(i,j) = hlo(i,j) + hho(i,j)10continue
         call rinter(li, lo, raw, col/2)
         call rcfilter(coff4, lo, limq, raw, col, ltap4, mtap4)
         call rinter(hi, ho, raw, col/2)
         call rcfilter(coff3, ho, himq, raw, col, ltap3, mtap3)
         do 20 i=1, raw
           do 20 j=1, col
             outimg(i, j) = 1 * (ling(i, j) + himg(i, j))if(outimg(i,j).gt.255)outimg(i,j)=255
\mathbf C\overline{c}if(outimg(i,j).lt.0) outimg(i,j)=0
   20
          continue
         return
         end
```
 \mathbf{C}

```
subroutine readimage(pic, nx, ny)
       real pic(nx, ny)character*1 pim(65600), header(64)
       integer ilady(256,256)
       character*80 inputimg
       common /image/ inputimg
         open(1,file='/images/'//inputimg,access='direct',
     + form='unformatted', recl=65600)
         write(*,*) 'reading the inputimage','/images/'//inputimg
         read(1, rec=1) (pim(j), j=1, 65600)read(1, rec=1) (header(j), j=1,64)
         close (1)do 1000 i=1, nx
          k = (i-1) * ny1 = k + 6511=1+255mm=0do 1001 m=1,11
         mm = mm + 1ilady(i, mm) = ichar(pim(m))if(ilady(i, mm).lt.0) ilddy(i, mm)=256+ilady(i, mm)pic(i, mm) = float(ilady(i, mm))1001
         continue
```

```
1000
         continue
       return
       end
```

```
subroutine segma (image, raw, col, output, rmean)
     integer raw, col
     real image(raw, col), output
     double precision s1, s2
     s1=0.0s2=0.0do 10 i=1, raw
         do 10 j=1, col
          s1 = s1 + image(i, j)s2 = s2 + (image(i, j)) **210<sub>1</sub>continue
     rmean = s1/(raw*col)output=s2/(raw*col)-(rmean)**2
     return
     end
```

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```
subroutine m error(image1,image2,raw,col,output)
     integer raw, col
     real image1(raw, col), image2(raw, col), output
     s = 0.0do 10 i=1, raw
        do 10 j=1, cols = s + (image1(i, j) - image2(i, j))**2
10continue
     output=s/(raw*col)
     return
     end
```
subroutine write(pic, nx, ny) real $pic(nx, ny)$

```
character*1 pim(65600), header(64)
      common /header/ header
        open(1,file='image.img',access='direct',
    + form='unformatted', recl=65600)
      do 10 i=1,64pim(i) = header(i)10continue
      do 20 i=1, nxdo 20 j=1, nyip=int(pic(i,j))if(ip.get.128) ip=ip-256
            mm = 64 + j + (i - 1) * nypim(mm) = char(ip)20<sub>o</sub>continue
      write(1, rec=1) (pim(j), j=1, 65600)close (1)return
      end
```

```
\mathbf{C}
```

```
subroutine writeimg(pic, nx, ny, name)
      real pic(nx, ny)character*1 pim(65600)
      character*20 name
        open(1,file=name,access='direct',
    + form='unformatted', recl=nx*ny)
      do 20 i=1, nxdo 20 j=1, ny
          ip=int(pic(i,j))if(ip.get.128) ip=ip-256mm = j + (i - 1) * nypim(mm) = char(ip)20
      continue
      write(1, rec=1) (pim(j), j=1, nx*ny)close (1)return
      end
       function xlog2(x)
```

```
real x
xlog2 = log(x)/alog(2.0)return
end
```
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```
c quantizer subroutine
c for odd number of levels (3 to 35), 0, 64 and 128 levels
         subroutine quantizer(a, nb, c, n, mean, var2, level, raw, col)c a:real image
c c: integer quantized image
c n:quantizer order
c nb:band number
c mean: image mean
c var: image variance
c level, raw, col: output array dimension
         integer raw, col, level
        real a(raw,col), p(75), p1(75), mean, var
         integer c(level, raw, col)
\mathbf Cif(n.eq.0) thendo 4 i=1, raw
           do 4 \neq -1, col
             c(nb,i,j)=04
           continue
        go to 44
        endif
\mathbf{C}m=ncall readq(n, p, pl)if(n.ne.m) thenwrite(*,*) 'error in reading quantizer data'
        stop
        end if
         nn = n/2var=sqrt(var2)
        do 10 i=1, raw
           do 10 j=1, col
               xx=(a(i,j)-mean)do 20 k=nn, 1, -1px=p(k)*varif(xx.get.py) then
                 c(nb,i,j)=kgo to 100
               else if(xx.It.(-px)) then
                 c(nb,i,j) = -kgo to 100
               endif
20
        continue
        c(nb,i,j)=0100
        continue
        continue
10<sup>1</sup>
```
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```
44
         continue
         return
         end
\mathbf C\mathbf{C}dequantizer subroutine
\mathbf{C}c for odd number of levels (3 to 35), 0, 64 and 128 levels
         subroutine dequantizer(b, nb, ic, n, mean, var2, level, raw, col)
c b:real quantized image
c ic: integer quantized image 3D
c n: quantizer order
c nb:band number
c mean: image mean
c var: image variance
         integer raw, col, level
         real b(raw,col),p(75),p1(75), mean, var
         integer ic(level, raw, col)
\mathbf Cif(n.eq.0) thendo 4 i=1, raw
           do 4 j=1, col
             b(i,j)=mean
 \overline{4}continue
         go to 44
         endif
\mathbf{C}m=ncall readq(n, p, pl)if(n.ne.m) thenwrite(*,*) 'error in reading quantizer data'
         stop
         end if
         var=sqrt(var2)
         do 10 i=1, raw
           do 10 j=1, col
            k=ic(nb,i,j)if(k.get.0)then
              b(i,j) = pl(k) * var + meanelse if(k.lt.0)then
              k=-kb(i,j) = -pl(k) * var + meanelse if(k.eq.0)then
              b(i,j)=mean
        endif
10
        continue
44
        continue
```

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103
```

```
return
end
```

```
\mathbf C\mathbf Cc quantizer subroutine
c for odd number of levels (3 to 35), 0, 64 and 128 levels
        subroutine quant(a,b,c,n,mean,var,raw,col)
c a:real image
c b:real quantized image
c c:integer quantized image
c n:quantizer order
c mean: image mean
c var: image variance
         integer raw, col
        real a(raw,col), b(raw,col), p(75), p1(75), mean, var
         integer c(raw,col)\mathbf Cif(n.eq.0)then
           do 4 \neq i=1, raw
           do 4 j=1, colb(i,j)=mean
             c(i,j)=0\overline{4}continue
        go to 44
        endif
\mathbf Cm=ncall readq(n, p, pl)if(n.ne.m)then
        write(*,*) 'error in reading quantizer data'
        stop
        end if
         nn = n/2var=sqrt(var)
        do 10 i=1, raw
          do 10 j=1, col
               xx=(a(i,j)-mean)do 20 k=nn, 1, -1px=p(k)*varif(xx.get.py) then
                 b(i,j)=(pl(k)*var)+mean
                 c(i,j)=kgo to 100
               else if(xx.lt.(-px)) then
                 b(i,j) = (-pl(k) * var) + meanc(i,j) = -kgo to 100
```

```
20<sub>o</sub>continue
         b(i,j)=0+meanc(i, j) = 0100
         continue
 10<sup>°</sup>continue
 44
         continue
         return
         end
\mathbf{C}c read subroutine which
c read quantizer data
         subroutine readq(n, p, pl)real p(75), p1(75)open(32, file='../quantizer.dir/quant2.dat'
               , access='direct', form='unformatted', recl=1000)
      \deltann=(n)/2if(n.le.35.and.n.ge.3) then
            number=n/2+1else if(n.eq.64) then
            number=20else if(n.eq.128) then
            number=21else
            write(*,*) 'the order of the filter is rong'
            STOP
         endif
         read(32, rec=number) n, (p(i), i=1, nn), (p1(i), i=1, nn)
         close(32)return
         end
\mathbf C\mathbf C\mathbf{C}c This subroutine allocate the bitrate for each subband
c by finding the suitable quantizers
c Number of band is 64
\mathbf Csubroutine entropy (iq, var, mean)
         parameter(nn=64)
         real var(nn)
         real mean(nn), br(nn)integer iq(nn)
```
endif

double precision xx

```
logical ql(nn)
         common /bitrate/ temprav
         prav=temprav
\mathbf{C}\mathbf{C}check for negative br and cancell the corespondant band
\mathbf{C}do 10 i=1, nnql(i)=.true.10continue
 100
         continue
         call gain(var, ql, xx, nn, ncount)
         rav=prav*nn/real(ncount)
         do 30 i=nn, 1, -1if(q1(i)) thenx1=var(i)/xxbr(i)=rav+0.5*xloq2(x1)if(br(i).It.0) thenbr(i)=0ql(i)=. false.
                 iq(i)=0go to 100
              endif
            endif
         continue
 30
\mathbf Cc optimize the bit rates
\mathbf C200
         continue
\mathbf Cwrite(*, *) 'rav=', rav
         call gain(var, ql, xx, nn, ncount)
        write(*, *) xx
\mathbf Cdo 40 i=nn, 1, -1if(ql(i)) thenx1=var(i)/xxbr(i)=rav+0.5*xlog2(x1)in=int(2.0**br(i))iq(i) = (in+1)/2*2-1if(iq(i).lt.3) theniq(i)=0else if(iq(i).ge.128) then
                  iq(i) = 128else if(iq(i).ge.64) then
                  iq(i) = 64else if(iq(i).ge.35) then
                  iq(i)=35endif
             ql(i)=. false.
             if(ncount.gt.2)then
                if(iq(i).ge.3) then
```

```
xiq=real(iq(i))xr = xlog2(xiq)else
                   xr=0endif
                  xc=real(ncount)
                  xrav=rav
                  rav=(xrav*xc-xr)/(xc-1.0)\mathbf{C}write(55,*) br(i), iq(i), xr, var(i), ravgo to 200
             endif
           endif
         continue
 40
c calculate the actuall bit rate
        sum=0.0do 50 i=1, nnwrite(55, *) i, iq(i), var(i), mean(i)\overline{c}if(iq(i).gt.1) thenbr(i)=xlog2 (real(iq(i)))else
              br(i)=0endif
            sum = sum + br(i)50
        continue
        rav=sum/nn\mathbf Cwrite(55,*) 'rav =', rav
        return
        end
\mathbf C\mathbf Cc This subroutine calculates the gain
        subroutine gain(var, ql, xx, nband, ncount)
        real var(nband)
        logical ql(nband)
        double precision xx
        xx=1.0count=1.0do 10 i=1, nband
           if(ql(i))then
            xx=xx*var(i)count=count+1
          endif
  10
        continue
        XX=XX** (1.0/count)
        ncount=int(count+.5)
```
return end

```
\mathbf{C}c this subroutine calculate the entropy of each band
c and find the probability of each code
         subroutine bitrates(ic, ln, entropy, pr, level, raw, col)
         integer ic(level, raw, col), ln, level, raw, col
         real entropy(level), sum(-64:64), pr(-64:64)do 20 n=-64,64sum(n)=0.020
         continue
        do 10 i=1, raw
        do 10 j=1, colk=ic(ln, i, j)sum(k) = sum(k) + 110continue
         entropy(1n)=0.0total=real(raw*col)
        do 30 n=-64,64pr(n) = sum(n) / totalif(pr(n).qt.0) thenbr=pr(n)*xlog2(1.0/pr(n))entropy(ln) = entropy(ln) + brendif
 30
        continue
        write(*,*) 'entropy = ', entropy
\mathbf Creturn
        end
\mathbf{C}
```
subroutine convert(pr, prb, prnum, ns)

```
real pr(-64:64), prb(128)integer prnum(128), ns
n=0do 10 i=-64,64if(pr(i).qt.0.0) thenn=n+1prb(n)=pr(i)prnum(n)=i
```

```
endif
 10 continue
     ns=nreturn
     end
C
\overline{\mathbf{C}}c Program to sort the probability table
\mathbf{C}subroutine sort (pr, prnum, ns)
         real pr(128), temppr
         integer prnum(128), tempprn
\mathbf Cc sorting algorithm
         do 10 k=ns-1, 1, -1do 20 i=1,kif(pr(i).lt,pr(i+1)) then
               temppr=pr(i)pr(i)=pr(i+1)pr(i+1) = tempprtempprn=prnum(i)
              prnum(i) = prnum(i+1)prnum(i+1) = tempprnendif
 20continue
         continue
 10
         return
         end
\mathbf{C}\mathbf{C}c subroutine to find the huffman codebook
\overline{c}subroutine hcbook(prb, codbook, ns)
     real pr(128*129/2), prb(128)
         integer pointer(128*129/2)
     character*1 code(128*129/2)character*80 codbook(128), tempc, codebook
         logical flag
         do 19 k=1,128
19
            codbook(k) ='
         do 50 k=1,nspr(k) = prb(k)
```

```
pointer(k)=0code(k) = 'x'50 continue
        do 51 k=ns+1,128*129/2
          pr(k)=0.0pointer(k)=0code(k) = 'x'51
        continue
          code(ns*(ns+1)/2-1)=1'code(ns*(ns+1)/2-2)='0'last=0do 20 j=1, ns-2
          nn = ns - jxx=pr(last+nn)+pr(last+nn+1)code(last+nn) = '0'code(last+nn+1)=1'flag=.false.
          do 10 i=last+1, last+nn
             if(flag) then
              pr(i+nn+1)=pr(i-1)pointer(i-1)=i+nn+1else
               if(xx.get.pr(i))then
               flag=.true.
                 pr(i+nn+1)=xxpointer(last+nn)=i+nn+1
                 pointer(last+nn+1)=i+nn+1
              else
                 pr(i+nn+1)=pr(i)pointer(i)=i+nn+1endif
            endif
 10
          continue
          last=last+nn+1
 20
        continue
        do 11 i=1,nsk = i\text{codbook}(i) ='
          tempc=''
          do 12 j=1, ns-1
          if(code(k).eq.'0'.or.code(k).eq.'1')then
            write(*, *) code(k)\mathbf Ccodebook=code(k)//tempc
            write(*, *) codebook
Ċ
            tempc=codebook
```
 \mathbf{t}

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```
endif
\mathbf Cwrite(*,*) 'k=', k
            k = pointer(k)write(*,*) 'k=', k
\mathbf{C}12
          continue
          codbook(i) = codebook//'\mathcal{L}11continue
         write(*,*) (pr(i), i=1, 14)<br>write(*,*) (pointer(i), i=1, 14)
\mathbf C\mathbf Cwrite(*,*) (code(i), i=1, 14)
\mathbf C\mathbf Cwrite(51,*) (codbook(i), i=1, 5)return
      end
\mathbf{C}_{\_}c subroutine to code an image to hoffman
         subroutine codehc(image, level, codbook, prnum, ns, hc, 11, lev
                                , raw, col)
      s.
         integer lev, raw, col
         integer image(lev,raw,col),prnum(128),ns
         character*80 codbook(128), ccode
         character*1 hc(77000)
         do 5 i=1,77000hc(i) ='5
      continue
         ncount=0do 10 i=1, raw
         do 10 j=1, col
            icode = image(level, i, j)do 20 k=1,nsif (prnum(k).eq.icode) thenncode=k
                  qo to 21
               endif
 20
            continue
            write(*,*) 'ERROR IN READING CODBOOK'
            stop
21continue
```

```
111
```

```
length=index(codbook(ncode),'')-1
           ccode=codbook(ncode)
           do 30 k=1, length
             ncount=ncount+1
             hc (ncount) = ccode (k: k)
 30
           continue
 10<sup>°</sup>continue
         11=ncount
         do 50 il=1,77000
         if(hc(il).ne.'0'.and.hc(il).ne.'1') go to 60
 50 continue
 60
         112 = i1 - 1if(112.ne.11) thenwrite(*,*) ' length of hc is not correct', 112
         endif
         return
         end
\mathbf{C} and \mathbf{C}c program to decode the hoffman code
\overline{c}subroutine decodehc (image, level, codbook, prnum, ns, hc, 11,
     s.
                                lev, raw, col)
         integer lev, col, raw
         integer image(lev,raw,col),icode(128*128),prnum(128),ns
         character*80 codbook(128), ccode
         character*1 hc(77000), code1, code2
         logical flag(128)
         npexel=0
         ncount=0
         do 5 i=1, ns
           flag(i) = true.if(flag(i)) thenwrite(76, *) 'true'
\mathbf Cendif
 5
         continue
         do 10 ii=1,11
           code1=hc(ii)
```
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```
subroutine unifquan(pic, 11, image, order, min, step, level, raw, co
integer raw, col
integer image(level, raw, col), order
real pic(raw, col)
real min, max
```

```
assume the minimum value and the maximum
\mathbf Cmin =pic(1,1)
```

```
max = pic(1, 1)check if all pixels values are between min and max
\mathbf{C}do 10 i=1, raw
          do 10 j=1, col
             if(pic(i,j).gt.max) thenmax = pic(i, j)else if(pic(i,j)).It.min)then
               min = pic(i, j)end if
        continue
 10<sup>°</sup>rescale the image and convert it into image array
\mathbf{C}step = (max-min)/real(order)ixx=order/2
          do 20 i=1, raw
            do 20 j=1, col
               image(11, i, j) = int((pic(i, j) - min)/step+.5) - ixx20
          continue
         return
         end
\mathbf{C}subroutine unifdequan(pic, 11, image, order, min, step, level, raw,
        integer raw.col
        integer image(level, raw, col), order
        real pic(raw, col)
        real step, min
    rescale the image and convert it into pic array
\mathbf{C}ixx=order/2
          do 20 i=1, raw
            do 20 j=1, col
             pic(i, j) = real(image(ll, i, j) + ixx) * step + mincontinue
 20
         return
         end
\mathbf{C}
```
 \bar{V}

```
subroutine dpcdecoder(e, img, raw, col)
         parameter(ii=32, jj=32)integer raw, col
         real e(raw, col), img(raw, col)
         real temp(0:ii+1,0:jj+1)do 10 i=0, raw+1
         do 10 j=0, col+1
              temp(i, j) = 0continue
 10<sup>°</sup>do 20 i=1, raw
         do 20 j=1, col
             temp(i, j) = e(i, j) + 0.25*(temp(i, j-1) + temp(i-1, j-1) +\deltatemp(i-1,j) + temp(i-1,j+1))
             img(i, j) = temp(i, j)20
     continue
         return
     end
\mathbf{C}_{\perp}\mathbf{C}c program for noise less dpcm coding algorithm
\mathbf Csubroutine qdpcm(u, ie, iq, var, mean, raw, col, p, pl)
         parameter(ii=32, jj=32)integer raw, col
     real u(\text{raw}, \text{col}), ud(0:i i+1, 0:j j+1), udbreal e, ed
         real p(75), p1(75)integer ie(4, raw, col)real mean, var
         do 10 i=0, raw+1
         do 10 j=0, raw+1
          ud(i,j)=0.010 continue
           a1=0.25a2=0.25a3=0.25a4=0.25do 20 m=1, raw
     do 20 n=1, col
          udb=a1*ud(m-1,n)+a2*ud(m,n-1)+a3*ud(m-1,n-1)+a4*ud(m-1,n+1
          e=u(m,n)-udbcall quantizer1(e,ed,ie(1,m,n),iq,mean,var,p,pl)
```

```
ud(m, n) = udb + ed20 continue
      return
      end
\mathbf{C}_{\perp}\mathbf{c}c quantizer subroutine
c for odd number of levels (3 to 35), 0, 64 and 128 levels
         subroutine quantizer1(a,b,c,n,mean,var2,p,pl)
c a:real image
c b: real quantized image
c c: integer quantized
c n:quantizer order
c mean: image mean
c var: image variance
\mathbf Creal a, p(75), p1(75), mean, var
         integer c
C
         if(n.eq.0) thenc=0go to 44
         endif
\mathbf Cnn = n/2var=sqrt(var2)
                xx=(a-mean)do 20 k=nn, 1, -1px=p(k)*varif(xx.get.py) then
                   b = (p1(k) * var) + meanc = kgo to 100
                else if(xx.lt.(-px))then
                  b = (-p1(k) * var) + meanc=-kgo to 100
                endif
 20
         continue
         b=mean
         c=0100
         continue
 44
         continue
         return
         end
\mathbf C\mathbf C
```
 $\overline{1}$

```
c program for dpcm coding algorithm
     subroutine dpcm1(x,y,raw, col)parameter(ii=32, jj=32)integer raw, col
     real x(raw, col), temp(0:ii+1, 0:ii+1), y(raw, col)write(*,*) 'dpcml'
         do 10 i=0, i:i+1do 10 j=0,jj+1if(i.eq.0.or.j.eq.0.or.i.gt.ii.or.j.gt.jj) then
                temp(i, j) = 0else
                temp(i, j) = x(i, j)endif
    continue
 10<sup>°</sup>do 20 i=1, raw
     do 20 j=1, col
       y(i,j) = temp(i,j) - 0.25*(temp(i-1,j) + temp(i,j-1))+temp(i-1,j-1)+temp(i-1,j+1))
     s.
     continue
 20
     return
     end
\mathbf{C}_{\_}\Lambda
```
 $\ddot{}$

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